Problem Set 7

514 - Networks and Combinatorial Optimization

Autumn 2023

Exercise 7.1 (10pts)

Let G = (V, E) be a 3-regular graph without any bridge. Show that *G* has a perfect matching. (A *bridge* is an edge *e* not contained in any circuit; equivalently, deleting *e* increases the number of components; equivalently $\{e\}$ is a cut. A graph *G* is *k*-regular if all vertices have degree *k*.)

Exercise 7.2 (10pts)

Let G = (V, E) be a graph and let $T \subseteq V$. Then *G* has a matching covering *T* if and only if the number of odd components of $G \setminus W$ contained in *T* is at most |W|, for each $W \subseteq V$.

Exercise 7.3 (10pts)

Recall that for a graph G = (V, E) we have defined $P_{\text{matching}}(G) = \text{conv}\{\chi^M \in \mathbb{R}^E \mid M \subseteq E \text{ is a matching}\}$. Prove that for any $k \in \mathbb{N}$,

 $P_{\text{matching}}(G) \cap \{x \in \mathbb{R}^E \mid \mathbf{1}^T x = k\} = \text{conv}\{\chi^M \mid M \subseteq E \text{ is a matching with } |M| = k\}$

Hint: Let $\mathcal{F} := \{M \subseteq E \mid M \text{ is a matching}\}$. For the non-trivial direction, take a vector $x^* \in P_{\text{matching}}(G) \cap \{x \in \mathbb{R}^E \mid \mathbf{1}^T x = k\}$ and consider the vector $\lambda \in \mathbb{R}_{\geq 0}^{\mathcal{F}}$ that minimizes the function $G(\lambda) := \sum_{M \in \mathcal{F}} \lambda_M \cdot |M| - k|$ subject to $\sum_{M \in \mathcal{F}} \lambda_M = 1$ and $x^* = \sum_{M \in \mathcal{F}} \lambda_M \chi^M$.

Remark. All three exercises are taken from A. Schrijver's lecture notes.