# OPTIMIZATION AND PSEUDOSPECTRA, WITH APPLICATIONS TO ROBUST STABILITY* 

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#### Abstract

The $\epsilon$-pseudospectrum of a matrix $A$ is the subset of the complex plane consisting of all eigenvalues of all complex matrices within a distance $\epsilon$ of $A$. We are interested in two aspects of "optimization and pseudospectra." The first concerns maximizing the function "real part" over an $\epsilon$-pseudospectrum of a fixed matrix: this defines a function known as the $\epsilon$-pseudospectral abscissa of a matrix. We present a bisection algorithm to compute this function. Our second interest is in minimizing the $\epsilon$-pseudospectral abscissa over a set of feasible matrices. A prerequisite for local optimization of this function is an understanding of its variational properties, the study of which is the main focus of the paper. We show that, in a neighborhood of any nonderogatory matrix, the $\epsilon$-pseudospectral abscissa is a nonsmooth but locally Lipschitz and subdifferentially regular function for sufficiently small $\epsilon$; in fact, it can be expressed locally as the maximum of a finite number of smooth functions. Along the way we obtain an eigenvalue perturbation result: near a nonderogatory matrix, the eigenvalues satisfy a Hölder continuity property on matrix space-a property that is well known when only a single perturbation parameter is considered. The pseudospectral abscissa is a powerful modeling tool: not only is it a robust measure of stability, but it also reveals the transient (as opposed to asymptotic) behavior of associated dynamical systems.


Key words. pseudospectrum, eigenvalue optimization, spectral abscissa, nonsmooth analysis, subdifferential regularity, robust optimization, robust control, stability radius, distance to instability, $\mathbf{H}_{\infty}$ norm

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1. Introduction. The $\epsilon$-pseudospectrum of a matrix $A$, denoted $\Lambda_{\epsilon}(A)$, is the subset of the complex plane consisting of all eigenvalues of all complex matrices within a distance $\epsilon$ of $A$ (see $[20,39,40]$ ). We are interested in two aspects of "optimization and pseudospectra." The first concerns maximizing a simple real-valued function over a fixed pseudospectrum $\Lambda_{\epsilon}(A)$. We focus specifically on the case where this function is simply "real part." Then the optimal value defines the $\epsilon$-pseudospectral abscissa of $A$, denoted $\alpha_{\epsilon}(A)$. Just as the spectral abscissa of a matrix provides a measure of its stability, that is, the asymptotic decay of associated dynamical systems, so the $\epsilon$-pseudospectral abscissa provides a measure of robust stability, where by robust we mean with respect to complex perturbations in the matrix. One of the contributions of this paper is a bisection algorithm that computes $\alpha_{\epsilon}(A)$ for any $A$; this algorithm also identifies all maximizing points in the pseudospectrum.

In many applications, matrices are not fixed but dependent on parameters that may be adjusted. Our second interest in optimization concerns minimizing the $\epsilon$ pseudospectral abscissa $\alpha_{\epsilon}$ over a feasible set of matrices. A prerequisite for local

[^0]minimization of $\alpha_{\epsilon}$ is an understanding of its variational properties as a function of the matrix $A$. This provides the focus for most of the paper. Our main result shows that, in a neighborhood of any nonderogatory matrix, the $\epsilon$-pseudospectral abscissa is a nonsmooth but locally Lipschitz and subdifferentially regular function for sufficiently small $\epsilon$; in fact, it can be expressed locally as the maximum of a finite number of smooth functions. Such a property is desirable from the point of view of numerical methods for local optimization, but we defer a computational study to future work.

The paper is organized as follows. After setting up some notation in section 2, we begin in section 3 by discussing related ideas in the robust control literature. We review the connections between the pseudospectral abscissa and the "distance to instability" [28, 16], or "complex stability radius" [21], and the $\mathbf{H}_{\infty}$ norm of a transfer function [8]. The outcome of minimization of $\alpha_{\epsilon}$ over a set of matrices obviously depends on the crucial issue of the choice of $\epsilon$. We show that as $\epsilon$ is increased from zero to an arbitrarily large quantity, the corresponding optimization problem evolves from minimization of the spectral abscissa (enhancing the asymptotic decay rate of the associated dynamical system) to the minimization of the largest eigenvalue of the symmetric part of the matrix (minimizing the initial growth rate of the associated system). Regarding the first of these extremes (optimization of the spectral abscissa), variational analysis of this non-Lipschitz function is well understood [15, 12], global optimization is known to be hard [5], and some progress has been made in local optimization methods [14]. Regarding the second extreme, optimization of this convex function over a polyhedral feasible set is a semidefinite programming problem, and the global minimum can be found by standard methods [4, 38]. Intermediate choices of $\epsilon$ control transient peaking in the dynamical system associated with the matrix, and one particular choice corresponds exactly to the complex stability radius (or $\mathbf{H}_{\infty}$ norm) optimization problem. Thus the pseudospectral approach gives a whole range of stabilizing optimization problems, each with a quantifiable interpretation in terms of the allowable perturbations. Furthermore, unlike maximization of the complex stability radius, which simply optimizes the "robustness" of the stability, minimizing the pseudospectral abscissa preserves some explicit emphasis on optimizing the asymptotic decay rate of the system.

In section 4, we analyze the topology of the pseudospectrum, observing that points on the boundary are accessible from the interior by analytic paths, and discussing conditions under which the boundary is differentiable at points that maximize the real part. This sets the stage for the description of a simple bisection algorithm to compute $\alpha_{\epsilon}(A)$, the pseudospectral abscissa of a fixed matrix, in section 5 . In section 6 , we show that the bisection algorithm locates all maximizers of the real part over the pseudospectrum. The bisection algorithm is very much analogous to Byers' algorithm for measuring the distance to instability [16], which has spawned more sophisticated variants for the calculation of stability radii (real as well as complex) and $\mathbf{H}_{\infty}$ norms, both globally and quadratically (or higher order) convergent; see $[8,7,11,18,36]$. Along similar lines, we have also developed a quadratically convergent variant algorithm for computing $\alpha_{\epsilon}$, described and analyzed in a companion paper [13].

In section 7, we continue our study of analytical properties of the pseudospectrum. As is well known, the pseudospectrum of a matrix $A$ is defined by an inequality on $\sigma_{\min }(A-z I)$, the least singular value of $A-z I$. In Theorem 7.4 (growth near an eigenvalue) we give an interesting estimate relating $\sigma_{\min }(A-z I)$ to $\left|z-\lambda_{0}\right|^{m}$, where
$\lambda_{0}$ is a nonderogatory eigenvalue (one whose geometric multiplicity is one), and $m$ is its algebraic multiplicity. The coefficient relating these quantities is a ratio of two products, of eigenvalue separations and of singular values, respectively. One corollary of this result is that an $\epsilon$-pseudospectrum component around a nonderogatory eigenvalue is strictly convex for sufficiently small $\epsilon$, an intuitively appealing but apparently nontrivial fact. Another corollary is that the nonderogatory eigenvalue $\lambda_{0}$ satisfies a Hölder continuity property in a neighborhood of $A$ in matrix space, with Hölder exponent equal to $1 / \mathrm{m}$. While this result might not seem surprising, in light of wellknown classical spectral perturbation theory [27, 23, 3, 30], we have not seen it in the literature. The classical analysis focuses almost exclusively on single perturbation parameters.

The analytical results of section 7 allow us to achieve our primary goal in section 8: a detailed variational analysis of the pseudospectral abscissa $\alpha_{\epsilon}$. The main result has already been mentioned above. Finally, in section 9, we examine the boundary properties of the pseudospectrum at points where the boundary is not smooth, using techniques from modern variational analysis [17, 34]. We show that, under a nondegeneracy condition, the complement of the pseudospectrum is Clarke regular at such a point, and give a formula for the normal cone.
2. Notation. We consider a matrix $A$ in the space of $n \times n$ complex matrices $\mathbf{M}^{n}$. We denote the spectrum of $A$ by $\Lambda=\Lambda(A)$, and we denote by $\alpha=\alpha(A)$ the spectral abscissa of $A$, which is the largest of the real parts of the eigenvalues.

For a real $\epsilon>0$, the $\epsilon$-pseudospectrum of $A$ is the set

$$
\Lambda_{\epsilon}=\{z \in \mathbf{C}: z \in \Lambda(X) \text { where }\|X-A\| \leq \epsilon\}
$$

(Throughout, $\|\cdot\|$ denotes the operator 2-norm on $\mathbf{M}^{n}$.) For the most part, $\epsilon$ is fixed, so where it is understood we drop it from the terminology. Any element of the pseudospectrum is called a pseudoeigenvalue. Unless otherwise stated, we shall always assume $\epsilon>0$, but it is occasionally helpful to extend our notation to allow $\epsilon=0$, so $\Lambda_{0}=\Lambda$. Analogously, the strict pseudospectrum is the set

$$
\Lambda_{\epsilon}^{\prime}=\{z \in \mathbf{C}: z \in \Lambda(X) \text { where }\|X-A\|<\epsilon\}
$$

The pseudospectral abscissa $\alpha_{\epsilon}$ is the maximum value of the real part over the pseudospectrum:

$$
\begin{equation*}
\alpha_{\epsilon}=\sup \left\{\operatorname{Re} z: z \in \Lambda_{\epsilon}\right\} \tag{2.1}
\end{equation*}
$$

We call this optimization problem the pseudospectral abscissa problem. Note $\alpha_{0}=\alpha$.
The function $\sigma_{\min }: \mathbf{M}^{n} \rightarrow \mathbf{R}$ denotes the smallest singular value. We define a function $g: \mathbf{C} \rightarrow \mathbf{R}$ by

$$
g(z)=\sigma_{\min }(A-z I)=\left\|(A-z I)^{-1}\right\|^{-1}
$$

where we interpret the right-hand side as zero when $z \in \Lambda(A)$. Thus $g$ is the reciprocal of the norm of the resolvent. Using this notation, a useful characterization of the pseudospectrum is

$$
\Lambda_{\epsilon}=\{z \in \mathbf{C}: g(z) \leq \epsilon\}
$$

and analogously

$$
\Lambda_{\epsilon}^{\prime}=\{z \in \mathbf{C}: g(z)<\epsilon\}
$$

(see [39]). Clearly as $\epsilon$ increases, both families of sets are monotonic increasing.
We will sometimes want to allow the matrix $A$ (and the parameter $\epsilon$ ) to vary. We therefore define the pseudospectral abscissa function $\alpha_{\epsilon}: \mathbf{M}^{n} \rightarrow \mathbf{R}$ by

$$
\alpha_{\epsilon}(Z)=\sup \left\{\operatorname{Re} z: \sigma_{\min }(Z-z I) \leq \epsilon\right\} .
$$

3. Related ideas. The pseudospectral abscissa is related to several other functions important for stability analysis. In this section we briefly sketch the connections with two such functions, in particular, the "distance to instability" and the $\mathbf{H}_{\infty}$ norm.

A matrix $A$ is stable if all its eigenvalues have strictly negative real parts; in other words, the spectral abscissa of $A$ satisfies $\alpha(A)<0$. From any given matrix $A$, the distance to the set of matrices which are not stable [28, 19] (also known as the complex stability radius [21]) is

$$
\beta(A)=\min \left\{\|X-A\|: X \in \mathbf{M}^{n}, \alpha(X) \geq 0\right\} .
$$

Since the set of matrices which are not stable is closed, this minimum is attained. Notice in particular that $\beta(A)=0$ if and only if $A$ is not stable. It is now easy to check the relationship

$$
\begin{equation*}
\beta(A) \leq \epsilon \quad \Leftrightarrow \quad \alpha_{\epsilon}(A) \geq 0, \tag{3.1}
\end{equation*}
$$

and more generally, for any real $x$,

$$
\alpha_{\epsilon}(A) \geq x \quad \Leftrightarrow \quad \alpha_{\epsilon}(A-x I) \geq 0 \quad \Leftrightarrow \quad \beta(A-x I) \leq \epsilon .
$$

Notice that we can write the pseudospectral abscissa in the form

$$
\alpha_{\epsilon}(A)=\max \{\alpha(X):\|X-A\| \leq \epsilon\},
$$

a special case of "robust regularization" [26] and "minimum stability degree" [2]. Since the spectral abscissa $\alpha$ is continuous, standard arguments [26] show that the function

$$
\begin{equation*}
(\epsilon, A) \in \mathbf{R}_{+} \times \mathbf{M}^{n} \mapsto \alpha_{\epsilon}(A) \tag{3.2}
\end{equation*}
$$

is continuous.
In this paper we consider almost exclusively a fixed choice of the parameter $\epsilon$, but for the moment let us consider the effect of varying $\epsilon$ on the solution of a pseudospectral abscissa minimization problem. For any fixed set of feasible matrices $F \subset \mathbf{M}^{n}$, the continuity of the map (3.2) guarantees various useful continuity properties of the optimal value and solutions of the optimization problem $\inf _{F} \alpha_{\epsilon}$ (see [34, Chap. 7]). In particular, if $F$ is nonempty and compact, then

$$
\lim _{\epsilon \rightarrow \bar{\epsilon}} \inf _{F} \alpha_{\epsilon}=\inf _{F} \alpha_{\bar{\epsilon}},
$$

and any cluster point of a sequence of matrices $A_{r}$ minimizing $\alpha_{\epsilon_{r}}$ over $F$, where $\epsilon_{r} \rightarrow \bar{\epsilon}$, must minimize $\alpha_{\bar{\epsilon}}$ over $F$.

Notice that any stable matrix $A$ satisfies

$$
\alpha_{\beta(A)}(A)=0 .
$$

To see this, note that the implication (3.1) shows $\alpha_{\beta(A)} \geq 0$, while if $\alpha_{\beta(A)}>0$, then by the continuity of $\alpha_{\epsilon}$ with respect to $\epsilon$, there would exist $\epsilon \in(0, \beta(A))$ such that $\alpha_{\epsilon}(A) \geq 0$, whence we get the contradiction $\beta(A) \leq \epsilon<\beta(A)$.

We return to our pseudospectral abscissa minimization problem $\inf _{F} \alpha_{\epsilon}$. The following easy result shows that, under reasonable conditions, for a particular choice of $\epsilon$, this problem is equivalent to maximizing the distance to instability over the same set of feasible matrices.

Proposition 3.1 (maximizing the distance to instability). If the optimal value $\bar{\beta}=\max _{F} \beta$ is attained by some stable matrix, then $\min _{F} \alpha_{\bar{\beta}}=0$ and

$$
\operatorname{argmin}\left\{\alpha_{\bar{\beta}}(X): X \in F\right\}=\operatorname{argmax}\{\beta(X): X \in F\} .
$$

Proof. Any matrix $A \in F$ satisfies $\beta(A) \leq \bar{\beta}$. If $A$ is stable, then $\alpha_{\bar{\beta}}(A) \geq$ $\alpha_{\beta(A)}(A)=0$, while on the other hand, if $A$ is not stable, then $\alpha_{\bar{\beta}}(A) \geq \alpha_{0}(A) \geq 0$. Hence $\inf _{F} \alpha_{\bar{\beta}} \geq 0$.

By assumption, $\bar{\beta}$ is finite and strictly positive, so clearly every matrix in the (nonempty) set of optimal solutions $\operatorname{argmax}_{F} \beta$ is stable. Any such matrix $A$ satisfies $\alpha_{\bar{\beta}}(A)=\alpha_{\beta(A)}(A)=0$, and hence $A \in \operatorname{argmin}_{F} \alpha_{\bar{\beta}}$. We deduce $\operatorname{argmax}_{F} \beta \subset$ $\operatorname{argmin}_{F} \alpha_{\bar{\beta}}$ and $\min _{F} \alpha_{\bar{\beta}}=0$.

Consider, conversely, a matrix $A \in F$ such that $A \notin \operatorname{argmax}_{F} \beta$. Suppose first that $A$ is stable. Since $\beta(A)<\bar{\beta}$, we know $\alpha_{\bar{\beta}}(A)>\alpha_{\beta(A)}(A)=0$, because as we shall see in the next section, $\alpha_{\epsilon}(A)$ is strictly increasing in $\epsilon$. On the other hand, if $A$ is not stable, then the same reasoning shows $\alpha_{\bar{\beta}}(A)>\alpha_{0}(A) \geq 0$. In either case, we have shown $A \notin \operatorname{argmin}_{F} \alpha_{\bar{\beta}}$, so $\operatorname{argmin}_{F} \alpha_{\bar{\beta}} \subset \operatorname{argmax}_{F} \beta$ as required.

We thus see that, under reasonable conditions, as $\epsilon$ increases from zero, the set of optimal solutions $\operatorname{argmax}_{F} \alpha_{\epsilon}$ evolves from the set of minimizers of the spectral abscissa through the set of maximizers of the stability radius. This raises the question of what happens for large $\epsilon$. The following result shows that the limiting version of $\inf _{F} \alpha_{\epsilon}$ as $\epsilon \rightarrow+\infty$ is the optimization problem

$$
\inf _{X \in F} \lambda_{\max }\left(\frac{X+X^{*}}{2}\right),
$$

where $\lambda_{\text {max }}$ denotes the largest eigenvalue of a Hermitian matrix.
Theorem 3.2 (large $\epsilon$ ). For any matrix $A \in \mathbf{M}^{n}$,

$$
\left[\alpha_{\epsilon}(X)-\epsilon\right] \rightarrow \lambda_{\max }\left(\frac{A+A^{*}}{2}\right) \text { as } \epsilon \rightarrow+\infty \text { and } X \rightarrow A \text {. }
$$

Proof. If we denote the right-hand side by $\lambda$, then there is a unit vector $u \in \mathbf{C}^{n}$ satisfying $u^{*}\left(A+A^{*}\right) u=2 \lambda$. Consider any sequence $\epsilon_{r} \rightarrow+\infty$ and $X_{r} \rightarrow A$. Since $\left\|u u^{*}\right\|=1$, we know

$$
\alpha_{\epsilon_{r}}\left(X_{r}\right)-\epsilon_{r} \geq \alpha\left(X_{r}+\epsilon_{r} u u^{*}\right)-\epsilon_{r}=\epsilon_{r}\left(\alpha\left(u u^{*}+\frac{1}{\epsilon_{r}} X_{r}\right)-1\right) .
$$

Now standard perturbation theory [23] shows $\alpha$ is analytic around the matrix $u u^{*}$ with gradient $\nabla \alpha\left(u u^{*}\right)=u u^{*}$, so as $r \rightarrow \infty$, the right-hand side in the above relationship converges to

$$
\operatorname{Re}\left(\operatorname{tr}\left(u u^{*} A\right)\right)=\operatorname{Re} u^{*} A u=\lambda .
$$

We have thus shown

$$
\liminf _{r}\left(\alpha_{\epsilon_{r}}\left(X_{r}\right)-\epsilon_{r}\right) \geq \lambda .
$$

Now suppose

$$
\limsup _{r}\left(\alpha_{\epsilon_{r}}\left(X_{r}\right)-\epsilon_{r}\right)>\lambda
$$

We will derive a contradiction. Without loss of generality, there exists a real $\delta>0$ such that

$$
\alpha_{\epsilon_{r}}\left(X_{r}\right)-\epsilon_{r}>\lambda+\delta \text { for all } r .
$$

For each $r$ we can choose a matrix $D_{r}$ satisfying $\left\|D_{r}\right\| \leq 1$ and

$$
\alpha_{\epsilon_{r}}\left(X_{r}\right)=\alpha\left(X_{r}+\epsilon_{r} D_{r}\right),
$$

and a unit vector $w_{r} \in \mathbf{C}^{n}$ satisfying

$$
\alpha\left(X_{r}+\epsilon_{r} D_{r}\right)=\operatorname{Re}\left(w_{r}^{*}\left(X_{r}+\epsilon_{r} D_{r}\right) w_{r}\right)
$$

Hence

$$
\begin{aligned}
\lambda+\delta & <\operatorname{Re}\left(w_{r}^{*} X_{r} w_{r}\right)+\epsilon_{r}\left(\operatorname{Re}\left(w_{r}^{*} D_{r} w_{r}\right)-1\right) \\
& \leq \operatorname{Re}\left(w_{r}^{*} X_{r} w_{r}\right)=w_{r}^{*}\left(\frac{X_{r}+X_{r}^{*}}{2}\right) w_{r} \\
& \leq \lambda_{\max }\left(\frac{X_{r}+X_{r}^{*}}{2}\right) .
\end{aligned}
$$

But as $r \rightarrow \infty$, the right-hand side above converges to $\lambda$, which is the desired contradiction.

We see from this result that, for example, if the set $F$ is a polyhedron, then the limiting version of the optimization problem $\inf _{F} \alpha_{\epsilon}$ as $\epsilon \rightarrow \infty$ is a computationally straightforward, convex minimization problem, whereas when $\epsilon=0$ the problem may be hard [5].

The idea of the $\mathbf{H}_{\infty}$ norm of a transfer matrix is also closely related to the complex stability radius. Consider the linear time-invariant dynamical system

$$
\dot{p}=A p+u
$$

where $p$ denotes the state vector (in this simple case coinciding with the output) and $u$ denotes the input vector. The "transfer matrix" of this system is the function $H(s)=(s I-A)^{-1}$ (where $s$ is a complex variable). Assuming the matrix $A$ is stable, the corresponding $\mathbf{H}_{\infty}$ norm is defined by

$$
\|H\|_{\infty}=\sup _{\omega \in \mathbf{R}} \sigma_{\max }(H(i \omega))
$$

where $\sigma_{\max }$ denotes the largest singular value. Clearly

$$
\|H\|_{\infty}=\sup _{\omega \in \mathbf{R}} \frac{1}{\sigma_{\min }(A-i \omega I)}
$$

so $\|H\|_{\infty}<\epsilon^{-1}$ if and only if we have

$$
\sigma_{\min }(A-i \omega I)>\epsilon \text { for all } \omega \in \mathbf{R}
$$

As a consequence of Theorem 5.4 below, for example, this is equivalent to $\alpha_{\epsilon}(A)<0$. In summary, for a stable matrix $A$, we have

$$
\begin{equation*}
\alpha_{\epsilon}(A)<0 \Leftrightarrow \beta(A)>\epsilon \Leftrightarrow\|H\|_{\infty}<\frac{1}{\epsilon} \tag{3.3}
\end{equation*}
$$

We can characterize the condition $\alpha_{\epsilon}(A)<x$ analogously in terms of a "shifted" $\mathbf{H}_{\infty}$ norm [9, p. 67].

An important topic in robust control has been the design of controllers which minimize the $\mathbf{H}_{\infty}$ norm [44, 43]. In the language above, this corresponds to choosing the parameters defining the stable matrix $A$ in order to maximize the minimum value of $\sigma_{\min }(A-z I)$ as $z$ varies along the imaginary axis. Our ultimate aim of optimizing the pseudospectral abscissa is related, but rather different, being motivated by the broad idea of robust optimization [4]. We first fix the "level of robustness" $\epsilon$ (precisely the quantity that we seek to maximize in an $\mathbf{H}_{\infty}$ norm problem) and then vary $A$ to move the corresponding pseudospectrum as far as possible to the left in the complex plane. In other words, we try to maximize a real parameter $x$ such that the $\mathbf{H}_{\infty}$ norm corresponding to the shifted matrix $A-x I$ is not more than $\epsilon^{-1}$.

What are the relative merits of different choices of $\epsilon$ in a pseudospectral minimization problem $\inf _{F} \alpha_{\epsilon}$ ? Here we are motivated by Trefethen's well-known viewpoint [39, 40], but we add an optimization "twist." When $\epsilon=0$, optimization amounts to minimizing the spectral abscissa of a matrix $A \in F$, in other words, optimizing the asymptotic rate of decay of trajectories of the dynamical system $\dot{p}=A p$. On the other hand, for large $\epsilon$, by Theorem 3.2 (large $\epsilon$ ), optimization amounts to minimizing $\lambda_{\max }\left(A+A^{*}\right) / 2$. This corresponds to optimizing the initial decay rate of the dynamical system, since at time $t=0$,

$$
\frac{d}{d t} \frac{\|p\|^{2}}{2}=p(0)^{*}\left(\frac{A+A^{*}}{2}\right) p(0) \leq \lambda_{\max }\left(\frac{A+A^{*}}{2}\right)\|p(0)\|^{2}
$$

with equality if $p(0)$ is an eigenvector corresponding to the largest eigenvalue. For intermediate choices of $\epsilon$, minimizing the pseudospectral abscissa balances the two objectives of improving asymptotic stability and restricting the size of transient peaks in the trajectories. In particular, Proposition 3.1 (maximizing the distance to instability) shows that, under reasonable conditions, for some choice of $\epsilon$, minimizing $\alpha_{\epsilon}$ is equivalent to minimizing the $\mathbf{H}_{\infty}$ norm, that is, maximizing the complex stability radius.

To summarize, minimizing the $\mathbf{H}_{\infty}$ norm of a matrix $A$ optimizes the robustness of the stability of the dynamical system $\dot{p}=A p$, but with no explicit reference to its asymptotic decay rate. By minimizing the pseudospectral abscissa $\alpha_{\epsilon}$ instead, for different choices of the parameter $\epsilon$ we obtain a range of different balances between robustness and asymptotic decay, one choice giving exactly the $\mathbf{H}_{\infty}$ norm problem. One could achieve a similar range of balances by minimizing the $\mathbf{H}_{\infty}$ norm corresponding to the shifted matrix $A-x I$ as the real parameter $x$ varies; however, working with $\epsilon$-pseudospectra for fixed $\epsilon$ provides a natural interpretation in terms of allowable perturbations to $A$. Yet another range of balances is achieved by the "robust spectral abscissa" defined in [14].

Just as with the $\mathbf{H}_{\infty}$ norm, the pseudospectral abscissa can be characterized via semidefinite programming. Specifically, by [9, p. 67] or [4, Prop. 4.4.2], a real $x$ satisfies

$$
\alpha_{\epsilon}(A)<x
$$

if and only if there exist reals $\mu<0$ and $\lambda$, and an $n \times n$ positive definite Hermitian matrix $P$ such that the matrix

$$
\left[\begin{array}{cc}
(\mu-\lambda) I+2 x P-A^{*} P-P A & -\epsilon P \\
-\epsilon P & \lambda I
\end{array}\right]
$$

is positive semidefinite. As discussed in [9, pp. 3-4], the power of such semidefinite characterizations derives from their amenability to efficient interior point methods for convex optimization, pioneered in [31]. The disadvantage is the appearance of subsidiary semidefinite matrix variables: if the underlying matrices $A$ are large, and we need to calculate the pseudospectral abscissa for many different matrices (in an optimization routine, for example), involving these subsidiary variables may be prohibitive computationally; see, for example, [33, 14, 42]. For this reason, in this work we consider more direct approaches to the pseudospectral abscissa.
4. Boundary properties. We begin our direct, geometric approach to the pseudospectral abscissa by studying the boundary of the pseudospectrum.

Proposition 4.1 (compactness). The pseudospectrum $\Lambda_{\epsilon}$ is a compact set contained in the ball of radius $\|A\|+\epsilon$. It contains the strict pseudospectrum $\Lambda_{\epsilon}^{\prime}$, which is nonempty and open.

Proof. The strict pseudospectrum is nonempty since it contains the spectrum. It is open since $\sigma_{\min }$, and hence $g$ are continuous. This also shows that the pseudospectrum is closed. For any point $z \in \Lambda_{\epsilon}$ there is a unit vector $u \in \mathbf{C}^{n}$ satisfying $\|(A-z I) u\| \leq \epsilon$. On the other hand, $\|A u\| \leq\|A\|$, so we have the inequality

$$
\begin{equation*}
|z|=\|z u\| \leq\|(A-z I) u\|+\|A u\| \leq\|A\|+\epsilon \tag{4.1}
\end{equation*}
$$

which shows boundedness.
The next result is slightly less immediate.
THEOREM 4.2 (local minima). The only local minimizers of the function

$$
g(z)=\sigma_{\min }(A-z I)
$$

are the eigenvalues of the matrix $A$.
Proof. Suppose the point $z_{0}$ is a local minimizer that is not an eigenvalue. Then $z_{0}$ is a local maximizer of the norm of the resolvent $\left\|(A-z I)^{-1}\right\|$. We can choose unit vectors $u, v \in \mathbf{C}^{n}$ satisfying

$$
\left\|\left(A-z_{0} I\right)^{-1}\right\|=\left|u^{*}\left(A-z_{0} I\right)^{-1} v\right|
$$

and then we have, for all points $z$ close to $z_{0}$, the inequalities

$$
\left|u^{*}(A-z I)^{-1} v\right| \leq\left\|(A-z I)^{-1}\right\| \leq\left\|\left(A-z_{0} I\right)^{-1}\right\|=\left|u^{*}\left(A-z_{0} I\right)^{-1} v\right|
$$

Hence the modulus of the function $u^{*}(A-z I)^{-1} v$ has a local maximum at $z_{0}$. But this contradicts the maximum modulus principle, since this function is analytic and nonconstant near $z_{0}$.

Corollary 4.3 (closure of strict pseudospectrum). The closure of the strict pseudospectrum is the pseudospectrum, so for $\epsilon>0$ the pseudospectral abscissa is

$$
\alpha_{\epsilon}=\sup \left\{\operatorname{Re} z: z \in \Lambda_{\epsilon}^{\prime}\right\}
$$

Proof. A point in the pseudospectrum that is outside the closure of the strict pseudospectrum must be a local minimizer of the function $g$.

An easy exercise now shows that the pseudospectral abscissa $\alpha_{\epsilon}$ is a continuous, strictly increasing function of $\epsilon \in[0,+\infty)$. Note also that, by contrast with the above result, the function $g$ may have local maximizers and, consequently, the strict pseudospectrum may not equal the interior of the pseudospectrum.

We can refine the above corollary with a more delicate argument, showing that we can "access" any point in the pseudospectrum via a smooth path through the strict pseudospectrum.

THEOREM 4.4 (accessibility). Given any point $z_{0}$ in the pseudospectrum, there is a real-analytic path $p:[0,1] \rightarrow \mathbf{C}$ such that $p(0)=z_{0}$ and $p(t)$ lies in the strict pseudospectrum for all $t \in(0,1]$.

Proof. We may as well assume $g\left(z_{0}\right)=\epsilon$. By Corollary 4.3, there exists a sequence $z_{r} \in \Lambda_{\epsilon}^{\prime}$ approaching $z_{0}$. For each index $r$ there exists a vector $u^{r} \in \mathbf{C}^{n}$ satisfying the inequalities

$$
1<\left\|u^{r}\right\|<1+\frac{1}{r} \text { and }\left\|\left(A-z_{r} I\right) u^{r}\right\|<\epsilon
$$

By taking a subsequence, we may as well assume that the sequence $\left\{u^{r}\right\}$ converges to a limit $u^{0}$, and then we have $\left(z_{0}, u^{0}\right) \in \operatorname{cl} S$, where

$$
S=\left\{(z, u):\|u\|^{2}>1,\|(A-z I) u\|^{2}<\epsilon^{2}\right\}
$$

Since the set $S$ is defined by a finite number of strict algebraic inequalities, we can apply the accessibility lemma [29]. Hence there is a real-analytic path $q:[0,1] \rightarrow$ $\mathbf{C} \times \mathbf{C}^{n}$ such that $q(0)=\left(z_{0}, u^{0}\right)$ and $q(t) \in S$ for all $t \in(0,1]$. The result now follows by taking $p$ to be the first component of $q$.

In most cases the boundary of the pseudospectrum is straightforward to analyze without recourse to the above result. We make the following definition.

Definition 4.5. A point $z \in \mathbf{C}$ is degenerate if the smallest singular value of $A-z I$ is nonzero and simple (that is, has multiplicity one) and the corresponding right singular vector $u$ satisfies $u^{*}(A-z I) u=0$.

We need the following elementary identity.
Lemma 4.6. Given any unit vector $u \in \mathbf{C}^{n}$, matrix $B \in \mathbf{M}^{n}$, and scalar $w \in \mathbf{C}$, we have

$$
\|(B+w I) u\|^{2}-\|B u\|^{2}=\left|u^{*}(B+w I) u\right|^{2}-\left|u^{*} B u\right|^{2}
$$

The next result shows that, except possibly at degenerate points, the pseudospectrum can never be "pointed" outwards.

Proposition 4.7 (pointedness). Any nondegenerate point in the pseudospectrum lies on the boundary of an open disk contained in the strict pseudospectrum.

Proof. Consider a nondegenerate point $z_{0} \in \Lambda_{\epsilon}$. We may as well assume $g\left(z_{0}\right)=\epsilon$. Choose a unit right singular vector $u \in \mathbf{C}^{n}$ satisfying the condition $u^{*}\left(A-z_{0} I\right) u \neq 0$. We now claim

$$
\left|z-u^{*} A u\right|<\left|z_{0}-u^{*} A u\right| \quad \Rightarrow \quad z \in \Lambda_{\epsilon}^{\prime}
$$

To see this, observe that if $z$ satisfies the left-hand side, then

$$
\begin{aligned}
\sigma_{\min }^{2}(A-z I)-\epsilon^{2} & \leq\|(A-z I) u\|^{2}-\left\|\left(A-z_{0} I\right) u\right\|^{2} \\
& =\left|u^{*}(A-z I) u\right|^{2}-\left|u^{*}\left(A-z_{0} I\right) u\right|^{2} \\
& =\left|z-u^{*} A u\right|^{2}-\left|z_{0}-u^{*} A u\right|^{2} \\
& <0,
\end{aligned}
$$

using the preceding lemma.
In particular, this result shows that Theorem 4.4 is elementary in the case when the point of interest $z_{0}$ is nondegenerate.

Thus the pseudospectrum is not pointed outward, except possibly at a degenerate point. In fact, a more detailed analysis due to Trefethen shows that the pseudospectrum is never pointed outward [41]. However, it can certainly be pointed inward, as the following example shows.

Example 1 (nonsmooth points). Consider the matrix

$$
A=\left[\begin{array}{cc}
0 & 1 \\
-1 & 0
\end{array}\right]
$$

The pseudospectrum $\Lambda_{\epsilon}$ consists of the union of two disks of radius $\epsilon$, centered at the two eigenvalues, $\pm i$. For example, if $\epsilon=\sqrt{2}$, then $z=1$ is a nonsmooth point where the boundary of the pseudospectrum is pointed inward. In the case where $\epsilon=1$, the pseudospectrum consists of two disks tangent to each other at the origin.

Nonetheless, even though the boundary of the pseudospectrum can be nonsmooth, this cannot occur at any nondegenerate optimal solution of the pseudospectral abscissa problem.

Proposition 4.8 (optimal solutions). Any locally optimal solution $z_{0}$ of the pseudospectral abscissa problem (2.1) must lie on the boundary of the pseudospectrum. Furthermore, unless $z_{0}$ is degenerate, the boundary is differentiable there.

Proof. The fact that $z_{0}$ cannot lie in the interior of $\Lambda_{\epsilon}$ is immediate. Now assume $z_{0}$ is nondegenerate. Since $z_{0}$ is optimal, $\Lambda_{\epsilon}$ lies on or to the left of the vertical line through $z_{0}$. But since $z_{0}$ is nondegenerate, $\Lambda_{\epsilon}$ contains a closed disk whose boundary contains $z_{0}$, by Proposition 4.7 (pointedness). Thus the boundary of $\Lambda_{\epsilon}$ lies between the disk and the vertical line, which are tangent at $z_{0}$. This completes the proof.

Again, the nondegeneracy hypothesis may be dropped using the more general result on pointedness mentioned above [41].
5. Components of the pseudospectrum. We recall some basic ideas from plane topology. A domain is a nonempty, open, arcwise connected subset of C. Given a point $z$ in an open set $\Omega \subset \mathbf{C}$, a particular example of a domain is the component of $z$, which consists of all points that can be joined to $z$ by a continuous path in $\Omega$ [35].

The following result is in essence well known (see, for example, [10]).
ThEOREM 5.1 (eigenvalues and components). Every component of the strict pseudospectrum of the matrix $A$ contains an eigenvalue of $A$.

Proof. Suppose the set $S$ is a component of the strict pseudospectrum $\Lambda_{\epsilon}^{\prime}$ that contains no eigenvalues of $A$. The function $g$ attains its minimum on the compact set $\operatorname{cl} S$ at some point $z$, and clearly $g(z)<\epsilon$, so $z \in \Lambda_{\epsilon}^{\prime}$. Since $S$ is open and contains no eigenvalues, Theorem 4.2 (local minima) implies $z \notin S$.

But since $\Lambda_{\epsilon}^{\prime}$ is open, it contains an open disk $D$ centered at $z$. Since $z \in \operatorname{cl} S$, we know $D \cap S \neq \emptyset$, and hence $D \cup S$ is an arcwise connected subset of $\Lambda_{\epsilon}^{\prime}$ strictly larger than $S$. But this contradicts the definition of $S$.

In Example 1 (nonsmooth points), when $\epsilon=1$ the strict pseudospectrum consists of two components, namely the two open disks centered at the two eigenvalues, $\pm i$. By contrast, the pseudospectrum is arcwise connected.

The simplest case of the above result occurs when each eigenvalue has geometric multiplicity one and $\epsilon$ is small. In this case we show later (Corollary 7.5) that the pseudospectrum consists of disjoint compact convex neighborhoods of each eigenvalue (cf. [32]).

Our next aim is to try to bracket the pseudospectral abscissa. We first need a subsidiary result.

Lemma 5.2 (moving to the boundary). For any point $z_{1}$ in $\Lambda_{\epsilon}$ there exists a point $z_{2}$ satisfying $\operatorname{Re} z_{1}=\operatorname{Re} z_{2}$ and $g\left(z_{2}\right)=\epsilon$.

Proof. We simply take $z_{2}$ on the boundary of the intersection of the vertical line through $z_{1}$ and the pseudospectrum $\Lambda_{\epsilon}$ (which is compact).

Byers' algorithm for calculating the distance to instability [16] and its subsequent variants (see the introduction) all depend on versions of the following easy piece of linear algebra, relating singular values to imaginary eigenvalues of a certain Hamiltonian matrix. We include a proof for completeness.

Lemma 5.3 (imaginary eigenvalues). For real numbers $x$ and $y$, and $\epsilon \geq 0$, the matrix $A-(x+i y) I$ has a singular value $\epsilon$ if and only if the matrix

$$
\left[\begin{array}{cc}
x I-A^{*} & \epsilon I \\
-\epsilon I & A-x I
\end{array}\right]
$$

has an eigenvalue iy.
Proof. Plus and minus the singular values of any matrix $B \in \mathbf{M}^{n}$ are exactly the eigenvalues of the matrix

$$
\left[\begin{array}{cc}
0 & B \\
B^{*} & 0
\end{array}\right]
$$

Thus the matrix $A-(x+i y) I$ has a singular value $\epsilon$ if and only if $\epsilon$ is an eigenvalue of the matrix

$$
\left[\begin{array}{cc}
0 & A-(x+i y) I \\
A^{*}-(x-i y) I & 0
\end{array}\right]
$$

or, in other words, if and only if the matrix

$$
\left[\begin{array}{cc}
-\epsilon I & A-(x+i y) I \\
A^{*}-(x-i y) I & -\epsilon I
\end{array}\right]
$$

is singular. Since

$$
\begin{gathered}
{\left[\begin{array}{cc}
0 & I \\
-I & 0
\end{array}\right]\left[\begin{array}{cc}
-\epsilon I & A-(x+i y) I \\
A^{*}-(x-i y) I & -\epsilon I
\end{array}\right]} \\
=\left[\begin{array}{cc}
\left(A^{*}-x I\right) & -\epsilon I \\
\epsilon I & (x I-A)
\end{array}\right]+i y\left[\begin{array}{ll}
I & 0 \\
0 & I
\end{array}\right]
\end{gathered}
$$

this is equivalent to $i y$ being an eigenvalue of the matrix

$$
\left[\begin{array}{cc}
x I-A^{*} & \epsilon I \\
-\epsilon I & A-x I
\end{array}\right] .
$$

The following result is our key test. Geometrically it states simply that a given real $x$ (bigger than the spectral abscissa $\alpha$ ) is less than the pseudospectral abscissa exactly when the vertical line through $x$ intersects the boundary of the pseudospectrum. As we shall see, this is a straightforward computational test.

THEOREM 5.4 (bracketing the pseudospectral abscissa). For any real $x \geq \alpha$, the following statements are equivalent:
(i) $x \leq \alpha_{\epsilon}$;
(ii) the equation

$$
\begin{equation*}
g(x+i y)=\epsilon, \quad y \in \mathbf{R} \tag{5.1}
\end{equation*}
$$

is solvable;
(iii) the system

$$
i y \in \Lambda\left[\begin{array}{cc}
x I-A^{*} & \epsilon I  \tag{5.2}\\
-\epsilon I & A-x I
\end{array}\right], y \in \mathbf{R}
$$

is solvable.
Proof. We first show (i) $\Rightarrow$ (ii). If $x=\alpha_{\epsilon}$, then choose any point $z$ solving the pseudospectral abscissa problem (2.1). Clearly $z=x+i y$ for some real $y$, and $g(z)=\epsilon$, so we have shown that (5.1) has a solution.

We can therefore assume $x<\alpha_{\epsilon}$, in which case there exists a point $z_{1}$ such that $\operatorname{Re} z_{1}>x$ and $g\left(z_{1}\right)<\epsilon$. The component of $z_{1}$ in the strict pseudospectral abscissa $\Lambda_{\epsilon}^{\prime}$ contains an eigenvalue $z_{2}$ by Theorem 5.1 (eigenvalues and components). Hence there is an arc in this component connecting $z_{1}$ and $z_{2}$. But since $\operatorname{Re} z_{1}>x \geq \operatorname{Re} z_{2}$, this arc must contain a point $z_{3}$ with $\operatorname{Re} z_{3}=x$. Now applying Lemma 5.2 (moving to the boundary) gives a solution to (5.1).

The implication (ii) $\Rightarrow$ (iii) is immediate from Lemma 5.3 (imaginary eigenvalues), so it remains to show (iii) $\Rightarrow$ (i). But this is again an easy consequence of Lemma 5.3: if system (5.2) holds, then $\epsilon$ is a singular value of the matrix $A-(x+i y) I$, and hence the smallest singular value of this matrix is no greater than $\epsilon$, whence we get the result.

Using this result, the relationship (3.3) between the pseudospectral abscissa and the $\mathbf{H}_{\infty}$ norm is an easy exercise.

We can now approximate the pseudospectral abscissa $\alpha_{\epsilon}$ by a bisection search as follows.

Algorithm 5.5 (bisection method). We begin with the initial interval

$$
[\alpha,\|A\|+\epsilon] .
$$

We know $\alpha_{\epsilon}$ lies in this interval by the argument of Proposition 4.1 (compactness). Now at each iteration we let $x$ be the midpoint of the current interval and compute all the eigenvalues of the matrix

$$
\left[\begin{array}{cc}
x I-A^{*} & \epsilon I  \tag{5.3}\\
-\epsilon I & A-x I
\end{array}\right]
$$

If any of the eigenvalues are purely imaginary, then we deduce $x \leq \alpha_{\epsilon}$ and replace the current interval with its right half. Otherwise, by Theorem 5.4 (bracketing the pseudospectral abscissa), we know $x>\alpha_{\epsilon}$, so we replace the current interval with its left half. The intervals generated by this algorithm are guaranteed to converge to the pseudospectral abscissa $\alpha_{\epsilon}$.

The difference between this algorithm and Byers' bisection method for the distance to instability [16] is that the former searches for $x$ by bisection, while the latter searches for $\epsilon$ by bisection.

Notice that at each iteration of the bisection method we can easily solve (5.1). We first list the purely imaginary eigenvalues of the matrix (5.5), namely $\left\{i y_{1}, i y_{2}, \ldots, i y_{k}\right\}$. We then form the index set

$$
J=\left\{j: \sigma_{\min }\left(A-\left(x+i y_{j}\right) I\right)=\epsilon\right\}
$$

The set of solutions of (5.1) is then simply $\left\{y_{j}: j \in J\right\}$. As we shall see in the next section, the points $x+i y_{j}$ (for $j \in J$ ) provide good approximations to all the solutions of the pseudospectral abscissa problem (2.1).

A more sophisticated, quadratically convergent algorithm for the pseudospectral abscissa, based on similar ideas and analogous to $\mathbf{H}_{\infty}$ norm algorithms such as [7, 11, 24], is developed in [13].
6. Approximate solutions. The results generated by the bisection algorithm (or the algorithm in [13]) approximate all the global maximizers in the pseudospectral abscissa problem (2.1). To make this precise we use the following standard notion of set convergence [34]. We say that a sequence of sets $Y^{1}, Y^{2}, \ldots \subset \mathbf{R}$ converges to a set $Y \subset \mathbf{R}$ if the following properties hold:
(i) For any number $y \in Y$ there exists a sequence of numbers $y^{r} \in Y^{r}$ converging to $y$;
(ii) any cluster point of a sequence of numbers $y^{r} \in Y^{r}$ lies in $Y$.
(This notion is weaker than the idea of convergence with respect to the PompeiuHausdorff distance [34, Ex 4.13], although it is equivalent in the case when the sets $Y^{r}$ and $Y$ are uniformly bounded, as will be the case in our application below.)

We now prove a rather general result.
ThEOREM 6.1 (global maximizers). The number of global maximizers of the pseudospectral abscissa problem (2.1) does not exceed n. Denote these

$$
\left\{\alpha_{\epsilon}+i y: y \in Y\right\}
$$

where $Y \subset \mathbf{R}$. Consider any real sequence $\alpha \leq x^{r} \uparrow \alpha_{\epsilon}$. Then the sets

$$
Y_{r}=\left\{y \in \mathbf{R}: g\left(x^{r}+i y\right)=\epsilon\right\}
$$

converge to $Y$.
Proof. The pseudospectral abscissa problem (2.1) has at least one maximizer, by compactness. Furthermore, any solution $z=\alpha_{\epsilon}+i y$ must satisfy the equation $g(z)=\epsilon$. Just as in the proof of Theorem 5.4 (bracketing the pseudospectral abscissa), this implies that $y$ must satisfy the equation

$$
\operatorname{det}\left[\begin{array}{cc}
-\epsilon I & A-\left(\alpha_{\epsilon}+i y\right) I \\
A^{*}-\left(\alpha_{\epsilon}-i y\right) I & -\epsilon I
\end{array}\right]=0
$$

But this polynomial equation has at most $2 n$ solutions, so we can write

$$
Y=\left\{y_{1}, y_{2}, \ldots, y_{m}\right\}
$$

where $1 \leq m \leq 2 n$.
Fix the index $j \in\{1,2, \ldots, m\}$. Theorem 4.4 (accessibility) and Theorem 5.1 (eigenvalues and components) together imply the existence of a continuous function $p:[0,1] \rightarrow \mathbf{C}$ such that $p(0)=\alpha_{\epsilon}+i y_{j}, p(1)$ is an eigenvalue, and $p(t) \in \Lambda_{\epsilon}^{\prime}$ for all $t>0$.

Using the continuity of $p$, we can now iteratively construct a nonincreasing sequence $\left\{t_{r}\right\} \subset[0,1]$ such that $\operatorname{Re} p\left(t_{r}\right)=x^{r}$ for all $r$. Taking limits shows

$$
\operatorname{Re} p\left(\lim _{r} t_{r}\right)=\lim _{r} x^{r}=\alpha_{\epsilon} .
$$

But for $t>0$ we have $g(p(t))<\epsilon$, which implies $\operatorname{Re} p(t)<\alpha_{\epsilon}$, so we deduce $t_{r} \downarrow 0$. Hence if we define $v_{j}^{r}=\operatorname{Im} p\left(t_{r}\right)$, we have $v_{j}^{r} \rightarrow y_{j}$.

For each index $r$, consider the bounded open set

$$
\left\{y \in \mathbf{R}: g\left(x^{r}+i y\right)<\epsilon\right\}
$$

If $t_{r}=0$, this set is empty, and we define $l_{j}^{r}=u_{j}^{r}=y_{j}$. Otherwise, denote the component of $v_{j}^{r}$ in this set by the open interval $\left(l_{j}^{r}, u_{j}^{r}\right)$. By continuity, $l_{j}^{r}$ and $u_{j}^{r}$ are both zeros of the function

$$
\begin{equation*}
y \in \mathbf{R} \mapsto g\left(x^{r}+i y\right)-\epsilon \tag{6.1}
\end{equation*}
$$

We now claim both $l_{j}^{r} \rightarrow y_{j}$ and $u_{j}^{r} \rightarrow y_{j}$.
If this claim fails, then without loss of generality, after taking a subsequence, we can assume $l_{j}^{r} \rightarrow w<y_{j}$. By definition, we know

$$
g\left(x^{r}+i\left(s v_{j}^{r}+(1-s) l_{j}^{r}\right)\right)<\epsilon \text { for all } s \in(0,1], \quad r=1,2, \ldots
$$

so taking limits shows

$$
g\left(\alpha_{\epsilon}+i\left(s y_{j}+(1-s) w\right)\right) \leq \epsilon \text { for all } s \in[0,1]
$$

But in this case every point in the line segment $\alpha_{\epsilon}+i\left[w, y_{j}\right]$ solves the pseudospectral abscissa problem (2.1), contradicting the fact that there are only finitely many solutions. This proves the claim. We have thus shown property (i) in the definition of set convergence: the constructed sequence $\left(l_{j}^{r}\right)$ converges to the desired point $y_{j}$. Property (ii) is immediate.

Finally, suppose $m>n$. Choose any nondecreasing sequence $\left\{x^{r}\right\} \subset\left[\alpha, \alpha_{\epsilon}\right)$ converging to $\alpha_{\epsilon}$, and for each index $r$ construct the set

$$
\left\{l_{j}^{r}, u_{j}^{r}: j=1,2, \ldots, m\right\}
$$

as above. Then for $r$ sufficiently large, this is a set of $2 m$ distinct zeros of the function (6.1), and hence of the polynomial

$$
\operatorname{det}\left[\begin{array}{cc}
-\epsilon I & A-\left(x^{r}+i y\right) I \\
A^{*}-\left(x^{r}-i y\right) I & -\epsilon I
\end{array}\right]
$$

But this polynomial is not identically zero, and has degree $2 n$, which is a contradiction.

The algorithmic significance of the above result is this: Consider any algorithm that generates a sequence of lower approximations to the pseudospectral abscissa, $x^{r} \uparrow \alpha_{\epsilon}$. In particular, we could consider the bisection algorithm of the previous section. For each step $r$, an eigenvalue computation generates the set $Y^{r} \subset \mathbf{R}$, as described after Algorithm 5.5. The above result now shows that this set is a good approximation to $Y$, and hence gives us a good approximation to the set of all optimal solutions to the pseudospectral abscissa problem.
7. Smoothness. To study the smoothness of the function $g$, and hence the boundary of the pseudospectrum, we rely on the following well-known result. We consider $\mathbf{M}^{n}$ as a Euclidean space with inner product

$$
\langle X, Y\rangle=\operatorname{Retr}\left(X^{*} Y\right) \quad\left(X, Y \in \mathbf{M}^{n}\right)
$$

A real-valued function on a real vector space is real-analytic at zero if in some neighborhood of zero is can be written as the sum of an absolutely convergent power series
in the coordinates relative to some basis, and we make an analogous definition at other points. In particular, such functions are smooth $\left(C^{\infty}\right)$ near the point in question.

We call vectors $u, v \in \mathbf{C}^{n}$ minimal left and right singular vectors for a matrix $Z \in \mathbf{M}^{n}$ if

$$
Z v=\sigma_{\min }(Z) u \text { and } Z^{*} u=\sigma_{\min }(Z) v
$$

Theorem 7.1 (analytic singular value). If the matrix $Z$ has a simple smallest singular value, then the function $\sigma_{\min }^{2}$ is real-analytic at $Z$. If, furthermore, $\sigma_{\min }(Z)>$ 0 , then $\sigma_{\min }$ is real-analytic at $Z$, with gradient

$$
\nabla \sigma_{\min }(Z)=u v^{*}
$$

for any unit minimal left and right singular vectors $u, v \in \mathbf{C}^{n}$.
Proof. The matrix

$$
\left(X^{T}-i Y^{T}\right)(X+i Y)
$$

depends analytically on the matrices $X, Y \in \mathbf{M}^{n}$ and has a simple eigenvalue $\sigma_{\text {min }}^{2}(Z)$ when $(X, Y)=\left(X_{0}, Y_{0}\right)$ for real matrices $X_{0}, Y_{0} \in \mathbf{M}^{n}$ satisfying $Z=X_{0}+i Y_{0}$. Hence by standard perturbation theory [23], the above matrix has a unique eigenvalue near $\sigma_{\min }^{2}(Z)$ for all $(X, Y)$ close to $\left(X_{0}, Y_{0}\right)$, depending analytically on $(X, Y)$. When $X$ and $Y$ are real, this eigenvalue is exactly $\sigma_{\min }^{2}(X+i Y)$, so the first part follows. The second part follows by taking square roots. The gradient calculation is standard (see, for example, [37]).

We next turn to smoothness properties of the function $g: \mathbf{C} \rightarrow \mathbf{R}$ defined by

$$
g(z)=\sigma_{\min }(A-z I)
$$

We will often find it more convenient to work with the squared function $g^{2}(z)=$ $(g(z))^{2}$.

We can treat $\mathbf{C}$ as a Euclidean space, where we define the inner product by $\langle w, z\rangle=\operatorname{Re}\left(w^{*} z\right)$.

Corollary 7.2 (analytic boundary). If the singular value $\sigma_{\min }\left(A-z_{0} I\right)$ is simple, then the function $g^{2}$ is real-analytic at $z_{0}$. If, furthermore, this singular value is strictly positive, then $g$ is real-analytic at $z_{0}$, with gradient

$$
\nabla g\left(z_{0}\right)=-v^{*} u
$$

where the vectors $u, v \in \mathbf{C}^{n}$ are unit minimal left and right singular vectors for $A-$ $z_{0} I$.

Proof. This follows from the previous result by the chain rule.
Thus what we called "degenerate" points are simply smooth critical points of $g$, distinct from the eigenvalues. At a nondegenerate smooth point $z_{0}$ with $g\left(z_{0}\right)=\epsilon$, the gradient of $g$ is nonzero, and hence the boundary of the pseudospectrum

$$
\Lambda_{\epsilon}=\{z \in \mathbf{C}: g(z) \leq \epsilon\}
$$

is simply a smooth curve locally, with normal $u^{*}\left(A-z_{0}\right) u$ at $z_{0}$.
We call an eigenvalue of $A$ nonderogatory if it has geometric multiplicity one. This is the most common type of multiple eigenvalue (from the perspective of the dimensions of the corresponding manifolds in $\mathbf{M}^{n}[1]$ ). The following result is very well known.

Proposition 7.3 (nonderogatory eigenvalues). The point $\lambda_{0} \in \mathbf{C}$ is a nonderogatory eigenvalue of the matrix $A$ if and only if 0 is a simple singular value of $A-\lambda_{0} I$.

Proof. First, note that $\lambda_{0}$ is an eigenvalue of $A$ if and only if $A-\lambda_{0} I$ is singular, which is equivalent to 0 being a singular value of $A-\lambda_{0} I$. Second, $v$ is a corresponding eigenvector of $A$ exactly when $\left(A-\lambda_{0} I\right) v=0$, which says that $v$ is a right singular vector of $A-\lambda_{0} I$ corresponding to the singular value 0 . Thus the eigenspace of $A$ corresponding to the eigenvalue $\lambda_{0}$ coincides with the subspace of right singular vectors of $A-\lambda_{0} I$ corresponding to the singular value 0 , so in particular these spaces have the same dimension. The result now follows.

We can now show that the function $g$ is well behaved near any nonderogatory eigenvalue of $A$.

THEOREM 7.4 (growth near an eigenvalue). Let $\lambda_{0}$ be a nonderogatory eigenvalue of multiplicity $m$ for the matrix $A$. Then

$$
\sigma_{\min }(A-z I)=g(z)=\frac{\prod_{j=1}^{n-m}\left|\lambda_{j}-\lambda_{0}\right|}{\prod_{k=1}^{n-1} \sigma_{k}}\left|z-\lambda_{0}\right|^{m}+O\left(\left|z-\lambda_{0}\right|^{m+1}\right)
$$

for complex $z$ near $\lambda_{0}$, where $\lambda_{1}, \lambda_{2}, \ldots, \lambda_{n-m}$ are the eigenvalues of $A$ distinct from $\lambda_{0}$ (listed by multiplicity) and $\sigma_{1}, \sigma_{2}, \ldots, \sigma_{n-1}$ are the nonzero singular values of $A-\lambda_{0} I$ (listed by multiplicity). (In the case $n=1$ or $m=n$, we interpret the empty products appearing in the above expression as 1.)

Furthermore, the function $g^{2}$ has positive definite Hessian at all points $z \neq \lambda_{0}$ near $\lambda_{0}$.

Proof. We prove the case $\lambda_{0}=0$ : the general case follows by a simple transformation.

Since 0 is a nonderogatory eigenvalue of $A$, Proposition 7.3 (nonderogatory eigenvalues) shows 0 is a simple singular value of $A$. Hence by Corollary 7.2 (analytic boundary), the function $f: \mathbf{R}^{2} \rightarrow \mathbf{R}$ defined by

$$
f(x, y)=(g(x+i y))^{2}
$$

is real-analytic at $(0,0)$.
Consider any point $(x, y) \in \mathbf{R}^{2}$, and let $z=x+i y$. The matrix

$$
(A-z I)^{*}(A-z I)
$$

is Hermitian, so its characteristic polynomial

$$
p_{z}(\mu)=\operatorname{det}\left((A-z I)^{*}(A-z I)-\mu I\right)
$$

has all real coefficients. Hence we can write

$$
p_{z}(\mu)=\sum_{r=0}^{n} q_{r}(x, y) \mu^{r}
$$

for some real polynomials $q_{r}$. The smallest zero of $p_{z}$ is $f(x, y)$.
We concentrate on the two lowest-order coefficients of the above polynomial. First, note

$$
q_{0}(x, y)=p_{z}(0)
$$

$$
\begin{aligned}
& =\operatorname{det}\left((A-z I)^{*}(A-z I)\right) \\
& =|\operatorname{det}(A-z I)|^{2} \\
& =|z|^{2 m} \prod_{j=1}^{n-m}\left|\lambda_{j}-z\right|^{2}
\end{aligned}
$$

Hence for small $(x, y)$ we have

$$
\begin{equation*}
q_{0}(x, y)=\left(\left(x^{2}+y^{2}\right)^{m} \prod_{j=1}^{n-m}\left|\lambda_{j}\right|^{2}\right)+O\left(\|(x, y)\|^{2 m+1}\right) \tag{7.1}
\end{equation*}
$$

Turning to the coefficient of $\mu$, notice

$$
p_{0}(\mu)=\operatorname{det}\left(A^{*} A-\mu I\right)=-\mu \prod_{k=1}^{n-1}\left(\sigma_{k}^{2}-\mu\right)
$$

so

$$
q_{1}(0,0)=-\prod_{k=1}^{n-1} \sigma_{k}^{2}
$$

(notice this is nonzero), and hence

$$
\begin{equation*}
q_{1}(x, y)=-\prod_{k=1}^{n-1} \sigma_{k}^{2}+O(\|x, y\|) \tag{7.2}
\end{equation*}
$$

Since the function $f$ is real-analytic at $(0,0)$, we know for some integer $t=1,2, \ldots$,

$$
f(x, y)=s(x, y)+O\left(\|(x, y)\|^{t+1}\right)
$$

for some nonzero homogeneous polynomial $s$ of degree $t$. Now substituting into the relationship

$$
\sum_{r=0}^{n} q_{r}(x, y)(f(x, y))^{r}=0
$$

and using (7.1) and (7.2) shows $t=2 m$, and

$$
f(x, y)=\frac{\prod_{j=1}^{n-m}\left|\lambda_{j}\right|^{2}}{\prod_{k=1}^{n-1} \sigma_{k}^{2}}\left(x^{2}+y^{2}\right)^{m}+O\left(\|(x, y)\|^{2 m+1}\right)
$$

as required.
It remains to show that the Hessian $\nabla^{2} f(x, y)$ is positive definite for all small $(x, y) \neq(0,0)$. We have shown that $f$ is analytic at $(0,0)$ and

$$
f(x, y)=\tau\left(x^{2}+y^{2}\right)^{m}+O\left(\|(x, y)\|^{2 m+1}\right)
$$

for some nonzero constant $\tau$. Since we can differentiate the power series for $f$ term-by-term, a short calculation shows

$$
\nabla^{2} f(x, y)=2 m \tau\left(x^{2}+y^{2}\right)^{m-2} H(x, y)
$$

where

$$
H(x, y)=\left[\begin{array}{cc}
(2 m-1) x^{2}+y^{2} & 2(m-1) x y \\
2(m-1) x y & x^{2}+(2 m-1) y^{2}
\end{array}\right]+O\left(\|(x, y)\|^{3}\right)
$$

We deduce

$$
H_{11}(x, y)=(2 m-1) x^{2}+y^{2}+O\left(\|(x, y)\|^{3}\right)>0
$$

and furthermore

$$
\operatorname{det} H(x, y)=(2 m-1)\left(x^{2}+y^{2}\right)^{2}+O\left(\|(x, y)\|^{5}\right)>0
$$

so the matrix $H(x, y)$ is positive definite for all small $(x, y) \neq(0,0)$. The result now follows.

The following results are immediate consequences.
Corollary 7.5 (convexity). If $\lambda_{0}$ is a nonderogatory eigenvalue of the matrix A, then for all small $\epsilon>0$ the pseudospectrum $\Lambda_{\epsilon}$ near $\lambda_{0}$ consists of a compact, strictly convex neighborhood of $\lambda_{0}$.

Proof. We can consider the pseudospectrum as a level set of the function $g^{2}$, which is strictly convex near $\lambda_{0}$.

Corollary 7.6 (smoothness). If $\lambda_{0}$ is a nonderogatory eigenvalue of the matrix A, then the function $g$ is smooth with nonzero gradient at all nearby points distinct from $\lambda_{0}$.

Proof. Since the real-analytic function $g^{2}$ is strictly convex near the eigenvalue $\lambda_{0}$, with a strict local minimizer there, it follows that $\lambda_{0}$ is an isolated critical point of $g^{2}$. It is then easy to see that $g$ is smooth and noncritical near $\lambda_{0}$.

If the matrix $A$ has an eigenvalue $\lambda_{0}$ of multiplicity $m$, then, by continuity of the set of eigenvalues, any matrix close to $A$ will have exactly $m$ eigenvalues close to $\lambda_{0}$ (counted by multiplicity). Our last corollary bounds how far these eigenvalues can be from $\lambda_{0}$.

Corollary 7.7 (Hölder continuity). With the assumptions and notation of Theorem 7.4, consider any constant

$$
\kappa>\left(\frac{\prod_{k=1}^{n-1} \sigma_{k}}{\prod_{j=1}^{n-m}\left|\lambda_{j}-\lambda_{0}\right|}\right)^{1 / m}
$$

For any matrix $Z$ close to $A$, any eigenvalue $z$ of $Z$ close to $\lambda_{0}$ satisfies

$$
\left|z-\lambda_{0}\right| \leq \kappa\|Z-A\|^{1 / m} .
$$

Proof. This follows easily from Theorem 7.4 (growth near an eigenvalue), using the elementary property that

$$
\sigma_{\min }(A-z I) \leq\|Z-A\|
$$

for any eigenvalue $z$ of $Z$.
In the above result, if we specialize to the case of a perturbation $Z=A+t B$ (where $t$ is a complex parameter), then the result shows that the eigenvalues of $A+$ $t B$ near a nonderogatory eigenvalue of $A$ of multiplicity $m$ satisfy an $m^{-1}$-Hölder continuity condition in $t$. This is a well-known result; see [27, 23, 3, 30].
8. Smoothness and regularity of the pseudospectral abscissa. Our ultimate goal is an understanding of how the pseudospectral abscissa $\alpha_{\epsilon}$ depends on the underlying matrix $A$. We therefore now allow $A$ (and $\epsilon$ ) to vary. Recall that the pseudospectral abscissa function $\alpha_{\epsilon}: \mathbf{M}^{n} \rightarrow \mathbf{R}$ is given by

$$
\alpha_{\epsilon}(Z)=\max \left\{\operatorname{Re} z: \sigma_{\min }(Z-z I) \leq \epsilon\right\},
$$

and for any nonempty set $\Omega \subset \mathbf{C}$ we define the refinement

$$
\begin{equation*}
\alpha^{\Omega}(Z, \epsilon)=\sup \left\{\operatorname{Re} z: z \in \Omega, \sigma_{\min }(Z-z I) \leq \epsilon\right\} \tag{8.1}
\end{equation*}
$$

Thus for $\Omega=\mathbf{C}$, we obtain exactly the pseudospectral abscissa. We now apply classical sensitivity analysis to differentiate this function.

THEOREM 8.1 (smoothness of pseudospectral abscissa). Suppose that, for $\epsilon=$ $\epsilon_{0}>0$ and $Z=A$, the supremum (8.1) is attained by a point $z_{0} \in \operatorname{int} \Omega$, where the singular value $\sigma_{\min }\left(A-z_{0} I\right)$ is simple. Then for any corresponding unit minimal left and right singular vectors $u, v \in \mathbf{C}^{n}$, the number $v^{*} u$ is real and nonpositive.

Now suppose furthermore that $z_{0}$ is the unique attaining point in (8.1), that it is nondegenerate (or, in other words, $v^{*} u \neq 0$ ), and that the Hessian $\nabla^{2}\left(g^{2}\right)\left(z_{0}\right)$ is nonsingular. Then the function $\alpha^{\Omega}$ is smooth around the point $\left(A, \epsilon_{0}\right)$, with

$$
\nabla_{Z} \alpha^{\Omega}\left(A, \epsilon_{0}\right)=\frac{u v^{*}}{v^{*} u} \text { and } \nabla_{\epsilon} \alpha^{\Omega}\left(A, \epsilon_{0}\right)=-\frac{1}{v^{*} u} .
$$

Proof. Consider the optimization problem

$$
\left\{\right.
$$

When $(Z, \epsilon)=\left(A, \epsilon_{0}\right)$ this problem becomes

$$
\left\{\begin{array}{lr}
\sup & \operatorname{Re} z \\
& \\
\text { subject to } & g^{2}(z)
\end{array} \leq \epsilon_{0}^{2},\right.
$$

with optimal solution $z_{0}$. By Corollary 7.2 (analytic boundary), the function $g^{2}$ is smooth near $z_{0}$, with gradient

$$
\nabla g^{2}\left(z_{0}\right)=2 g\left(z_{0}\right) \nabla g\left(z_{0}\right)=-2 \epsilon_{0} v^{*} u
$$

Either this gradient is zero or there is a Lagrange multiplier $\mu \in \mathbf{R}_{+}$such that the gradient of the Lagrangian

$$
z \mapsto \operatorname{Re} z-\mu\left(g^{2}(z)-\epsilon_{0}^{2}\right)
$$

at $z=z_{0}$ is zero. In this case,

$$
\begin{equation*}
1+2 \mu \epsilon_{0} v^{*} u=0 \tag{8.2}
\end{equation*}
$$

so the first part follows.
Moving to the second part, (8.2) implies $\mu=-\left(2 \epsilon_{0} v^{*} u\right)^{-1}$. Under the additional assumptions we can apply a standard sensitivity result (for example, [6, Thm 5.5.3])
to deduce that the gradient of the optimal value of the original optimization problem at $\left(A, \epsilon_{0}\right)$ equals the gradient of the Lagrangian

$$
(Z, \epsilon) \mapsto \operatorname{Re} z_{0}+\left(2 \epsilon_{0} v^{*} u\right)^{-1}\left(\sigma_{\min }^{2}\left(Z-z_{0} I\right)-\epsilon^{2}\right)
$$

at $\left(A, \epsilon_{0}\right)$. The result now follows by Theorem 7.1 (analytic singular value).
An eigenvalue of the matrix $A$ with real part equal to the spectral abscissa $\alpha$ is called active.

Theorem 8.2 (regular representation). If the matrix $A$ has $s$ distinct active eigenvalues, all of which are nonderogatory, then there exist s functions

$$
\gamma_{j}: \mathbf{M}^{n} \times \mathbf{R}_{++} \rightarrow \mathbf{R} \quad(j=1,2, \ldots, s)
$$

such that for small $\epsilon>0$ and matrices $Z$ close to $A$, each map

$$
(Z, \epsilon) \mapsto \gamma_{j}(Z, \epsilon)
$$

is smooth and satisfies $\gamma_{j}(A, 0)=\alpha(A)$, the pseudospectral abscissa can be expressed as

$$
\alpha_{\epsilon}(Z)=\max \left\{\gamma_{j}(Z, \epsilon): j=1,2, \ldots, s\right\}
$$

and the set of gradients

$$
\left\{\nabla_{Z} \gamma_{j}(A, \epsilon): j=1,2, \ldots, s\right\}
$$

is linearly independent.
Proof. Denote the distinct eigenvalues of $A$ by $\lambda_{1}, \lambda_{2}, \ldots, \lambda_{k}$, where

$$
\operatorname{Re} \lambda_{j}\left\{\begin{array}{lll}
= & \alpha & (j \leq s) \\
< & \alpha & (j>s)
\end{array}\right.
$$

Let $D$ denote the open unit disk in $\mathbf{C}$. Providing we choose a radius $\delta>0$ sufficiently small, we have

$$
\begin{aligned}
2 \delta & <\left|\lambda_{p}-\lambda_{q}\right| \text { for all } p \neq q \\
\delta+\operatorname{Re} \lambda_{j} & <\alpha \text { for all } j>s
\end{aligned}
$$

and so the open disks $\lambda_{j}+\delta D$ are disjoint, and those with $j>m$ lie in the half-plane $\operatorname{Re} z<\alpha$. Furthermore, again by reducing $\delta$ if necessary, Theorem 7.4 (growth near eigenvalues) guarantees that each of the functions

$$
\left.g^{2}\right|_{\lambda_{j}+\delta D} \quad(j=1,2, \ldots, s)
$$

is smooth, with everywhere positive definite Hessian except possibly at $\lambda_{j}$.
We claim that the small pseudospectra (by which we mean pseudospectra corresponding to small $\epsilon$ ) of matrices close to $A$ lie in small disks around the eigenvalues of $A$. More precisely, for small $\epsilon \geq 0$ and matrices $Z$ close to $A$, we claim

$$
\begin{equation*}
\left\{z \in \mathbf{C}: \sigma_{\min }(Z-z I) \leq \epsilon\right\} \subset\left\{\lambda_{1}, \lambda_{2}, \ldots, \lambda_{k}\right\}+\delta D \tag{8.3}
\end{equation*}
$$

Otherwise there would exist sequences $\epsilon_{r} \rightarrow 0, Z_{r} \rightarrow A$, and $z_{r} \in \mathbf{C}$ satisfying, for all $r=1,2, \ldots$,

$$
\begin{aligned}
\sigma_{\min }\left(Z_{r}-z_{r} I\right) & \leq \epsilon_{r} \\
\left|z_{r}-\lambda_{j}\right| & \geq \delta \quad(j=1,2, \ldots, k)
\end{aligned}
$$

The first inequality above implies the sequence $\left\{z_{r}\right\}$ is bounded, so has a cluster point $z_{0}$, which must satisfy the inequalities

$$
\begin{aligned}
\sigma_{\min }\left(A-z_{0} I\right) & \leq 0 \\
\left|z_{0}-\lambda_{j}\right| & \geq \delta \quad(j=1,2, \ldots, k)
\end{aligned}
$$

The first inequality above can only hold if $z_{0}$ is an eigenvalue of $A$, which contradicts the second inequality. Hence inequality (8.3) holds, as we claimed.

Using the notation of (8.1), we can, for small $\epsilon>0$, matrices $Z$ close to $A$ and, for each $j=1,2, \ldots, m$, define functions

$$
\gamma_{j}(Z, \epsilon)=\alpha^{\lambda_{j}+\delta D}(Z, \epsilon)=\sup \left\{\operatorname{Re} z:\left|z-\lambda_{j}\right|<\delta, \sigma_{\min }(Z-z I) \leq \epsilon\right\}
$$

and as a consequence of inclusion (8.3), we can then write

$$
\alpha_{\epsilon}(Z)=\max \left\{\gamma_{j}(Z, \epsilon): j=1,2, \ldots, m\right\} .
$$

We claim each function $\gamma_{j}$ is smooth around the point $\left(A, \epsilon_{0}\right)$ for any small $\epsilon_{0}>0$.
To prove this claim, we use Theorem 8.1 (smoothness of pseudospectral abscissa). For any $j=1,2, \ldots, s$, consider the supremum

$$
\begin{aligned}
\gamma_{j}\left(A, \epsilon_{0}\right) & =\sup \left\{\operatorname{Re} z:\left|z-\lambda_{j}\right|<\delta, \sigma_{\min }(A-z I) \leq \epsilon_{0}\right\} \\
& =\sup \left\{\operatorname{Re} z:\left.g^{2}\right|_{\lambda_{j}+\delta D}(z) \leq \epsilon_{0}^{2}\right\}
\end{aligned}
$$

By our choice of the radius $\delta$, this supremum is attained at a unique point $z_{j}$ (cf. Corollary 7.5 (convexity)), which is nondegenerate (cf. Corollary 7.6 (smoothness)), and at which the Hessian $\nabla^{2}\left(g^{2}\right)\left(z_{j}\right)$ is positive definite. Hence the function $\gamma_{j}$ is smooth around $\left(A, \epsilon_{0}\right)$, with gradient

$$
\nabla_{Z} \gamma_{j}\left(A, \epsilon_{0}\right)=\frac{u_{j} v_{j}^{*}}{v_{j}^{*} u_{j}}
$$

where $u_{j}, v_{j}$ are unit minimal left and right singular vectors for $A-z_{j} I$, and $v_{j}^{*} u_{j}$ is real and strictly negative.

To complete the proof, it suffices to show that the set of matrices

$$
\left\{u_{j} v_{j}^{*}: j=1,2, \ldots, s\right\} \subset \mathbf{M}^{n}
$$

is linearly independent providing our choice of radius $\delta>0$ is sufficiently small. If this fails, then for each $j$ there is a sequence of points $z_{j}^{r} \rightarrow \lambda_{j}$ and sequences of unit minimal left and right singular vectors $u_{j}^{r}, v_{j}^{r}$ for $A-z_{j}^{r} I$ such that the set of matrices

$$
\left\{u_{j}^{r}\left(v_{j}^{r}\right)^{*}: j=1,2, \ldots, s\right\} \subset \mathbf{M}^{n}
$$

is linearly dependent. By taking subsequences, we can suppose $u_{j}^{r} \rightarrow u_{j}^{0}$ and $v_{j}^{r} \rightarrow v_{j}^{0}$ for each $j$, and then the set

$$
S=\left\{u_{j}^{0}\left(v_{j}^{0}\right)^{*}: j=1,2, \ldots, s\right\} \subset \mathbf{M}^{n}
$$

must be linearly dependent. But it also follows that $u_{j}^{0}, v_{j}^{0}$ are unit left and right eigenvectors for $A$ corresponding to the eigenvalue $\lambda_{j}$. Since the eigenvalues $\lambda_{1}, \lambda_{2}, \ldots, \lambda_{s}$ are distinct, the sets of eigenvectors

$$
\left\{u_{j}^{0}: j=1,2, \ldots, m\right\} \text { and }\left\{v_{j}^{0}: j=1,2, \ldots, s\right\}
$$

are each linearly independent, and a standard exercise then shows the contradiction that the set $S$ above is linearly independent.

As a consequence of this result, the pseudospectral abscissa must be a reasonably well-behaved nonsmooth function near a matrix with all nonderogatory eigenvalues. Specifically, we have the following result. We refer the reader to $[17,34]$ for standard nonsmooth terminology.

Corollary 8.3 (regularity). If all the active eigenvalues of a matrix $A \in \mathbf{M}^{n}$ are nonderogatory, then for all small $\epsilon>0$ the pseudospectral abscissa $\alpha_{\epsilon}$ is locally Lipschitz and subdifferentially regular around $A$.

Proof. This follows immediately from the representation as a maximum of smooth functions in the previous result [17, Prop. 2.3.12].

This corollary presents an interesting parallel with a key result in [15]. This result states that the spectral abscissa, even though non-Lipschitz, is a subdifferentially regular function around the matrix $A$ if and only if each active eigenvalue of $A$ is nonderogatory.

By combining the representation of $\alpha_{\epsilon}$ constructed in the proof of Theorem 8.2 with the growth estimate of Theorem 7.4 , we can also see how the pseudospectral abscissa depends on the parameter $\epsilon$.

Corollary 8.4 (dependence on $\epsilon$ ). If all the active eigenvalues of the matrix $A$ are nonderogatory, with maximum algebraic multiplicity $m$, then as a function of $\epsilon \geq 0$ we have

$$
\alpha_{\epsilon}-\alpha \sim \gamma \epsilon^{1 / m} \quad \text { as } \epsilon \downarrow 0
$$

for some constant $\gamma>0$.
9. Nonsmooth geometry. What about points $z_{0} \in \mathbf{C}$, where $\sigma_{\min }\left(A-z_{0} I\right)$ is multiple? The function $\sigma_{\min }$ is nonsmooth at any matrix with a multiple smallest singular value, so the function $g$ may be nonsmooth at $z_{0}$. An appropriate approach to studying the pseudospectrum near $z_{0}$ is therefore to use nonsmooth analysis. We again refer to $[17,34]$ for the standard concepts.

For any point $z \in \mathbf{C}$ we consider the subspace $U(z) \subset \mathbf{C}^{n}$ spanned by all right singular vectors corresponding to $\sigma_{\min }(A-z I)$, and we define a subset of $\mathbf{C}$ by

$$
G(z)=\left\{u^{*}(A-z I) u: u \in U(z),\|u\|=1\right\}
$$

Proposition 9.1 (convexity). The set $G(z)$ is nonempty, compact, and convex.
Proof. Define a linear map $B: U(z) \rightarrow U(z)$ by

$$
B u=P_{U(z)}((A-z I) u)
$$

where $P_{U(z)}: \mathbf{C}^{n} \rightarrow U(z)$ denotes the orthogonal projection. Now notice for all vectors $u \in U(z)$ we have

$$
\begin{array}{r}
\langle u, B u\rangle=\left\langle u, P_{U(z)}((A-z I) u)\right\rangle=\left\langle P_{U(z)}^{*} u,(A-z I) u\right\rangle \\
=\langle u,(A-z I) u\rangle=u^{*}(A-z I) u
\end{array}
$$

since the map $P_{U(z)}^{*}: U(z) \rightarrow \mathbf{C}^{n}$ is just the embedding. We deduce

$$
G(z)=\{\langle u, B u\rangle: u \in U(z),\|u\|=1\}
$$

and this set is nonempty, compact, and convex, by the Toeplitz-Hausdorff theorem [22].

The next result gives another perspective on the pointedness of the pseudospectrum (recall Proposition 4.7 (pointedness)).

THEOREM 9.2 (nonsmooth boundary behavior). For complex $z_{0}$ satisfying $g\left(z_{0}\right)=$ $\epsilon$ and $0 \notin G\left(z_{0}\right)$, the complement of the strict pseudospectrum,

$$
\{z \in \mathbf{C}: g(z) \geq \epsilon\}
$$

is Clarke regular at $z_{0}$, with normal cone cone $\left(G\left(z_{0}\right)\right)$.
Proof. The complement of the strict pseudospectrum is

$$
\begin{aligned}
& \left\{z \in \mathbf{C}: \sigma_{\min }(A-z I) \geq \epsilon\right\} \\
& =\left\{z: \lambda_{\min }\left((A-z I)^{*}(A-z I)\right) \geq \epsilon^{2}\right\} \\
& =\left\{z: F(z) \in \mathbf{H}_{+}^{n}\right\}=F^{-1}\left(\mathbf{H}_{+}^{n}\right),
\end{aligned}
$$

where $\mathbf{H}^{n}$ denotes the Euclidean space of $n \times n$ Hermitian matrices, with inner product $\langle X, Y\rangle=\operatorname{Re}(\operatorname{tr}(X Y))$ and positive semidefinite cone $\mathbf{H}_{+}^{n}$, the function $\lambda_{\min }: \mathbf{H}^{n} \rightarrow$ $\mathbf{R}$ is the smallest eigenvalue, and the function $F: \mathbf{C} \rightarrow \mathbf{H}^{n}$ is defined by

$$
F(z)=(A-z I)^{*}(A-z I)-\epsilon^{2} I
$$

The gradient map $\nabla F\left(z_{0}\right): \mathbf{C} \rightarrow \mathbf{H}^{n}$ is given by

$$
\nabla F\left(z_{0}\right)(w)=-w^{*} A-w A^{*}+2\left\langle w, z_{0}\right\rangle I
$$

and a short calculation shows that the adjoint map $\nabla F\left(z_{0}\right)^{*}: \mathbf{H}^{n} \rightarrow \mathbf{C}$ is given by

$$
\nabla F\left(z_{0}\right)^{*} X=2 \operatorname{tr}\left(\left(z_{0} I-A\right) X\right)
$$

It is well known (see, for example, [25]) that the positive semidefinite cone is Clarke regular at $F\left(z_{0}\right)$ (being convex), with normal cone

$$
\begin{aligned}
N_{\mathbf{H}_{+}^{n}}\left(F\left(z_{0}\right)\right) & =-\operatorname{cone}\left\{u u^{*}: F\left(z_{0}\right) u=0\right\} \\
& =-\operatorname{cone}\left\{u u^{*}: u \in U\left(z_{0}\right),\|u\|=1\right\}
\end{aligned}
$$

Now consider any matrix

$$
X \in N_{\mathbf{H}_{+}^{n}}\left(F\left(z_{0}\right)\right) \cap N\left(\nabla F\left(z_{0}\right)^{*}\right)
$$

By the calculations above, we deduce

$$
-X=\sum_{j=1}^{k} \mu_{j} u_{j} u_{j}^{*}
$$

for some integer $k$, reals $\mu_{j} \geq 0$, and unit vectors $u_{j} \in U\left(z_{0}\right)(j=1,2, \ldots, k)$, and

$$
0=\operatorname{tr}\left(\left(A-z_{0} I\right) X\right)=\sum_{j=1}^{k} \mu_{j} u_{j}^{*}\left(A-z_{0} I\right) u_{j}
$$

But since $0 \notin G\left(z_{0}\right)$, by Proposition 9.1 (convexity) this implies that each $\mu_{j}$ is zero. We have therefore proved the condition

$$
N_{\mathbf{H}_{+}^{n}}\left(F\left(z_{0}\right)\right) \cap N\left(\nabla F\left(z_{0}\right)^{*}\right)=\{0\} .
$$

Under this condition we can apply a standard chain rule [34] to the set of interest, $F^{-1}\left(\mathbf{H}_{+}^{n}\right)$, to deduce that it is Clarke regular at the point $z_{0}$, with normal cone

$$
N_{F^{-1}\left(\mathbf{H}_{+}^{n}\right)}\left(z_{0}\right)=\nabla F\left(z_{0}\right)^{*} N_{\mathbf{H}_{+}^{n}}\left(F\left(z_{0}\right)\right)=\operatorname{cone}\left(G\left(z_{0}\right)\right)
$$

as required.
Consider, for instance, Example 1 (nonsmooth points). When $\epsilon=\sqrt{2}$, we saw that the point $z_{0}=1$ is a nonsmooth point on the boundary of the pseudospectrum, which consists of the union of two disks of radius $\sqrt{2}$, centered at $\pm i$. A calculation shows that the set $G\left(z_{0}\right)$ in this case is the line segment $[1-i, 1+i]$, so according to the above result, the normal cone to the complement of the strict pseudospectrum is the cone $\{x+i y:|y| \leq-x\}$, as we expect.

By contrast, when $\epsilon=1$ we saw that the pseudospectrum consists of two unit disks, tangent at 0 . A calculation shows $G(0)$ is the line segment $[-i, i]$, which contains 0 , so the above theorem does not apply.

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