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Computational Noise: Uncertainty and Sensitivity

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Joint work with Stefan Wild

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Connections

- PhD, University of Illinois 1983
- ◊ Summer faculty position at Argonne, 1985.
- ◇ University of Washington, 1985 -

Publications

- On the identification of active constraints, 1988
- ◊ Convergence properties of trust region methods ... 1990
- Exposing contraints, 1994

Computational noise \rightarrow High Precision \rightarrow J. Borwein \rightarrow Jim Burke

High-precision computation: Mathematical physics and dynamics, Bailey, Barrio, Borwein, Applied Mathematics and Computation, 2012.



In the Beginning

The development of automatic digital computers has made it possible to carry out computations involving a very large number of arithmetic operations and this has stimulated a study of the cumulative effect of rounding errors.



J. H. Wilkinson, Rounding Errors in Algebraic Processes, 1963



The output of an algorithm \mathcal{A} is defined by $f : \mathbb{R}^n \mapsto \mathbb{R}$.

◇ f_∞(x): Output when A is executed in infinite precision
 ◇ f(x): Output when A is executed in finite precision

Rounding errors are measured by $|f_{\infty}(x) - f(x)|$

For backward-stable computations,

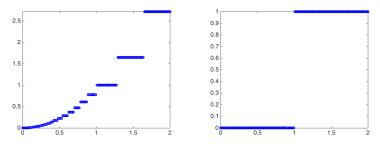
$$f_{\infty}(x+\delta x) = f(x)$$

for *small* perturbations δx .



Rounding Errors: A Cautionary Example

```
f = x;
for k = 1:L f = sqrt(f); end; for k = 1:L f = f<sup>2</sup>; end;
f = f<sup>2</sup>;
```



Plot of f for L = 50 (left) and $L \ge 55$ (right)

N. Higham, Accuracy and Stability of Numerical Algorithms, 2002W. Kahan, Interval arithmetic options in the proposed IEEE standard, 1980



Assessments of Roundoff

Repeat the computation but ...

- ◊ in higher precision
- with a different rounding mode
- with random rounding
- o use slightly different inputs
- use interval arithmetic

Monte Carlo arithmetic: How to gamble with floating point and win, D. S. Parker, B. Pierce and P. R. Eggert, Computing in Science Engineering, 2000

How futile are mindless assessments of roundoff in floating-point computation? W. Kahan, 2006. Work in progress, 56 pages.

CADNA: A library for estimating round-off error propagation, F. Jézéquél and J-M. Chesneaux, Computer Physics Communications, 2008.



The uncertainty in f is an estimate of

 $|f(x+\delta x) - f(x)|$

for a *small* perturbation δx .

 Δ If the computed f is backward-stable, then the uncertainty is

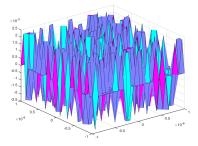
$$|f_{\infty}(x+\delta x_r) - f(x)|$$

for a small δx_r . This is an estimate of the rounding errors.

J. Moré and S. Wild, Estimating computational noise, SIAM Journal on Scientific Computing, 2011.



Research Issues



- \diamond What is a noisy function f?
- \diamond Determine the noise (uncertainty) in f with a few evaluations
- $\diamond\,$ Reliably approximate a derivative of f
- \diamond How do you optimize f?



Definition. The *noise level* of f in a region Ω is

$$\varepsilon_f = \mathrm{E}\left\{ rac{1}{2} \Big(f(\boldsymbol{x}_2) - f(\boldsymbol{x}_1) \Big)^2 \right\}^{1/2}, \quad \text{iid } \boldsymbol{x}_1, \boldsymbol{x}_2 \mapsto \Omega.$$

where Ω contains x_0 and all permissible perturbations of x_0

Leading causes of noise

- $\diamond \ 10^X \ {\rm flops}$
- ◊ Iterative calculations
- ◊ Adaptive algorithms
- Mixed precision



Theorem 1. If \mathcal{F} is the space of all $\mathbf{iid} \ \boldsymbol{x} \mapsto \Omega$,

$$arepsilon_f = \mathrm{Var}\left\{f(oldsymbol{x})
ight\}^{1/2} = \mathrm{E}\left\{|f(oldsymbol{x})-\mu|^2
ight\}^{1/2}, \qquad oldsymbol{x} \in \mathcal{F}$$

where μ is the expected value of $f(\boldsymbol{x})$.

Theorem 2. If f is a step function with values $v_1 \dots v_p$, then there are weights $w_k \ge 0$ with $\sum w_k = 1$ such that

$$\varepsilon_f = \left(\sum_{k=1}^p w_k (v_k - \mu)^2\right)^{1/2}$$

where μ is the (weighted) average of the function values.

The function $x \mapsto f[\operatorname{chop}(x,t)]$ is a step function.



Let μ be the expected value of $f(\pmb{x}).$

Chebyshev inequality

$$\mathcal{P}\left\{|f(\boldsymbol{x}) - \mu| \leq \gamma \varepsilon_f\right\} \geq 1 - \frac{1}{\gamma^2}, \qquad \gamma \geq 1.$$

Cauchy-Schwartz inequality

$$\mathrm{E}\left\{\left|f(\boldsymbol{x})-\boldsymbol{\mu}\right|\right\} \leq \varepsilon_f$$

Two Claims

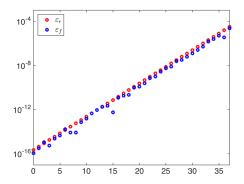
- $\diamond\,$ The noise level ε_f is a measure of the uncertainty of f
- $\diamond\,$ We can determine ε_f in a few function evaluations



Case Study: The Higham Function

$$f = x;$$

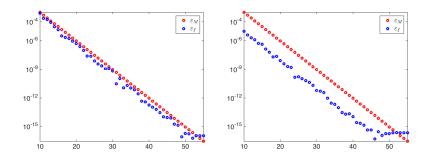
for k = 1:L f = sqrt(f); end; for k = 1:L f = f²; end;
f = f²;



$$\varepsilon_r(k) = 2^{(-k+52)}$$



Case Study: Mixed Precision Quadratics



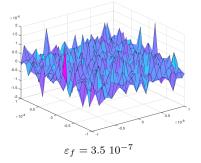
 $x \mapsto \|\operatorname{chop}(x,t)\|^2$ where $\operatorname{chop}(x,t)$ truncates $x \in \mathbb{R}^n$ to t bits n = 4 (left) and $n = 10^4$ (right)



Case Study: Eigenvalue Solvers

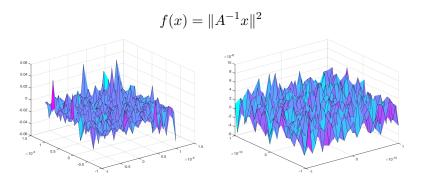
$$f(x) = \sum_{i=1}^{p} \lambda_i \Big(A + \operatorname{diag} (x) \Big), \quad p = 5$$

where $\lambda_i(\cdot)$ is the *i*-th smallest eigenvalue in magnitude.



A is the sparse Laplacian on an L-shaped region eigs with $\tau = 10^{-3}.$





bicgstab, $\varepsilon_f=7.8 \ 10^{-3}$ $$\rm pcg$, $\varepsilon_f=1.6 \ 10^{-8}$

Sparse matrix A from the University of Florida collection, $\tau=10^{-3}$ index = 38 (left) , index = 35 (right)



A realistic model of a finite-precision function is defined by

$$f(t) \equiv f_s[\boldsymbol{x}(t)], \qquad t \in [0, 1]$$

where $f_s : \mathbb{R} \mapsto \mathbb{R}$ is smooth $(f_s = f_\infty \text{ is an option})$

This model accounts for

- ◊ Changes in computer, software libraries, operating system,
- Code changes and reformulations
- Asynchronous, highly-concurrent algorithms
- Stochastic methods
- Variable/adaptive precision methods



$$f(t) = f_s(t) + \boldsymbol{\varepsilon}(t), \qquad t \in [0, 1]$$

 $\diamond~$ Construct the k-th order differences of f

$$\Delta^{k+1}f(t) = \Delta^k f(t+h) - \Delta^k f(t).$$

 $\diamond\,$ Note that ε_f can be determined from $\Delta^k \varepsilon(t)$

$$\gamma_k \mathbb{E}\left\{\left[\Delta^k \varepsilon(t)\right]^2\right\} = \varepsilon_f^2, \qquad \gamma_k = \frac{(k!)^2}{(2k)!}.$$

 \diamond Estimate the noise level of f from

$$\lim_{h\to 0} \gamma_k \, \mathbb{E}\left\{\left[\Delta^k f(t)\right]^2\right\} = \varepsilon_f^2,$$

R. W. Hamming, Introduction to Applied Numerical Analysis, 1971J. Moré and S. Wild, Estimating computational noise, 2011.



ECnoise: Difference Tables

1.56e+03 -6.92e+00 1.32e+00 3.40e+01 -1.21e+02 2.93e+02 -6.12e+02 1.56e+03 -5.61e+00 3.53e+01 -8.70e+01 1.72e+02 -3.19e+02 1.55e+03 2.97e+01 -5.17e+01 8.53e+01 -1.46e+02 1.58e+03 -2.20e+01 3.36e+01 -6.11e+01 1.56e+03 -1.59e+03 -1.59e+01 1.56e+03

Noise levels

1.24e+01 1.39e+01 1.57e+01 1.77e+01 1.93e+01 2.01e+01

bicgstab, index = 38, $\varepsilon_f = 7.8 \ 10^{-3}$

5.29e+02 7.53e-06 -5.74e-06 -8.38e-06 3.41e-05 -4.74e-05 -7.65e-06 5.29e+02 1.80e-06 -1.41e-05 2.58e-05 -1.33e-05 -5.51e-05 5.29e+02 -1.23e-05 1.16e-05 1.25e-05 -6.84e-05 5.29e+02 2.34e-07 2.41e-05 -5.59e-05 5.29e+02 2.34e-05 -3.18e-05 5.29e+02 -8.38e-06 5.29e+02 -8.38e-06

Noise levels

8.32e-06 8.08e-06 7.08e-06 5.35e-06 3.24e-06 2.52e-07

pcg, index = 35,
$$\varepsilon_f = 1.6 \ 10^{-8}$$



Define $f_{\tau}: \mathbb{R}^n \mapsto \mathbb{R}$ as the iterative solution of a Krylov solver,

$$f_{\tau}(x) = \|y_{\tau}(x)\|^2, \qquad \|Ay_{\tau}(x) - b\| \le \tau \|b\|,$$

where b is a function of the input x. We use b = x.

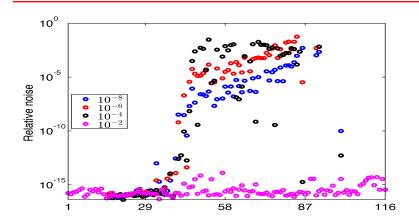
 Δ $y_{\tau}: \mathbb{R}^n \mapsto \mathbb{R}^n$ is continuously differentiable for almost all τ

- \diamond UF symmetric positive definite matrices (116) with $n \leq 10^4$
- ♦ Scaling: $A \leftarrow D^{-1/2}AD^{-1/2}$, $D = diag(a_{i,i})$
- ◇ Solvers: bicgstab (similar results for pcg, minres, gmres, ...)

♦ Tolerances:
$$\tau \in [10^{-8}, 10^{-1}]$$



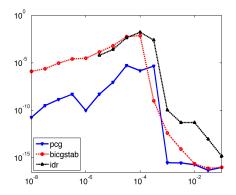
What is the Noise Level of Krylov Solvers?



Distribution of ε_f for f_{τ} (bicgstab)



Noise Level Transitions



 ε_f as a function of tolerance τ



We measure the uncertainty of f' with

$$\mathsf{re}(f') = \mathsf{re}\Big\{f'(x_0; p), f'(x_0; (1+\varepsilon)p)\Big\}, \qquad \varepsilon = \varepsilon_M.$$

- $\diamond~$ With infinite precision ${\rm re}(f')=\varepsilon/(1+\varepsilon)$ for any $\varepsilon>0$
- $\diamond\,$ We expect $\mathrm{re}(f')$ to be small in floating point arithmetic.

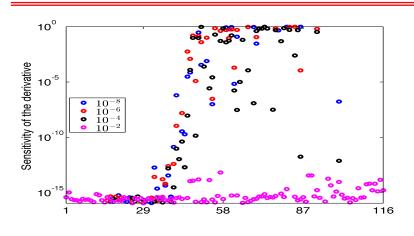
Two AD algorithms (forward mode) were used to compute f':

- IntLab (Siegfried Rump, Hamburg)
- AdiMat (Andre Vehreschild, Aachen)

IntLab was used in the numerical results.



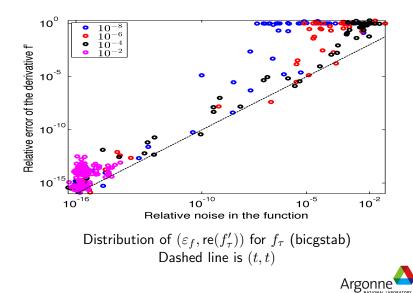
Can You Trust Derivatives?



Distribution of $re(f'_{\tau})$ for f_{τ} (bicgstab)



 $\mathrm{re}(f_\tau')\gg \varepsilon_f$



Define the RMS expected error in the derivative by

$$\mathcal{E}(h) = \mathbf{E} \left\{ \left(\frac{f(t_0 + h) - f(t_0)}{h} - f'_s(t_0) \right)^2 \right\}^{1/2}$$

Theorem. If $\mu_L = \min |f''|$, $\mu_M = \max |f''|$, and

$$h^* = \gamma_2 \left(\frac{\varepsilon_f}{\mu}\right)^{1/2}, \qquad \gamma_2 = 8^{1/4} \approx 1.68,$$

where μ is an estimate of $|f''(t_0)|$ in $[\mu_{\scriptscriptstyle L},\mu_{\scriptscriptstyle M}]$, then

$$\mathcal{E}(h^*) \le (\gamma_1 \mu_M \,\varepsilon_f)^{1/2} \le \left(\frac{\mu_M}{\mu_L}\right)^{1/2} \min_{0 < h \le h_0} \mathcal{E}(h)$$



The estimate μ of $|f''(t_0)|$ is obtained from

$$\mu = \frac{f(t_0 - h) - 2f(t_0) + f(t_0 + h)}{h^2}$$

where h is of order $\varepsilon_f^{1/4}$. Set

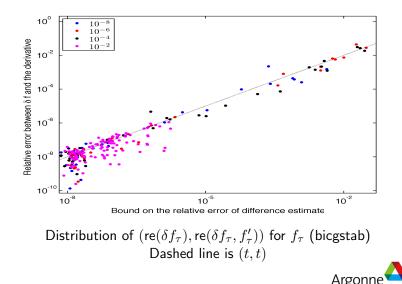
$$\delta f(t_0) = \frac{f(t_0 + h^*) - f(t_0)}{h^*}$$

Claim: If $re(\delta f, f')$ is the relative error between δf and f', then

$$\operatorname{re}(\delta f, f') \sim \operatorname{re}(\delta f) \equiv \frac{\mathcal{E}(h^*)}{|f'(t_0)|}$$



$\operatorname{re}(\delta f)\sim\operatorname{re}(\delta f,f')$



S. Wild, Estimating Computational Noise in Numerical Simulations www.mcs.anl.gov/~wild/cnoise

- J. Moré and S. Wild, *Estimating Computational Noise*, SIAM Journal on Scientific Computing, 33 (2011), 1292-1314.
- J. Moré and S. Wild, Estimating Derivatives of Noisy Simulations, ACM Trans. Mathematical Software, 38 (2012), 19:1–19:21
- J. Moré and S. Wild, *Do You Trust Derivatives or Differences?*, Journal of Computational Physics, 273 (2014), 268 – 277.

