

## Problem Set 8

**514 - Networks and Combinatorial Optimization**

Autumn 2022

**Exercise 5.2 (10pts)**

Let  $G = (V, E)$  be a 3-regular graph without any bridge. Show that  $G$  has a perfect matching. (A *bridge* is an edge  $e$  not contained in any circuit; equivalently, deleting  $e$  increases the number of components; equivalently  $\{e\}$  is a cut.)

**Exercise 5.4 (10pts)**

Let  $G = (V, E)$  be a graph and let  $T$  be a subset. Then  $G$  has a matching covering  $T$  if and only if the number of odd components of  $G \setminus W$  contained in  $T$  is at most  $|W|$ , for each  $W \subseteq V$ .

**Exercise 5.23 (10pts)**

Recall that for a graph  $G = (V, E)$  we have defined  $P_{\text{matching}}(G) = \text{conv}\{\chi^M \in \mathbb{R}^E \mid M \subseteq E \text{ is a matching}\}$ . Prove that for any  $k \in \mathbb{N}$ ,

$$P_{\text{matching}}(G) \cap \{x \in \mathbb{R}^E \mid \mathbf{1}^T x = k\} = \text{conv}\{\chi^M \mid M \subseteq E \text{ is a matching with } |M| = k\}$$

**Hint:** Let  $\mathcal{F} := \{M \subseteq E \mid M \text{ is a matching}\}$ . For the non-trivial direction, take a vector  $x^* \in P_{\text{matching}}(G) \cap \{x \in \mathbb{R}^E \mid \mathbf{1}^T x = k\}$  and consider the vector  $\lambda \in \mathbb{R}_{\geq 0}^{\mathcal{F}}$  that minimizes the function  $G(\lambda) := \sum_{M \in \mathcal{F}} \lambda_M \cdot \left| |M| - k \right|$  subject to  $\sum_{M \in \mathcal{F}} \lambda_M = 1$  and  $x^* = \sum_{M \in \mathcal{F}} \lambda_M \chi^M$ .

**Remark.** All three exercises are taken from A. Schrijver's lecture notes.