

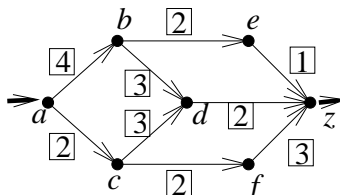
[Homework can be handed in to me or to my mail box in the Math Lounge (opposite the Math main office). Please show your work to receive full credit.]

**A.** Consider a digraph  $G = (V, A)$  and two vertices  $s \neq t \in V$ .

(a) To hedge against possible arc failures, we wish to find the largest number of paths from  $s$  to  $t$  that are arc-disjoint (i.e., no two paths share any arc). Formulate this as a maximum flow problem. Explain your answer. [Hint: Use unit capacities.]

(b) If instead of the paths being arc-disjoint, we require that at most  $k$  paths ( $k$  a positive integer) share any arc. How would your answer to (a) change?

**B.** Consider the capacitated digraph shown below:



(a) Apply the algorithm MAXFLOW to find a maximum flow and minimum cut from  $a$  to  $z$ .

(b) For the  $a$ - $z$  flow  $x_{ab} = x_{be} = x_{ez} = 1$ ,  $x_{ac} = x_{cd} = x_{dz} = 2$  (with zero flow on remaining arcs), find the residual digraph and an  $a$ - $z$  augmenting path. Augment flow along the augmenting path to obtain a new  $a$ - $z$  flow.

**C.** Prove that if  $\{x_{uv}\}_{(u,v) \in A}$  is an  $s$ - $t$  flow of value  $v$  and  $[S, T]$  is an  $s$ - $t$  cut of capacity  $c$  (with  $T = V \setminus S$ ), then  $v = c$  if and only if  $x_{uv} = c_{uv}$  for every  $(u, v) \in [S, T]$  and  $x_{vu} = 0$  for every  $(v, u) \in [T, S]$ . [You can use the fact that  $v = \sum_{(u,v) \in [S, T]} x_{uv} - \sum_{(v,u) \in [T, S]} x_{vu}$  for any  $s$ - $t$  flow  $(x_{uv})_{(u,v) \in A}$  of value  $v$ .]

**D.**

(a) Give an example of a digraph  $G = (V, A)$  with arc capacities of 1 and two vertices  $s, t$  such that  $|V| \leq 4$  and the  $s$ - $t$  flow of maximum value is not unique.

(b) For your answer to (a). Find an  $s$ - $t$  flow of maximum value whose arc flows are *not* all integer.

(c) Give an example of a digraph  $G = (V, A)$  with arc capacities of 1 and two vertices  $s, t$  such that  $|V| = 3$  and the  $s$ - $t$  cut of minimum capacity is not unique.

(d) Consider a digraph  $G = (V, A)$  with arc capacities  $c_{uv}$ ,  $(u, v) \in A$ , and  $s \neq t \in V$ . Suppose  $u_1, u_2, u_3, u_4$  is an  $s$ - $t$  augmenting path relative to an  $s$ - $t$  flow  $(x_{uv})_{(u,v) \in A}$ . Suppose  $(u_2, u_3)$  is a forward arc and  $(u_2, u_1), (u_4, u_3)$  are backward arcs. Write down a formula for the augmentation amount  $\Delta$  in terms of the flow and capacity of these arcs.

**Bonus.** Consider a digraph  $G = (V, A)$  with arc capacities  $c_{uv}$ ,  $(u, v) \in A$ , and  $s \neq t \in V$ . Suppose that  $[S_1, T_1]$  and  $[S_2, T_2]$  are two  $s$ - $t$  cuts of minimum capacity. Prove that  $[S_1 \cap S_2, T_1 \cup T_2]$  and  $[S_1 \cup S_2, T_1 \cap T_2]$  are also  $s$ - $t$  cuts of minimum capacity. [Hint: First show that they are  $s$ - $t$  cuts. Then show that their capacities equal that of  $[S_1, T_1]$  and  $[S_2, T_2]$  by considering arcs out of  $S_1 \cap S_2$ ,  $S_1 \setminus S_2$ ,  $S_2 \setminus S_1$ .]