RECONSTRUCTION THEORY

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ABSTRACT

Suppose that G is a compact group. Denote by Rep G the category whose objects are the continuous finite dimensional unitary representations of G and whose morphisms are the intertwining operators — then Rep G is a monoidal *-category with certain properties P_1, P_2, \ldots Conversely, if C is a monoidal *-category possessing properties P_1, P_2, \ldots can one find a compact group G, unique up to isomorphism, such that Rep G "is" C? The central conclusion of reconstruction theory is that the answer is affirmative.

CONTENTS

- §L. MONOIDAL CATEGORIES
- §2. MONOIDAL FUNCTORS
- §3. STRICTIFICATION
- §4. SYMMETRY
- §5. DUALITY
- §6. TWISTS
- §7. *-CATEGORIES
- §8. NATURAL TRANSFORMATIONS
- §9. CONJUGATES
- \$10. TANNAKIAN CATEGORIES
- §11. FIBER FUNCTORS
- §12. THE INTRINSIC GROUP
- §13. CLASSICAL THEORY

§1. MONOIDAL CATEGORIES

Given categories C,D, their product is the category $C \times D$ defined by

$$Ob(\underline{C} \times \underline{D}) = Ob \underline{C} \times Ob \underline{D}$$

$$Mor((X,Y),(X',Y')) = Mor(X,X') \times Mor(Y,Y')$$

$$id_{X \times Y} = id_{X} \times id_{Y},$$

with composition

$$(f',g') \circ (f,g) = (f' \circ f,g' \circ g).$$

Now take C = D — then a <u>monoidal category</u> is a category C equipped with a functor $C \times C \times C \to C$ (the <u>multiplication</u>) and an object $C \times C \to C$ (the <u>unit</u>), together with natural isomorphisms $C \times C \to C$ where

$$\begin{bmatrix} R_{X}:X \otimes e \to X \\ L_{X}:e \otimes X \to X \end{bmatrix}$$

and

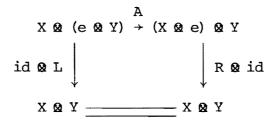
$$A_{X,Y,Z}$$
: X & (Y & Z) \rightarrow (X & Y) & Z,

subject to the following assumptions.

 (MC_1) The diagram

commutes.

(MC₂) The diagram



commutes.

[Note: The "coherency" principle then asserts that "all" diagrams built up from instances of R, L, A (or their inverses), and id by repeated application of 20 necessarily commute. In particular, the diagrams

commute and $L_e = R_e : e \otimes e \rightarrow e$.

N.B. Technically, the categories

$$(\overline{C} \times \overline{C}) \times \overline{C}$$

$$\overline{C} \times (\overline{C} \times \overline{C})$$

are not the same so it doesn't quite make sense to say that the functors

$$(\underline{C} \times \underline{C}) \times \underline{C} \rightarrow \underline{C}$$

$$((f,g),h) \rightarrow (f \otimes g) \otimes h$$

are naturally isomorphic. However, there is an obvious isomorphism

$$\overline{C} \times (\overline{C} \times \overline{C}) \xrightarrow{1} (\overline{C} \times \overline{C}) \times \overline{C}$$

and the assumption is that A:F \rightarrow G \circ 1 is a natural isomorphism, where

$$\begin{array}{ccc}
\underline{C} \times (\underline{C} \times \underline{C}) & \xrightarrow{F} \underline{C} \\
\downarrow & \downarrow \\
(\underline{C} \times \underline{C}) \times \underline{C} & \xrightarrow{G} \underline{C}.
\end{array}$$

Accordingly,

$$\forall$$
 (X,(Y,Z)) \in Ob $C \times (C \times C)$

and

$$\forall$$
 (f,(g,h)) \in Mor $\underline{C} \times (\underline{C} \times \underline{C})$,

the square

commutes.

Interchange Principle If

then

$$(f \otimes id_{X}) \circ (id_{X} \otimes g) = f \otimes g = (id_{X'} \otimes g) \circ (f \otimes id_{Y}).$$

[Note: Since $\mathfrak{A}: \underline{C} \times \underline{C} \to \underline{C}$ is a functor, in general

$$(f \circ f') \otimes (g \circ g') = (f \otimes g) \circ (f' \otimes g').$$

1.1 EXAMPLE Given a field \underline{k} , let $\underline{\text{VEC}}_{\underline{k}}$ be the category whose objects are the vector spaces over \underline{k} and whose morphisms are the linear transformations — then $\underline{\text{VEC}}_{\underline{k}}$ is monoidal: Take X $\underline{\text{W}}$ Y to be the algebraic tensor product and let e be \underline{k} .

[Note: If

$$f:X \to X'$$

$$g:Y \to Y',$$

then

$$\otimes$$
 (f,g) = f \otimes g:X \otimes Y \rightarrow X' \otimes Y'

sends $x \otimes y$ to $f(x) \otimes g(y)$.

Let H and K be complex Hilbert spaces — then their algebraic tensor product H Ω K can be equipped with an inner product given on elementary tensors by

$$\langle x_1 \otimes y_1, x_2 \otimes y_2 \rangle = \langle x_1, x_2 \rangle \langle y_1, y_2 \rangle$$

and its completion $\mathcal{H} \overset{\mathbf{Q}}{=} \mathcal{K}$ is a complex Hilbert space.

N.B. If

$$\begin{bmatrix} A \in \mathcal{B}(H_1, H_2) \\ B \in \mathcal{B}(K_1, K_2), \end{bmatrix}$$

then

$$A \otimes B: H_1 \otimes K_1 \rightarrow H_2 \otimes K_2$$
.

extends by continuity to a bounded linear operator

$$A \underline{\otimes} B: H_1 \underline{\otimes} K_1 \rightarrow H_2 \underline{\otimes} K_2.$$

Denote by <u>HILB</u> the category whose objects are the complex Hilbert spaces and whose morphisms are the bounded linear operators.

1.2 EXAMPLE HILB is a monoidal category.

PROOF Define a functor

by

aml

and let e be C.

1.3 REMARK Both $\underline{\text{Vec}}_{\underline{k}}$ and $\underline{\text{HILB}}$ admit a second monoidal structure: Take for the multiplication the direct sum Φ and take for the unit the zero object $\{0\}$.

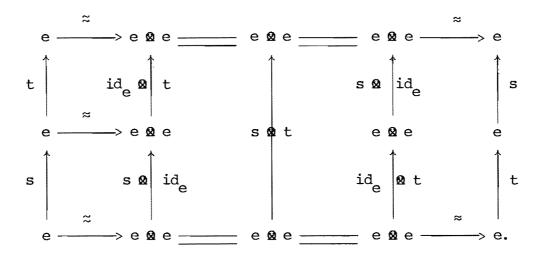
Put

$$M(C) = Mor(e,e)$$
.

Then $\underline{M}(\underline{C})$ is a monoid with categorical composition as monoid multiplication.

1.4 LEMMA The monoid $\underline{M}(\underline{C})$ is commutative.

PROOF Take $s,t \in M(C)$ and consider the commutative diagram



Then

$$R_e^{-1} \circ (s \circ t) \circ R_e = R_e^{-1} \circ (t \circ s) \circ R_e$$

$$\Rightarrow s \circ t = t \circ s.$$

Given $f \in Mor(X,Y)$ and $s \in \underline{M}(\underline{C})$, define $s \cdot f$ to be the composition

$$L^{-1} \qquad s \otimes f \qquad L$$

$$X \longrightarrow e \otimes X \longrightarrow e \otimes Y \longrightarrow Y.$$

1.5 LEMMA We have

$$id_{e} \cdot f = f$$

$$s \cdot (t \cdot f) = (s \circ t) \cdot f$$

$$(t \cdot g) \circ (s \cdot f) = (t \circ s) \cdot (g \circ f)$$

$$(s \cdot f) \otimes (t \cdot g) = (s \circ t) \cdot (f \otimes g).$$

A monoidal category \underline{C} is said to be <u>strict</u> if R, L, and A are identities. So, if \underline{C} is strict, then

$$X \otimes (Y \otimes Z) = (X \otimes Y) \otimes Z$$

ami

[Note: While monoidal, neither $\underline{\text{VEC}}_k$ nor $\underline{\text{HILB}}$ is strict monoidal.]

 $\underline{\text{N.B.}} \quad \text{Take \underline{C} strict and consider $\underline{M}(\underline{C})$ $--$ then \forall f,g \in \underline{M}(\underline{C})$,}$

$$f \Omega g = f \circ g = g \circ f = g \Omega f.$$

1.6 <u>EXAMPLE</u> Let \$ be the category whose objects are the nonnegative integers and whose morphisms are specified by the rule

composition in Mor(n,n) being group multiplication in $\$_n$. Define

on objects by

$$\Omega(n,m) = n + m$$

and on morphisms by

$$\otimes (n \rightarrow n, m \rightarrow m) = \rho_{n,m}(\sigma, \tau),$$

where

$$\rho_{n,m}: \mathcal{S}_n \times \mathcal{S}_m \rightarrow \mathcal{S}_{n+m}$$

is the canonical map, i.e.,

and let e = 0 -- then with these choices, \$ is a strict monoidal category.

[Note: \$ is equivalent to the category whose objects are the finite sets and whose morphisms are the bijective maps.]

1.7 $\underline{\text{EXAMPLE}}$ Let $\underline{\text{MAT}}_k$ be the category whose objects are the positive integers and whose morphisms are specified by the rule

$$Mor(n,m) = M_{n,m}(\underline{k}),$$

the n-by-m matrices with coefficients in \underline{k} . Here $\mathrm{id}_n: n \to n$ is the unit diagonal n-by-n matrix and composition

$$\circ: Mor(n,m) \times Mor(m,p) \rightarrow Mor(n,p)$$

is

$$B \circ A = AB$$
,

the product on the right being ordinary multiplication of matrices. Define

$$\mathbf{MAT}_{\underline{k}} \times \underline{\mathbf{MAT}}_{\underline{k}} \to \underline{\mathbf{MAT}}_{\underline{k}}$$

on objects by

$$Q(n,m) = nm$$

and on morphisms by

A B
$$\mathfrak{Q}(n \to m, p \to q)$$

$$= \begin{bmatrix} a_{11}^{B} & \cdots & a_{1m}^{B} \\ \vdots & & \vdots \\ a_{n1}^{B} & \cdots & a_{nm}^{B} \end{bmatrix} \in Mor(np,mq)$$

and let e = 1 -- then with these choices, \underline{MAT}_k is a strict monoidal category.

[Note: Write $\underline{FDVEC}_{\underline{k}}$ for the full subcategory of $\underline{VEC}_{\underline{k}}$ whose objects are finite dimensional — then there is an equivalence $\underline{MAT}_{\underline{k}} \to \underline{FDVEC}_{\underline{k}}$. Thus assign to each object n the vector space \underline{k}^n and to each morphism $n \to m$ the linear map from \underline{k}^n to \underline{k}^m that sends $(x_1, \ldots, x_n) \in \underline{k}^n$ to $(y_1, \ldots, y_m) \in \underline{k}^m$, where y_i is the i^{th} entry of the 1-by-m matrix $[x_1, \ldots, x_n]A$.]

1.8 EXAMPLE Given a C*-algebra A, let End A be the category whose objects are the unital *-homomorphisms $\Phi: A \to A$ and whose arrows $\Phi \to \Psi$ are the intertwiners, i.e.,

$$Mor(\Phi, \Psi) = \{ \mathbf{T} \in A : \mathbf{T}\Phi(\mathbf{A}) = \Psi(\mathbf{A})\mathbf{T} \ \forall \ \mathbf{A} \in A \}.$$

Here, the composition of arrows, when defined, is given by the product in A and $l_A \in Mor(\Phi, \Phi)$ is $l_{\bar{\Phi}}$. Define

on objects by

and on morphisms by

and let $e = id_A$ -- then with these choices, End A is a strict monoidal category.

[Note: $\forall A \in A$, we have

$$T\Phi(T')(\Phi \circ \Phi')(A)$$

$$= T\Phi(T')\Phi(\Phi'(A))$$

$$= T\Phi(T'\Phi'(A))$$

$$= \Psi(T'\Phi'(A))T$$

$$= \Psi(\Psi'(A)T')T$$

$$= \Psi(\Psi'(A))\Psi(T')T$$

$$= (\Psi \circ \Psi')(A)T\Phi(T').$$

1.9 EXAMPLE Given a category \underline{C} , let $[\underline{C},\underline{C}]$ be the metacategory whose objects are the functors $F:\underline{C} \to \underline{C}$ and whose morphisms are the natural transformations Ξ from F to G. Define

on objects by

and on morphisms by

where

$$(\exists \ \& \ \exists')_{X}$$

$$= \exists_{G'X} \circ F\exists_{X}' \ (= G\exists_{X}' \circ \exists_{F'X}),$$

and let $e = id_{\underline{C}}$ (the identity functor) — then with these choices, $[\underline{C},\underline{C}]$ is a strict monoidal category.

[Note: If

then

there are commutative diagrams

$$\begin{array}{ccc} & \xrightarrow{\Xi_{X}} & \text{GX} \\ & & \downarrow & \text{Gf} \\ & \text{FY} & \xrightarrow{\Xi_{Y}} & \text{GY} \end{array}$$

$$F'X' \xrightarrow{\Xi'_{X'}} G'X'$$

$$F'f' \downarrow \qquad \qquad \downarrow G'f'$$

$$F'Y' \xrightarrow{\Xi'_{Y'}} G'Y'.$$

In particular: The diagram

$$FF'X \xrightarrow{\Xi} GF'X$$

$$F\Xi'_{X} \downarrow \qquad \qquad \downarrow G\Xi'_{X}$$

$$FG'X \xrightarrow{\Xi} GG'X$$

$$G'X$$

commutes. This said, the claim is that

$$E \otimes E' \in Nat(F \circ F', G \circ G'),$$

i.e., that the diagram

$$FF'X \xrightarrow{(\Xi \ \Omega \ \Xi')_X} GG'X$$

$$FF'f \downarrow \qquad \qquad \downarrow GG'f$$

$$FF'Y \xrightarrow{(\Xi \ \Omega \ \Xi')_Y} GG'Y$$

commutes. In fact,

$$GG'f \circ (\Xi \boxtimes \Xi')_{X}$$

$$= GG'f \circ \Xi \qquad \circ F\Xi'_{X}$$

$$= GG'f \circ G\Xi'_{X} \circ \Xi$$

$$F'X$$

1.10 $\underline{\text{LEMMA}}$ Suppose that \underline{C} is monoidal and let e,e' be units -- then e and e' are isomorphic.

[There is an isomorphism $\phi:e \rightarrow e'$ for which the diagrams

commute, viz.

$$\phi = L \circ (R_e')^{-1} \quad (e \rightarrow e \otimes e' \rightarrow e').]$$

§2. MONOIDAL FUNCTORS

Let \underline{C} , \underline{C} ' be monoidal categories — then a <u>monoidal functor</u> is a triple (F,ξ,Ξ) , where $F:\underline{C}\to\underline{C}$ ' is a functor, $\xi:e'\to Fe$ is an isomorphism, and the

$$\Xi_{X,Y}$$
: FX \(\text{SY}\) \(\text{FY}\) \(\text{F}\) \(\text{X}\) \(\text{Y}\)

are isomorphisms, natural in X,Y, subject to the following assumptions.

(MF₁) The diagram

commutes.

(MF₂) The diagrams

commute.

N.B. A monoidal functor is said to be strict if ξ and Ξ are identities.

2.1 EXAMPLE Write FDHILB for the full subcategory of HILB whose objects are finite dimensional — then the forgetful functor

is strict monoidal.

[Take for

$$\Xi_{X,Y}$$
: UX \otimes UY \rightarrow U(X \otimes Y)

the identity $id_{X \otimes Y}$ and let $\xi = id_{C^*}$

[Note: A forgetful functor need not be monoidal, let alone strict monoidal. E.g.: Give \underline{AB} its monoidal structure per the tensor product, give \underline{SET} its monoidal structure per the cartesian product, and consider $\underline{U:AB} \rightarrow \underline{SET}$ -- then the canonical maps

are not isomorphisms.]

Let

(F,
$$\xi$$
, Ξ)
(G, θ , Θ)

be monoidal functors -- then a monoidal natural transformation

$$(F,\xi,\Xi) \rightarrow (G,\theta,\Theta)$$

is a natural transformation $\alpha:F \to G$ such that the diagrams

commute.

Write $[\underline{C},\underline{C}']^{\underline{Q}}$ for the metacategory whose objects are the monoidal functors $\underline{C} + \underline{C}'$ and whose morphisms are the monoidal natural transformations.

- N.B. A monoidal natural transformation is a monoidal natural isomorphism if α is a natural isomorphism.
- 2.2 <u>REMARK</u> Some authorities assume outright that Fe = e', the rationale being that this can always be achieved by replacing $F \in Ob \ [\underline{C},\underline{C}']^{\Omega}$ by an isomorphic $F' \in Ob \ [\underline{C},\underline{C}']^{\Omega}$ such that F'e = e' (on objects $X \neq e$, F'X = FX).

2.3 LEMMA Let

$$(F,\Xi,\xi) \qquad (F:\underline{C} \to \underline{C}')$$

$$(F',\Xi',\xi') \qquad (F':\underline{C}' \to \underline{C}'')$$

be moroidal functors — then their composition $F' \, \circ \, F$ is a monoidal functor.

[Consider the arrows e'' $\xrightarrow{\xi'}$ F'e' $\xrightarrow{F'\xi}$ F'Fe and

F'FX Q'' F'FY
$$\xrightarrow{\Xi'}$$
 F'(FX Q' FY) $\xrightarrow{F'\Xi}$ F'F(X Q Y).]

Write MONCAT for the metacategory whose objects are the monoidal categories and whose morphisms are the monoidal functors.

2.4 <u>RAPPEL</u> Let \underline{C} , \underline{D} be categories — then a functor $F:\underline{C} \to \underline{D}$ is said to be an <u>equivalence</u> if there exists a functor $G:\underline{D} \to \underline{C}$ such that $G \circ F \simeq \operatorname{id}_{\underline{C}}$ and $F \circ G \simeq \operatorname{id}_{\underline{D}}$, the symbol \simeq standing for natural isomorphism.

- 2.5 <u>LEMMA</u> A functor $F:\underline{C} \to \underline{D}$ is an equivalence iff it is full, faithful, and has a representative image (i.e., for any $Y \in Ob \ \underline{D}$, there exists an $X \in Ob \ \underline{C}$ such that FX is isomorphic to Y).
- N.B. Categories \underline{C} , \underline{D} are said to be <u>equivalent</u> provided there is an equivalence $F:\underline{C} \to \underline{D}$. The object isomorphism types of equivalent categories are in a one-to-one correspondence.

isomorphic, i.e., if it is possible to assign to each ordered pair $\begin{vmatrix} -& x \in Ob \ \underline{C} \\ & y \in Ob \ \underline{D} \end{vmatrix}$

a bijective map $\Xi_{X,Y}$:Mor(FX,Y) \rightarrow Mor(X,GY) which is functorial in X and Y. When this is so, F is a <u>left adjoint</u> for G and G is a <u>right adjoint</u> for F. Any two left (right) adjoints for G (F) are naturally isomorphic. In order that (F,G) be an adjoint pair, it is necessary and sufficient that there exist natural trans-

The data (F,G,μ,ν) is referred to as an <u>adjoint situation</u>, the natural trans-

alence of categories is an adjoint situation (F,G,μ,ν) in which both μ and ν are natural isomorphisms.

2.7 <u>LEMMA</u> A functor $F:C \rightarrow D$ is an equivalence iff F is part of an adjoint equivalence.

Let C, C' be monoidal categories — then C, C' are monoidally equivalent if there are monoidal functors

and monoidal natural isomorphisms

- 2.8 <u>LEMMA</u> Suppose that $F:\underline{C} \to \underline{C}'$ is a monoidal functor. Assume: F is an equivalence then F is a monoidal equivalence.
 - 2.9 REMARK Embed F in an adjoint situation (F,F',μ,μ') , where

$$\mu: \operatorname{id}_{\underline{C}} \to F' \circ F$$

$$\mu': F \circ F' \to \operatorname{id}_{\underline{C}'}$$

are the arrows of adjunction (cf. 2.7) — then one can equip F' with the structure of a monoidal functor in such a way that the natural isomorphisms μ , μ ' are monoidal natural isomorphisms. Thus first specify $\xi':e \to F'e'$ by taking it to

be the composition $e \xrightarrow{\mu_e} F'\xi^{-1}$ be the composition $e \xrightarrow{\mu_e} F'Fe \xrightarrow{F'} F'e'$. As for

$$\Xi^{\dagger}$$
; $F^{\dagger}X^{\dagger} \otimes F^{\dagger}Y^{\dagger} \rightarrow F^{\dagger}(X^{\dagger} \otimes^{\dagger} Y^{\dagger})$, X^{\dagger}, Y^{\dagger}

build it in three stages:

1.
$$F'X' \otimes F'Y' \rightarrow F'F(F'X' \otimes F'Y')$$
;

2.
$$F'F(F'X' \otimes F'Y') \xrightarrow{F'\Xi^{-1}} F'(FF'X' \otimes' FF'Y');$$

=>

=>

If C is monoidal, then $\underline{C}^{\mathsf{OP}}$ is monoidal when equipped with the same \mathbf{Q} and \mathbf{e} , taking

$$\begin{bmatrix}
 & R^{OP} = R^{-1} \\
 & L^{OP} = L^{-1} \\
 & A^{OP} = A^{-1}
 \end{bmatrix}$$

§3. STRICTIFICATION

A <u>strictification</u> of a moroidal category \underline{C} is a strict monoidal category which is monoidally equivalent to \underline{C} .

3.1 EXAMPLE $\underline{\text{MAT}}_{\underline{k}}$ is a strictification of $\underline{\text{FDVEC}}_{\underline{k}}$.

[The equivalence $\underline{\text{MAT}}_{\underline{k}} \to \underline{\text{FDVEC}}_{\underline{k}}$ constructed in 1.7 is a monoidal functor, hence is a monoidal equivalence (cf. 2.8).]

3.2 THEOREM Every monoidal category $\underline{\mathbf{C}}$ is monoidally equivalent to a strict monoidal category $\underline{\mathbf{C}}_{\text{str}}.$

The proof is constructive and best broken up into steps.

Step 1: Let \underline{S} be the class of all finite sequences $S = (X_1, \dots, X_n)$ of objects of \underline{C} , including the empty sequence \emptyset . Given nonempty

$$T = (Y_{1}, ..., Y_{m})$$

let

$$s * T = (x_1, ..., x_n, x_1, ..., x_m)$$

and write

$$S \star \emptyset = S = \emptyset \star S.$$

Step 2: The claim is that S is the object class of a strict monoidal

category $\underline{C}_{\underline{str}}$, i.e., $\underline{S} = Ob \ \underline{C}_{\underline{str}}$. In any event, the multiplication

$$\star:S \times S \rightarrow S$$

is associative, so we can take A to be the identity. Also, \emptyset serves as the unit and

$$\begin{bmatrix} - & R_S : S * \emptyset \rightarrow S \\ & L_S : \emptyset * S \rightarrow S \end{bmatrix}$$

are the identities.

Step 3: Given S, T, we need to specify Mor(S,T). For this purpose, define a map $\Gamma: S \to Ob \subset Dy \Gamma \emptyset = e$, $\Gamma((X)) = X$, and $\Gamma(S \star (X)) = \Gamma S \otimes X$, thus

$$\Gamma(X_1, \dots, X_n)$$

$$= (\dots(X_1 \otimes X_2) \otimes \dots) \otimes X_n,$$

where all opening parentheses are to the left of \mathbf{X}_1 . Definition:

$$Mor(S,T) = Mor(\Gamma S, \Gamma T)$$
.

This prescription then gives rise to a category \underline{C}_{str} with 0b $\underline{C}_{str} = \underline{S}$.

Step 4: We shall now define a functor $*:\underline{C}_{\underline{str}} \times \underline{C}_{\underline{str}} \to \underline{C}_{\underline{str}}$ that serves to render $\underline{C}_{\underline{str}}$ strict monoidal, the issue being the meaning of

$$u * u^{\dagger} = *(S \rightarrow T, S^{\dagger} \rightarrow T^{\dagger})$$

$$\in Mor(S * S^{\dagger}, T * T^{\dagger})$$

$$= Mor(\Gamma(S * S^{\dagger}), \Gamma(T * T^{\dagger})).$$

Bearing in mind that

let u * u' be the composite

where the outer arrows are the obvious canonical morphisms in \underline{C} . Accordingly, with this agreement, $\underline{C}_{\mathtt{str}}$ is strict monoidal.

Step 5: It is clear from its very construction that $\Gamma:\underline{C}_{\underline{\operatorname{str}}}\to\underline{C}$ is a functor which, moreover, is full, faithful, and is isomorphism dense. But $\Gamma\emptyset$ = e and there are isomorphisms

$$E_{S,T}$$
: $\Gamma S \otimes \Gamma T \rightarrow \Gamma (S * T)$,

natural in S, T and satisfying M Γ_1 , M Γ_2 of §2. Therefore Γ is monoidal. To finish, it remains only to quote 2.8.

[Note: It is not necessary to quote 2.8: Simply observe that there is an inclusion functor $\gamma\colon\!\underline C\to\underline C_{\hbox{\rm str}}$ and

$$\Gamma \circ \gamma = id_{\underline{C}}$$

$$\gamma \circ \Gamma \approx id_{\underline{C}}$$

$$- \underbrace{\text{str}}$$

Detail: From

$$Mor(\gamma\Gamma S,S) = Mor(\Gamma S,\Gamma S),$$

let

$$\alpha_{S} \iff id_{\Gamma S}$$

thus $\alpha_S:\gamma\Gamma S \to S$ and $\alpha:\gamma \circ \Gamma \to id_{\underbrace{CStr}}$ is a monoidal natural isomorphism.]

3.3 <u>REMARK</u> Let <u>C</u>, <u>C'</u> be monoidal categories — then each monoidal functor $F:C \to C'$ induces a strict monoidal functor $F:C \to C'$ induces a strict monoidal functor $F:C \to C'$ and there is a commutative diagram

$$\begin{array}{ccc}
\underline{C} & \xrightarrow{F} & \underline{C'} \\
 & \downarrow & \downarrow & \downarrow \\
 & \downarrow & \downarrow & \downarrow \\
 & \underline{C}_{\underline{\mathtt{str}}} & \xrightarrow{F}_{\underline{\mathtt{str}}} & \underline{C'}_{\underline{\mathtt{str}}}.
\end{array}$$

Here, on an object S,

$$F_{str}S = (FX_1, \dots, FX_n),$$

while on a morphism $u:S \rightarrow T$,

$$(FX_1,...,FX_n) \xrightarrow{F_{\underline{str}}^u} (FY_1,...,FY_m)$$

is that element of Mor(FFS,FFT) defined by requiring commutativity of the square

$$\begin{array}{ccc} \Gamma FS & \longrightarrow & \Gamma FT \\ \approx & & \downarrow & \approx \\ F\Gamma S & \longrightarrow & F\Gamma T, \end{array}$$

where $f \in Mor(\Gamma S, \Gamma T)$ corresponds to u.

[Note: Composition of monoidal functors is preserved by this construction.]

There are five ingredients figuring in the definition of a monoidal category: $\mathbf{\hat{Q}}$, e, R, L, A. Keeping track of R, L, A in calculations can be annoying and one way out is to pass from $\mathbf{\hat{C}}$ to $\mathbf{\hat{C}}_{\underline{\mathbf{str}}}$. But this too has its downside since $\mathbf{\hat{C}}_{\underline{\mathbf{str}}}$ is a more complicated entity than $\mathbf{\hat{C}}$. So, in what follows, we shall stick with $\mathbf{\hat{C}}$ and determine to what extent R, L, A can be eliminated from consideration (i.e., are identities).

Suppose that

are monoidal structures on \underline{C} — then these structures are deemed <u>isomorphic</u> if \exists a monoidal equivalence of the form $(id_{\underline{C}}, \xi, \Xi)$ between them.

N.B. Therefore $\xi:e' \rightarrow e$ is an isomorphism and the

$$\Xi_{X,Y}:X \otimes' Y \rightarrow X \otimes Y$$

are isomorphisms, subject to the coherence conditions of §2.

- 3.4 <u>REMARK</u> The philosophy is that replacing a given monoidal structure on <u>C</u> by another isomorphic to it is of no consequence for the underlying mathematics.
- 3.5 <u>LEMMA</u> Let (Q, e, R, L, A) be a monoidal structure on C. Suppose given a map $Q': Ob C \times Ob C \to Ob C$, an object $e' \in Ob C$, an isomorphism $\xi': e \to e'$, and

isomorphisms

$$\Xi_{X,Y}^{\bullet}: X \otimes Y \to X \otimes^{\bullet} Y.$$

Then there is a unique monoidal structure (Q', e', R', L', A') on C such that

$$(\mathrm{id}_{\underline{C}},\ \xi',\ \Xi'):(\underline{C},\ \boxtimes',\ e',\ R',\ L',\ A')\ \rightarrow\ (\underline{C},\ \boxtimes,\ e,\ R,\ L,\ A)$$

is an isomorphism.

PROOF Extend Q' to a functor $Q':C \times C \to C$ by the prescription

so $\Omega \simeq \Omega'$ (via $E' \in Nat(\Omega,\Omega')$). This done, define R', L', A' by the diagrams

3.6 <u>THEOREM</u> Let $(\mathfrak{Q}, e, R, L, A)$ be a monoidal structure on \mathfrak{C} . Suppose that e' is an object isomorphic to e, say $\xi:e' \to e$ — then there is an isomorphic monoidal structure $(\mathfrak{Q}', e', R', L', A')$ on \mathfrak{C} in which R', L' are identities.

PROOF Bearing in mind 3.5, put

$$X \otimes' Y = X \otimes Y \text{ if } X \neq e' \neq Y$$

and

$$X \otimes^{1} Y = \begin{bmatrix} - & \text{Y if } X = e^{1} \\ & & \\ & & \\ & & \text{X if } Y = e^{1}. \end{bmatrix}$$

Define

$$\Xi_{X,Y}^{\prime}:X \boxtimes Y \rightarrow X \boxtimes^{\prime} Y$$

by stipulating that $\Xi_{X,Y}^{1}$ is to be the identity if $X \neq e^{*} \neq Y$, otherwise let

$$\Xi' = R_X \circ (id_X \otimes \xi)$$

$$\Xi' = L_Y \circ (\xi \otimes id_Y).$$

To establish consistency, i.e., that

$$R \circ (id \otimes \xi) = L \circ (\xi \otimes id),$$
 $e' e' e'$

set $\xi' = \xi^{-1}$ -- then

is an isomorphism and due to the naturality of R, L, the diagrams

commute. Therefore

from which the contention. Finally, by construction (cf. 3.5), R', L' are identities. E.g.:

$$R_X' \circ E' \circ id_X \otimes \xi' = R_X$$

or still,

$$R_X' \circ R_X \circ (id_X \otimes \xi) \circ id_X \otimes \xi' = R_X$$

or still,

$$R_X^{\prime} \circ R_X = R_X = R_X^{\prime} = id_{X^{\bullet}}$$

[Note: If A is the identity and e' is not in the image of Q, then A' is

also the identity. Proof:

$$e' \in \{X,Y,Z\} \Rightarrow A'_{XYZ} = id$$

$$e' \notin \{X,Y,Z\} \& e' \notin Im \& \Rightarrow A'_{XYZ} = A'_{XYZ}.]$$

3.7 REMARK Take e' = e — then the preceding result implies that by passing to an isomorphic monoidal structure, it is always possible to arrange that $\forall \ X \in Ob \ C$,

$$X \otimes e = X = e \otimes X$$
.

The situation for the associativity constraint is more complicated and it will be necessary to impose some conditions on C.

Definition: A construct is a pair (C,U), where

$$U:\underline{C} \to SET$$

is a faithful functor.

3.8 EXAMPLE Define a functor $Q: \underline{SET}^{OP} \to \underline{SET}$ as follows: On objects, $QX = 2^X$ and on morphisms, $Q(A \to B): QA \to QB$ sends $X \subset A$ to the inverse image $f^{-1}(X) \subset B$. In this connection, recall that

$$\begin{array}{c} \texttt{f} \\ \texttt{A} \rightarrow \texttt{B} \in \texttt{Mor SET}^{\texttt{OP}} \end{array}$$

means that

f
$$B \rightarrow A \in Mor \underline{SET}$$
.

Therefore (\underline{SET}^{OP},Q) is a construct.

Let (\underline{C},U) be a construct -- then (\underline{C},U) is <u>amnestic</u> if a \underline{C} -isomorphism f is a \underline{C} -identity whenever Uf is a \underline{SET} -identity, i.e., if $X,Y \in Ob \ \underline{C}$, if $f:X \to Y$ is an isomorphism, if Uf = id, then X = Y and f = id.

Let (\underline{C},U) be a construct -- then (\underline{C},U) is <u>transportable</u> if \forall \underline{C} -object X and every bijection $UX \to S$, \exists a \underline{C} -object Y with UY = S and an isomorphism $\Phi: X \to Y$ such that $U\Phi = \Phi$.

3.9 <u>LEMMA</u> If (\underline{C},U) is amnestic and transportable, then the pair (Y,Φ) is unique.

PROOF Say we have

$$Y_1 \xrightarrow{\Phi_1^{-1}} X \xrightarrow{\Phi_2} Y_2.$$

Then $\Phi_2 \circ \Phi_1^{-1}$ is an isomorphism and

$$U(\Phi_2 \circ \Phi_1^{-1}) = U\Phi_2 \circ U\Phi_1^{-1} = \phi \circ \phi^{-1} = id.$$

Therefore by amnesticity, $Y_1 = Y_2$ and $\Phi_2 \circ \Phi_1^{-1} = id \Rightarrow \Phi_2 = \Phi_1$.

- 3.10 EXAMPLE The construct $\underline{FDVEC_k}$ is amnestic and transportable but the full subcategory of $\underline{FDVEC_k}$ whose objects are the \underline{k}^n , while amnestic, is not transportable.
- 3.11 <u>LEMMA</u> If $\zeta:\underline{SET} \to \underline{SET}$ is an isomorphism and if (\underline{C},U) is amnestic and transportable, then $(\underline{C},\zeta \circ U)$ is amnestic and transportable.

3.12 <u>THEOREM</u> Suppose that (\underline{C},U) is amnestic and transportable. Let $(\underline{Q}, e, R, L, A)$ be a monoidal structure on \underline{C} — then there is an isomorphic strict monoidal structure $(\underline{Q}', e, R', L', A')$ on \underline{C} .

The proof is lengthy, the point of departure being 3.2:

$$\Gamma: \underline{C}_{\underline{str}} \to \underline{C}$$

$$\gamma: \underline{C} \to \underline{C}_{\underline{str}},$$

where

$$\begin{array}{c|c}
 & \Gamma \circ \gamma = id_{\underline{C}} \\
 & \gamma \circ \Gamma \approx id_{\underline{C}}
\end{array}$$

 $\underline{\text{Step 1:}} \ \ \text{Given S} \in \text{Ob} \ \underline{\textbf{C}}_{\text{str}}\text{, consider}$

$$\{S\} \times UTS \in Ob SET.$$

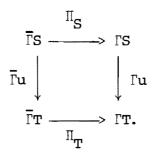
Then the projection

$$\{s\} \times urs \xrightarrow{\pi_S} urs$$

is bijective, so there exists a unique [S] \in Ob \underline{C} with U[S] = {S} \times UTS and a unique isomorphism Π_S :[S] \to TS such that U Π_S = π_S .

Step 2: There is a functor $\Gamma:\underline{C}_{\underline{str}}\to\underline{C}$ which on objects is the prescription $\Gamma:\underline{S}=[S]$

and on morphisms is dictated by requiring that $\Pi \in \text{Nat}(\overline{\Gamma}, \Gamma)$:



Step 3: $\Gamma:\underline{C}_{\underline{\mathtt{str}}}\to\underline{C}$ is an equivalence of categories $(\Pi:\overline{\Gamma}\to\Gamma)$ being a natural isomorphism). In addition, Γ is injective on objects.

Step 4: Define a functor $\overline{\gamma}:\underline{C}\to\underline{C}_{\underline{\operatorname{str}}}$ on objects by taking $\overline{\gamma}X=\gamma X$ if X is not in the image of $\overline{\Gamma}$ and letting $\overline{\gamma}[S]=S$ otherwise. Next, define

$$v_X: \overline{\Gamma}\overline{\gamma}X \to X$$

by

$$[\gamma X] \xrightarrow{\Pi} [\gamma X] \Rightarrow \Gamma \gamma X = X$$

if X is not in the image of $\overline{\Gamma}$ and let $\nu_X=\mathrm{id}_X$ if X = [S] for some S. Since $\overline{\Gamma}$ is fully faithful, we can then define $\overline{\gamma}$ on morphisms by requiring that $\nu\colon\overline{\Gamma}\circ\overline{\gamma}\to\mathrm{id}_{\underline{C}}$ be a natural isomorphism.

Step 5: The arrow

$$\mu = id:id_{\underbrace{C_{str}}} \rightarrow \overline{\gamma} \circ \overline{\Gamma}$$

is a natural isomorphism.

Step 6: The data $(\overline{\Gamma}, \overline{\gamma}, \mu, \nu)$ is an adjoint situation:

$$(\nabla \overline{\Gamma}) \circ (\overline{\Gamma}\mu) = id$$

$$\overline{\Gamma}$$

$$(\nabla \overline{\gamma}\nu) \circ (\mu \overline{\gamma}) = id$$

$$\overline{\gamma}$$

Explicated:

$$\begin{array}{cccc}
 & \nabla_{\overline{\Gamma}S} & \circ \overline{\Gamma}\mu_{S} & = & \mathrm{id} \\
 & \overline{\Gamma}S & & \overline{\Gamma}S & & \\
 & \overline{\gamma}\nu_{X} & \circ & \mu_{X} & = & \mathrm{id} \\
 & \overline{\gamma}X & & \overline{\gamma}X & & \overline{\gamma}X & & \\
\end{array}$$

Claim:

But

$$\overline{\Gamma}$$
S = [S] => $\nu_{\overline{\Gamma}}$ S = $id_{\overline{\Gamma}}$ S ($\overline{\Gamma}\mu_{S}$).

As for the relation

$$\overline{\gamma}_{V_{X}} = id_{\overline{\gamma}X} (\equiv \mu_{\overline{\gamma}X}),$$

since $\overline{\Gamma}$ is faithful, it suffices to show that

$$\overline{\Gamma}\overline{\gamma}\nu_{X} = id_{\overline{\Gamma}\overline{\gamma}X}$$

for all $X \in Ob$ C. But from the definitions, \forall $f \in Mor(\overline{\Gamma\gamma}X,X)$, there is a commutative diagram

Now take $f = v_X$ to get

$$v_X \circ \overline{\Gamma} \overline{\gamma} v_X = v_X \circ v_{\overline{\Gamma} \overline{\gamma} X}$$

or still,

$$\overline{\Gamma}\overline{\gamma}v_{X} = v_{\overline{\Gamma}\overline{\gamma}X}$$

or still,

$$\overline{\Gamma}_{Y}^{-}v_{X} = id_{\overline{\Gamma}_{Y}^{-}X}$$

as desired.

Step 7: The adjoint situation $(\overline{\Gamma}, \overline{\gamma}, \mu, \nu)$ is an adjoint equivalence of categories (μ and ν are natural isomorphisms).

Step 8: Put

$$X \boxtimes' Y = \overline{\Gamma}(\overline{\gamma}X * \overline{\gamma}Y)$$

and let $e' = \overline{\Gamma}\emptyset$ -- then

$$\overline{\gamma}(X \otimes^{\mathbf{1}} Y) = \overline{\gamma}\overline{\Gamma}(\overline{\gamma}X \star \overline{\gamma}Y)$$

$$= \overline{\gamma}X \star \overline{\gamma}Y$$

and

$$\overline{\gamma}e' = \overline{\gamma}\overline{\Gamma}\emptyset = \emptyset$$
.

Step 9: We have

$$X \boxtimes^{\mathbf{I}} (Y \boxtimes^{\mathbf{I}} Z) = \overline{\Gamma}(\overline{\gamma}X * \overline{\gamma}(Y \boxtimes^{\mathbf{I}} Z))$$

$$= \overline{\Gamma}(\overline{\gamma}X * \overline{\gamma}Y * \overline{\gamma}Z)$$

$$= \overline{\Gamma}(\overline{\gamma}(X \boxtimes^{\mathbf{I}} Y) * \overline{\gamma}Z)$$

$$= (X \boxtimes^{\mathbf{I}} Y) \boxtimes^{\mathbf{I}} Z,$$

so A' = id will work.

Step 10: Let

$$\begin{array}{ccc}
R_{X}^{\prime} &= \nu_{X} \\
\vdots & \overline{\Gamma \gamma} X \rightarrow X. \\
L_{X}^{\prime} &= \nu_{X}
\end{array}$$

Then this makes sense:

Furthermore, the diagram

$$X \otimes' e' \otimes' Y \xrightarrow{A' = id} X \otimes' e' \otimes' Y$$

id $\otimes' L' \downarrow \qquad \qquad \downarrow R' \otimes' id$
 $X \otimes' Y \xrightarrow{} X \otimes' Y$

commutes. To see this, note first that

$$X \otimes ' e' \otimes ' Y = \overline{\Gamma}(\overline{\gamma}X * \overline{\gamma}e' * \overline{\gamma}Y)$$

$$= \overline{\Gamma}(\overline{\gamma}X * \emptyset * \overline{\gamma}Y)$$

$$= \overline{\Gamma}(\overline{\gamma}X * \overline{\gamma}Y)$$

$$= X \otimes ' Y.$$

And the arrows

$$\begin{bmatrix} - & R'_X & \mathbf{0'} & \mathrm{id}_{Y} : X & \mathbf{0'} & \mathrm{e'} & \mathbf{0'} & Y \to X & \mathbf{0'} & Y \\ & \mathrm{id}_{X} & \mathbf{0'} & L'_{Y} : X & \mathbf{0'} & \mathrm{e'} & \mathbf{0'} & Y \to X & \mathbf{0'} & Y \end{bmatrix}$$

are identities. E.g.:

$$R_{X}' \boxtimes' id_{Y} = \overline{\Gamma}(\overline{\gamma}v_{X} * \overline{\gamma} id_{Y})$$

$$= \overline{\Gamma}(id_{\overline{\gamma}X} * id_{\overline{\gamma}Y})$$

$$= \overline{\Gamma}(id_{\overline{\gamma}X} * \overline{\gamma}Y)$$

$$= \overline{\Gamma}(id_{\overline{\gamma}X} * \overline{\gamma}Y)$$

$$= id_{\overline{\Gamma}(\overline{\gamma}X} * \overline{\gamma}Y)$$

$$= id_{\overline{X}} ...$$

$$X \boxtimes' Y$$

Step 11: It is clear that

$$\overline{\gamma}$$
: $(\underline{C},\underline{\omega}',e',R',L',A') \rightarrow (\underline{C}_{str},*,\emptyset,R,L,A)$

is a monoidal equivalence (cf. 2.8), thus the same is true of

$$\Gamma \overline{\gamma}: (\underline{C}, \underline{\otimes}', e', R', L', A') \rightarrow (\underline{C}, \underline{\otimes}, e, R, L, A)$$
 (cf. 2.3).

But there is a monoidal natural isomorphism $\Gamma \overline{\gamma} \approx id_{\underline{C}} \colon \ \forall \ X \in Ob \ \underline{C}$,

$$\Gamma_{\overline{Y}X} \xrightarrow{\overline{\Gamma_{\overline{Y}X}}} \overline{\Gamma_{\overline{Y}X}} \xrightarrow{\nu_X} X.$$

Therefore the monoidal structure (@',e',R',L',A') is isomorphic to (@,e,R,L,A).

Step 12: To complete the proof, it is necessary to fine tune (@',e',R',L',A') by an application of 3.6:

$$(\Omega', e', R', L', A') \rightarrow (\Omega'', e'', R'', L'', A'')$$

choosing e'' = e (cf. 1.10). So, R'', L'' are identities. However, by construction, A' is the identity, thus if e is not in the image of Ω' , then A'' is also the identity. To ensure that e is not in the image of Ω' , it is enough that e is not

in the image of $\bar{\Gamma}$. Suppose it were -- then

$$\text{Ue} = \{\text{S}\} \times \text{UTS} \quad (\exists \text{S} \in \text{Ob} \ \underline{\text{C}}_{\underline{\text{str}}}).$$

Now use 3.11 and replace U by ζU , where ζ has the property that ζUe is not a cartesian product of two sets.

3.13 EXAMPLE Consider the construct $\underline{FDVEC_k}$ — then the failure of the tensor product to be associative "on the nose" is an artifact of its definition by a universal property which determines it only up to isomorphism. While the usual procedures do not lead to an associative tensor product, the lesson to be drawn from 3.12 is that it is possible to find a tensor product on $\underline{FDVEC_k}$ such that

$$\begin{bmatrix} & X & \mathbf{\underline{\omega}}_{\underline{k}} & \underline{k} = X \\ & \underline{k} & \mathbf{\underline{\omega}}_{\underline{k}} & X = X \end{bmatrix}$$

and

$$(x \ \underline{\otimes}_{\underline{k}} \ Y) \ \underline{\otimes}_{\underline{k}} \ z = x \ \underline{\otimes}_{\underline{k}} \ (Y \ \underline{\otimes}_{\underline{k}} \ z) = x \ \underline{\otimes}_{\underline{k}} \ Y \ \underline{\otimes}_{\underline{k}} \ z.$$

§4. SYMMETRY

A symmetry for a monoidal category C is a natural isomorphism T, where

$$T_{X,Y}: X \otimes Y \rightarrow Y \otimes X$$

such that

$$T_{Y,X} \circ T_{X,Y} : X \boxtimes Y \to X \boxtimes Y$$

is the identity, $R_X = L_X \circ \tau_{X,e}$, and the diagram

commutes. A symmetric monoidal category is a monoidal category \underline{C} endowed with a symmetry τ . A monoidal category can have more than one symmetry (or none at all).

[Note: The "coherency" principle then asserts that "all" diagrams built up from instances of R, L, A, τ (or their inverses), and id by repeated application of Ω necessarily commute.]

N.B. Let

$$f:C \times C \rightarrow C \times C$$

be the interchange — then f is an isomorphism and $\tau: \Omega \to \Omega$ of is a natural isomorphism.

E.g.: $\underline{\text{VEC}}_{K}$ and $\underline{\text{HILB}}$ are symmetric monoidal.

4.1 EXAMPLE Let C*ALG be the category whose objects are the C*-algebras

and whose morphisms are the *-homomorphisms -- then under the minimal tensor product or the maximal tensor product, C*ALG is a symmetric monoidal category.

4.2 EXAMPLE Let CHX be the category of chain complexes of abelian groups and chain maps — then CHX is monoidal: Take X $\underline{\omega}$ Y to be the tensor product and let $e = \{e_n\}$ be the chain complex defined by $e_0 = \underline{z}$ and $e_n = 0$ (n \neq 0). Further-

4.3 REMARK In the strict situation, matters reduce to the relations $\tau_{e,X} = \tau_{X,e} = id_X$ and

$$T_{X \boxtimes Y,Z} = (T_{X,Z} \boxtimes id_{Y}) \circ (id_{X} \boxtimes T_{Y,Z}).$$

[Note: Therefore

$$T_{X \otimes Y,Z} \circ T_{Y \otimes Z,X} \circ T_{Z \otimes X,Y} = id.$$

4.4 EXAMPLE Let \$ be the permutation category introduced in 1.6 -- then \$ is symmetric monoidal. To establish this, one must exhibit isomorphisms

$$T_{n,m} \in Mor(n \otimes m, m \otimes n)$$

$$= s_{n+m}$$

fulfilling the various conditions. Definition:

with the understanding that $T_{n,0} = id_n = T_{0,n}$, thus

$$T_{m,n} \circ T_{n,m} = id_{n \otimes m}$$

As for the remaining details, it is simplest to work with permutation matrices, so take n > 0, m > 0, and note that

$$T_{n,m} = \begin{bmatrix} - & 0 & I_{m} & - \\ & & & \\ & I_{n} & 0 \end{bmatrix}$$

Then

$$(\mathsf{T}_{\mathsf{n},\mathsf{p}} \boxtimes \mathsf{id}_{\mathsf{m}}) \circ (\mathsf{id}_{\mathsf{n}} \boxtimes \mathsf{T}_{\mathsf{m},\mathsf{p}})$$

[Note:

$$\begin{bmatrix} 0 & \mathbf{I}_{\mathbf{m}} \\ \mathbf{I}_{\mathbf{n}} & 0 \end{bmatrix} \begin{bmatrix} \sigma & 0 \\ 0 & \tau \end{bmatrix} = \begin{bmatrix} 0 & \tau \\ \sigma & 0 \end{bmatrix}$$

$$= \begin{bmatrix} \tau & 0 \\ 0 & \sigma \end{bmatrix} \begin{bmatrix} 0 & \mathbf{I}_{\mathbf{m}} \\ \mathbf{I}_{\mathbf{n}} & 0 \end{bmatrix}.$$

Therefore naturality is manifest, i.e.,

$$\tau_{n,m} \circ (\sigma \otimes \tau) = (\tau \otimes \sigma) \circ \tau_{n,m}$$
.

Let \underline{C} , \underline{C}' be symmetric monoidal categories — then a <u>symmetric monoidal</u> functor is a monoidal functor (F, ξ, Ξ) such that the diagram

$$FX \otimes' FY \xrightarrow{\Xi_{X,Y}} F(X \otimes Y)$$

$$\uparrow^!_{FX,FY} \downarrow \qquad \qquad \downarrow^{F^{\top}_{X,Y}}$$

$$FY \otimes' FX \xrightarrow{\Xi_{Y,X}} F(Y \otimes X)$$

commutes.

N.B. The monoidal natural transformations between symmetric monoidal functors are, by definition, "symmetric monoidal" (i.e., no further conditions are imposed

that reflect the presence of a symmetry).

[Note: Therefore the subcategory $[\underline{C},\underline{C}']^{\underline{\omega},T}$ of $[\underline{C},\underline{C}']^{\underline{\omega}}$ whose objects are the symmetric monoidal natural transformations is, by definition, a full subcategory.]

4.5 EXAMPLE Recall that \mathfrak{S}_n has the following presentation: It is generated by $\sigma_1,\ldots,\sigma_{n-1}$ subject to the relations

$$\sigma_{\mathbf{i}}^2 = 1$$
, $\sigma_{\mathbf{i}}\sigma_{\mathbf{i}+1}\sigma_{\mathbf{i}} = \sigma_{\mathbf{i}+1}\sigma_{\mathbf{i}}\sigma_{\mathbf{i}+1}$, $\sigma_{\mathbf{i}}\sigma_{\mathbf{j}} = \sigma_{\mathbf{j}}\sigma_{\mathbf{i}}$ ($|\mathbf{i}-\mathbf{j}| > 1$).

Suppose now that \underline{C} is symmetric strict monoidal and fix $X \in Ob$ \underline{C} . Define automorphisms $\mathbb{I}^1, \ldots, \mathbb{I}^{n-1}$ of $X^{\boxtimes n}$ by

$$\Pi^{i} = id_{X^{\Omega(i-1)}} \otimes T_{X,X} \otimes id_{X^{\Omega(n-i-1)}}.$$

Then there exists a unique homomorphism

$$\Pi_n^X: \mathcal{S}_n \to \text{Aut } X^{\Omega n}$$

of groups such that

$$\Pi_{\mathbf{n}}^{\mathbf{X}}(\sigma_{\mathbf{i}}) = \Pi^{\mathbf{i}} \quad (\mathbf{i} = 1, \dots, \mathbf{n-1}).$$

Combining the \prod_n^X then leads to a symmetric monoidal functor $F:\mathcal{G}\to \underline{C}$ such that $Fn=X^{\underline{\Omega}n}$.

4.6 <u>LEMMA</u> Let $F: C \to C'$ be a monoidal equivalence. Assume: C is symmetric—then the symmetry τ on C can be transferred to a symmetry τ' on C' in such a way as to render F symmetric monoidal.

[Define
$$\tau_{FX,FY}$$
 by

and recall that F has a representative image (cf. 2.5).]

- 4.7 EXAMPLE If C is symmetric monoidal, then $\underline{C}_{\underline{\mathtt{str}}}$ is symmetric monoidal and $\gamma:\underline{C}\to\underline{C}_{\underline{\mathtt{str}}}$ is a symmetric monoidal equivalence.
- 4.8 <u>IEMMA</u> Let \underline{C} , \underline{C} ' be symmetric monoidal and let (F,F',μ,μ') be an adjoint equivalence. Assume: F is symmetric monoidal then F' is symmetric monoidal (cf. 2.9).

§5. DUALITY

Let \underline{C} be a monoidal category — then each $X \in Ob \ \underline{C}$ defines functors

Definition: C is

left closed

if $\forall X \in Ob C$,

— ② X admits a right adjoint, denoted lhom(X,—)

X ② — admits a right adjoint, denoted rhom (X,—).

[Note: C is closed if it is both left closed and right closed.]
So:

_ C left closed => Mor(Y ⊗ X,Z) ≈ Mor(Y,lhom(X,Z)) C right closed => Mor(X ⊗ Y,Z) ≈ Mor(Y,rhom(X,Z))

for all $Y,Z \in Ob C$.

N.B. The functor

1hom(X,---)

is called the internal hom functor attached to X. right

- 5.1 REMARK If \underline{C} is symmetric monoidal, then left and right internal homs are naturally isomorphic and if \underline{C} is left or right closed, then \underline{C} is closed.
- 5.2 EXAMPLE Given a commutative ring \underline{k} , let $\underline{MOD}_{\underline{k}}$ be the category whose objects are the left \underline{k} -modules and whose morphisms are the \underline{k} -linear maps then $\underline{MOD}_{\underline{k}}$ is symmetric monoidal. Moreover, $\underline{MOD}_{\underline{k}}$ is closed and

Thom(X,Z)
$$\approx$$
 Hom _{\underline{k}} (X,Z)
rhom(X,Z) \approx Hom _{\underline{k}} (X,Z).

- 5.3 <u>LFMMA</u> Suppose that <u>C</u> is left closed then \forall X \in Ob <u>C</u>, the functor \otimes X preserves colimits (being a left adjoint) and the functor lhom(X,—) preserves limits (being a right adjoint).
- 5.4 <u>LEMMA</u> Suppose that <u>C</u> is left closed then \forall Z \in Ob <u>C</u>, the cofunctor lhom(—,Z) converts colimits to limits.

<u>PROOF</u> Let \underline{I} be a small category, $\Delta:\underline{I}\to \underline{C}$ a diagram for which colim \underline{I} $\Delta_{\underline{I}}$ exists — then \forall $Y\in Ob$ \underline{C} ,

$$\begin{aligned} & \text{Mor} (Y, \text{lhom}(\text{colim}_{\underline{I}} \ \Delta_{\underline{i}}, Z)) \\ & \approx & \text{Mor} (Y \ \Omega \ \text{colim}_{\underline{I}} \ \Delta_{\underline{i}}, Z) \\ & \approx & \text{Mor}(\text{colim}_{\underline{I}} \ (Y \ \Omega \ \Delta_{\underline{i}}), Z) \end{aligned}$$

$$\approx \lim_{\underline{I}} \operatorname{Mor}(Y \otimes \Delta_{\underline{i}}, Z)$$

$$\approx \lim_{\underline{I}} \operatorname{Mor}(Y, \operatorname{lhom}(\Delta_{\underline{i}}, Z))$$

$$\approx \operatorname{Mor}(Y, \lim_{\underline{I}} \operatorname{lhom}(\Delta_{\underline{i}}, Z))$$

$$\Longrightarrow$$

$$\operatorname{lhom}(\operatorname{colim}_{\underline{I}} \Delta_{\underline{i}}, Z) \approx \lim_{\underline{I}} \operatorname{lhom}(\Delta_{\underline{i}}, Z).$$

Let \underline{C} be a monoidal category. Given $X\in Ob\ \underline{C}$, an object ${}^VX\in Ob\ \underline{C}$ is said to be a <u>left dual</u> of X if \exists morphisms

$$\begin{bmatrix} - & \varepsilon_{X} : ^{V} X \boxtimes X \rightarrow e \\ & \eta_{X} : e \rightarrow X \boxtimes ^{V} X \end{bmatrix}$$

and commutative diagrams

N.B. When C is strict, these diagrams reduce to the relations

$$(id_X \otimes \varepsilon_X) \circ (\eta_X \otimes id_X) = id_X$$

$$(\varepsilon_X \otimes id_Y) \circ (id_Y \otimes \eta_X) = id_Y.$$

5.5 <u>LEMMA</u> Suppose that $^{\vee}X$ is a left dual of X — then the functor — \mathbf{Q} $^{\vee}X$ is a right adjoint for the functor — \mathbf{Q} X and the functor $^{\vee}X$ \mathbf{Q} — is a left adjoint for the functor X \mathbf{Q} —.

In brief: $\forall Y,Z \in Ob C$,

Mor(Y
$$\mathbf{Q}$$
 X,Z) \approx Mor(Y,Z \mathbf{Q} $^{\vee}$ X)

Mor($^{\vee}$ X \mathbf{Q} Y,Z) \approx Mor(Y,X \mathbf{Q} Z).

<u>PROOF</u> It will be enough to show that — \mathbf{Q} $^{\mathsf{V}}X$ is a right adjoint for — \mathbf{Q} X, the proof that $^{\mathsf{V}}X$ \mathbf{Q} — is a left adjoint for X \mathbf{Q} — being similar. So let

and to simplify the writing, take $\underline{\mathbf{C}}$ strict. Define

$$\mu \in \text{Nat}(\text{id}_{\underline{C}}, G \circ F)$$

$$\nu \in \text{Nat}(F \circ G, \text{id}_{\underline{C}})$$

by

$$\mu_{W} \in Mor(W, W \otimes X \otimes ^{\vee}X)$$

$$\mu_{W} = id_{W} \otimes \eta_{X}$$

Consider

$$(\nabla F) \circ (F\mu)$$
.

Thus

$$((\vee F) \circ (F\mu))_W = (\vee F)_W \circ (F\mu)_W$$
.

And

$$(\mu) \begin{array}{|c|c|c|c|}\hline & F\mu \in Nat(F,FGF)\\ & & \\ & (F\mu)_W:FW \to FGFW \end{array}$$

or still,

(v)
$$\bigvee_{(vF)_{W}: FGFW \rightarrow FW}$$

or still,

$$(\vee F)_{W} : W \boxtimes X \boxtimes {}^{V}X \boxtimes X \xrightarrow{id_{W} \boxtimes id_{X} \boxtimes \epsilon_{X}} W \boxtimes X.$$

Therefore

$$\left(\mathsf{vF} \right)_{W} \, \circ \, \left(\mathsf{F} \boldsymbol{\mu} \right)_{W} \in \, \mathsf{Mor} (\mathsf{W} \, \boldsymbol{\boxtimes} \, \mathsf{X}, \! \mathsf{W} \, \boldsymbol{\boxtimes} \, \mathsf{X})$$

is the composition

$$(\mathrm{id}_{W} \otimes \mathrm{id}_{X} \otimes \varepsilon_{X}) \circ (\mathrm{id}_{W} \otimes n_{X} \otimes \mathrm{id}_{X})$$

$$= (\mathrm{id}_{W} \circ \mathrm{id}_{W}) \otimes ((\mathrm{id}_{X} \otimes \varepsilon_{X}) \circ (n_{X} \otimes \mathrm{id}_{X}))$$

$$= \mathrm{id}_{W} \otimes \mathrm{id}_{X}$$

$$= \mathrm{id}_{W} \otimes X$$

$$= \mathrm{id}_{FW}$$

$$= (\mathrm{id}_{F})_{W}.$$

I.e.:

$$(\lor F) \circ (F\mu) = id_{F}.$$

The verification that

(Gv)
$$\circ$$
 (μ G) = id_{G}

is analogous.

5.6 <u>LFMMA</u> A left dual of X, if it exists, is unique up to isomorphism.

PROOF Suppose that

are two left duals of X -- then the functors

are naturally isomorphic (both being right adjoints for — $\mbox{0}$ X), so \forall W \in Ob \mbox{C} ,

$$W \otimes {}^{\vee}X_1 \approx W \otimes {}^{\vee}X_2.$$

Now specialize and take W = e to get

$$e \otimes ^{\vee} x_1 \approx e \otimes ^{\vee} x_2$$

=>

$$^{\mathsf{v}}\mathbf{x}_{1} \approx {}^{\mathsf{v}}\mathbf{x}_{2}$$
.

[Note: Explicated,

5.7 <u>REMARK</u> Suppose that $({}^{V}X, \epsilon_{X}, \eta_{X})$ is a left dual of X. Let $\phi: {}^{V}X \rightarrow {}^{V}X'$ be an isomorphism and put

Then the triple $({}^{\vee}X^{\bullet}, \epsilon_{X}^{\bullet}, \eta_{X}^{\bullet})$ is a left dual of X.

[Consider first the case when \underline{C} is strict, thus, e.g.,

$$\begin{split} &(\mathrm{id}_X \ \underline{\otimes} \ \epsilon_X') \ \circ \ (\eta_X' \ \underline{\otimes} \ \mathrm{id}_X) \\ &= \ \mathrm{id}_X \ \underline{\otimes} \ (\epsilon_X \ \circ \ (\varphi^{-1} \ \underline{\otimes} \ \mathrm{id}_X)) \ \circ \ ((\mathrm{id}_X \ \underline{\otimes} \ \varphi) \ \circ \ \eta_X) \ \underline{\otimes} \ \mathrm{id}_X \\ &= \ \mathrm{id}_X \ \underline{\otimes} \ \epsilon_X \ \circ \ \mathrm{id}_X \ \underline{\otimes} \ (\varphi^{-1} \ \underline{\otimes} \ \mathrm{id}_X) \ \circ \ (\mathrm{id}_X \ \underline{\otimes} \ \varphi) \ \underline{\otimes} \ \mathrm{id}_X \ \circ \ \eta_X \ \underline{\otimes} \ \mathrm{id}_X. \end{split}$$

But

$$(id_X \otimes \phi) \otimes id_X = id_X \otimes (\phi \otimes id_X)$$

=>

$$id_{X} \boxtimes (\phi^{-1} \boxtimes id_{X}) \circ (id_{X} \boxtimes \phi) \boxtimes id_{X}$$

$$= id_{X} \boxtimes (\phi^{-1} \boxtimes id_{X}) \circ id_{X} \boxtimes (\phi \boxtimes id_{X})$$

$$= id_{X} \boxtimes (\phi^{-1} \boxtimes id_{X}) \circ (\phi \boxtimes id_{X})$$

$$= id_{X} \boxtimes id_{Y}$$

$$= id_{X} \boxtimes id_{Y}$$

$$= id_{X} \boxtimes X$$

=>

$$id_X \Omega \epsilon_X^{1} \circ \eta_X^{1} \Omega id_X$$

$$= (\mathrm{id}_X \ \mathfrak{Q} \ \epsilon_X) \ \circ \ (\eta_X \ \mathfrak{Q} \ \mathrm{id}_X) \ = \ \mathrm{id}_X.$$

In general, the claim is that id_{χ} equals

$$R \circ (id_{X} \otimes \epsilon_{X}^{*}) \circ A^{-1} \circ (\eta_{X}^{*} \otimes id_{X}) \circ L^{-1}$$

or still,

$$\text{R} \circ \operatorname{id}_X \, \underline{\text{O}} \, \left(\varepsilon_X \, \circ \, \left(\phi^{-1} \, \underline{\text{O}} \, \operatorname{id}_X \right) \right) \, \circ \, A^{-1} \, \circ \, \left(\left(\operatorname{id}_X \, \underline{\text{O}} \, \phi \right) \, \circ \, \eta_X \right) \, \underline{\text{O}} \, \operatorname{id}_X \, \circ \, L^{-1}$$
 or still,

$$\mathsf{R} \, \circ \, \operatorname{id}_{\mathsf{X}} \, \boldsymbol{\otimes} \, \, \boldsymbol{\varepsilon}_{\mathsf{X}} \, \circ \, \operatorname{id}_{\mathsf{X}} \, \boldsymbol{\otimes} \, \, (\boldsymbol{\varphi}^{-1} \, \, \boldsymbol{\otimes} \, \operatorname{id}_{\mathsf{X}}) \, \circ \, \boldsymbol{\mathsf{A}}^{-1} \, \circ \, (\operatorname{id}_{\mathsf{X}} \, \boldsymbol{\otimes} \, \boldsymbol{\varphi}) \, \, \boldsymbol{\otimes} \, \operatorname{id}_{\mathsf{X}} \, \circ \, \boldsymbol{\eta}_{\mathsf{X}} \, \boldsymbol{\otimes} \, \operatorname{id}_{\mathsf{X}} \, \circ \, \boldsymbol{\mathsf{L}}^{-1}.$$

Here

$$A^{-1}: (X \boxtimes {}^{\vee}X^{*}) \boxtimes X \to X \boxtimes ({}^{\vee}X^{*} \boxtimes X).$$

So, to complete the verification, one has only to show that the composition

$$(X \otimes ^{\vee}X) \otimes X \xrightarrow{\text{(id } \otimes \varphi) \otimes \text{id}} (X \otimes ^{\vee}X') \otimes X$$

$$\xrightarrow{A^{-1}} X \otimes (^{\vee}X' \otimes X)$$

$$\xrightarrow{\text{id } \otimes (\varphi^{-1} \otimes \text{id})} X \otimes (^{\vee}X \otimes X)$$

is

$$(X \otimes ^{\vee}X) \otimes X \longrightarrow A^{-1} \longrightarrow X \otimes (^{\vee}X \otimes X).$$

However, due to the naturality of the associativity constraint, there is a

commutative diagram

$$(X \otimes ^{\vee}X) \otimes X \xrightarrow{A^{-1}} X \otimes (^{\vee}X \otimes X)$$

$$(id \otimes \varphi) \otimes id \qquad \qquad \downarrow id \otimes (\varphi \otimes id)$$

$$(X \otimes ^{\vee}X^{i}) \otimes X \xrightarrow{A^{-1}} X \otimes (^{\vee}X^{i} \otimes X).$$

And

$$(id \otimes (\phi \otimes id))^{-1} = id \otimes (\phi^{-1} \otimes id).$$

A monoidal category \underline{C} is said to be <u>left autonomous</u> if each object in \underline{C} admits a left dual.

N.B. Suppose that C is left autonomous. Given $f \in Mor(X,Y)$, define $^{V}f \in Mor(^{V}Y, ^{V}X)$ by

Then the assignment

$$\begin{bmatrix} - & x \rightarrow {}^{\vee}x \\ & & \\ & f \rightarrow {}^{\vee}f \end{bmatrix}$$

defines a cofunctor $C \rightarrow C$.

[Note: The specific form of $^{\vee}f$ depends on the choices of $^{\vee}X$ and $^{\vee}Y$.]

5.8 REMARK If C is left autonomous and if $X,Y \in Ob$ C, then $(X \boxtimes Y)$ is isomorphic to $Y \boxtimes X$.

[We have

$$Mor(^{\vee}(X \otimes Y) \otimes W, Z) \approx Mor(W, (X \otimes Y) \otimes Z)$$

$$\approx Mor(^{\vee}X \otimes W, Y \otimes Z))$$

$$\approx Mor(^{\vee}X \otimes W, Y \otimes Z)$$

$$\approx Mor(^{\vee}Y \otimes (^{\vee}X \otimes W), Z)$$

$$\approx Mor((^{\vee}Y \otimes ^{\vee}X) \otimes W, Z)$$

$$\approx Mor((^{\vee}Y \otimes ^{\vee}X) \otimes W, Z)$$

$$=>$$

5.9 <u>LEMMA</u> Suppose that \underline{C} is left autonomous — then \underline{C} is left closed. <u>PROOF</u> In fact, $\forall \ X \in Ob \ \underline{C}$,

$$1hom(X, ---) = --- \otimes ^{\vee}X.$$

One can also introduce the notion of a right dual $\boldsymbol{X}^{\boldsymbol{V}}$ of \boldsymbol{X} , where this time

subject to the obvious commutativity conditions. Here the functor — $\mathbf{Q} \times^{\mathbf{V}}$ is a left adjoint for the functor — $\mathbf{Q} \times \mathbf{X}$ and the functor $\mathbf{X}^{\mathbf{V}} \times \mathbf{Q}$ — is a right adjoint for the functor $\mathbf{X} \times \mathbf{Q}$ —.

[Note: If X admits a left dual ^{V}X and a right dual X^{V} , then in general ^{V}X and X^{V} are not isomorphic. On the other hand, it is true that

$$(^{\vee}X)^{\vee} \approx X \approx ^{\vee}(X^{\vee}).$$

E.g.:

$$Mor(Y \otimes (^{\vee}X)^{\vee}, Z) \approx Mor(Y, Z \otimes ^{\vee}X) \approx Mor(Y \otimes X, Z)$$

=>

$$(^{\mathsf{V}}\mathsf{X})^{\mathsf{V}} \approx \mathsf{X}.]$$

The definition of "right autonomous" is clear and we shall term \underline{C} autonomous if it is both left and right autonomous.

5.10 <u>LEMMA</u> Suppose that \underline{C} is right autonomous -- then \underline{C} is right closed. PROOF In fact, \forall X \in Ob \underline{C} ,

$$\operatorname{rhom}(X, --) = X^{\vee} \otimes --.$$

5.11 REMARK If \underline{C} is autonomous, then $-\underline{\omega}$ — preserves colimits in both variables.

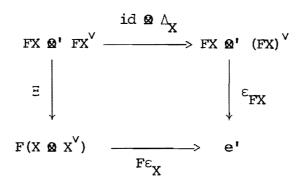
Suppose that $F:\underline{C}\to\underline{C}'$ is a monoidal functor. Assume: X^V is a right dual of X — then FX^V is a right dual of FX. Proof: Consider the arrows

[Note: Assume that $\underline{\mathbf{C}}$, $\underline{\mathbf{C}}'$ are right autonomous — then there is an isomorphism

$$\triangle_{\mathbf{X}}: \mathbf{FX}^{\vee} \rightarrow (\mathbf{FX})^{\vee},$$

namely the composition

and the diagram



commutes.]

N.B. One can, of course, work equally well with left duals.

5.12 LEMMA Let

(F,
$$\xi$$
, Ξ)
(G, θ , Θ)

be monoidal functors and let α :F \rightarrow G be a monoidal natural transformation. Assume: The source C of F and G is autonomous — then α is a monoidal natural isomorphism.

PROOF The claim is that $\forall X \in Ob C$,

$$\alpha_{X}:FX \to GX$$

is an isomorphism. From the above, FX^{\vee} (GX $^{\vee}$) is a right dual of FX (GX) or still, FX (GX) is a left dual of FX^{\vee} (GX $^{\vee}$). This said, form

$$\alpha_{X^{\vee}}:FX^{\vee} \to GX^{\vee}$$

and consider

$$(\alpha_{X^{\vee}}):GX \to FX.$$

[Note: Accordingly, if \underline{C} is autonomous, then the metacategory $[\underline{C},\underline{C}']^{\underline{Q}}$ is a groupoid.]

Suppose that \underline{C} is symmetric monoidal and left autonomous — then \underline{C} is right autonomous, hence \underline{C} is autonomous. Proof: Given $X \in Ob \ \underline{C}$, take $X^V = {}^V X$ and define morphisms

$$\begin{bmatrix} - & \mathbf{X} & \mathbf{X} & \mathbf{X}^{\mathsf{V}} \to \mathbf{e} \\ \\ & \mathbf{e} \to \mathbf{X}^{\mathsf{V}} & \mathbf{\Omega} & \mathbf{X} \end{bmatrix}$$

by

5.13 EXAMPLE FDVECT $_{\underline{k}}$ is autonomous. In fact, FDVECT $_{\underline{k}}$ is symmetric monoidal, so it suffices to set up a left duality. Thus given X, let ^{V}X be its dual and define

$$\varepsilon_{\mathbf{X}}^{\mathsf{Y}} \times \mathbf{X} \otimes \mathbf{X} \to \underline{\mathbf{k}}$$

by

$$\varepsilon_{\mathbf{X}}(\lambda,\mathbf{x}) = \lambda(\mathbf{x})$$
.

On the other hand, there is a canonical isomorphism

$$\phi$$
:Hom(X,X) \rightarrow Hom(k , $^{\vee}$ X \otimes X)

and we let

$$\eta_{\mathbf{X}} = \phi(\mathrm{id}_{\mathbf{X}})$$
.

[Note: An object X in $\underline{\underline{\text{VEC}}}_{\underline{\underline{k}}}$ admits a left dual iff it is finite dimensional.]

5.14 EXAMPLE The full subcategory of $\underline{\text{MOD}}_{\underline{k}}$ whose objects are finitely generated projective is autonomous (cf. 5.2).

Assume still that C is symmetric monoidal and left autonomous.

5.15 LEMMA There is a monoidal natural isomorphism

$$id_{\underline{C}} \rightarrow {}^{VV}(\underline{\hspace{1cm}}).$$

[To see this, consider the composition

N.B. Let

$$\delta_{\mathbf{x}}: \mathbf{X} \to {}^{\mathsf{vv}}\mathbf{X}$$

be the arrow constructed above - then

$$(\delta_{X})^{-1} = {}^{v}(\delta_{X}^{v})$$
 (cf. 5.12).

But here $X^{V} = {}^{V}X$, so

$$(\delta_{\mathbf{X}})^{-1} = (\delta_{\mathbf{Y}}).$$

[Note: To make sense of this, recall that

$$\begin{bmatrix} - & X \text{ is a left dual of } X^{V} \\ & - & X \text{ is a left dual of } V^{V}(X^{V}). \end{bmatrix}$$

And

$$\delta_{X^{\vee}}: X^{\vee} \to {}^{\vee\vee}(X^{\vee})$$

=>

$$(\delta_{X^{\vee}}): {}^{\vee \vee}X \to X.$$

§6. TWISTS

Let \underline{C} be symmetric monoidal and left autonomous — then a $\underline{\text{twist}}$ Ω is a monoidal natural isomorphism of the identity functor $\mathrm{id}_{\underline{C}}$ such that \forall X \in Ob \underline{C} ,

$$(\Omega_X \otimes id_{\vee_X}) \circ \eta_X = (id_X \otimes \Omega_{\vee_X}) \circ \eta_X$$

[Note: Tacitly, id_ is taken to be strict (ξ = id, Ξ = id), thus from the definitions

$$\Omega_{X \otimes Y} = \Omega_{X} \otimes \Omega_{Y} \text{ and } \Omega_{e} = id_{e}.$$

To consolidate the terminology, a symmetric monoidal \underline{C} which is left autonomous and has a twist Ω will be referred to as a ribbon category.

N.B. The choice $\Omega_{X} = id_{X}$ is permissible, in which case C is said to be even.

It was pointed out near the end of §5 that an even ribbon category is right autonomous. This fact is true in general. Proof: Given $X \in Ob \ \underline{C}$, take $X^{V} = {}^{V}X$ and define morphisms

$$\begin{bmatrix} - & \mathbf{X} & \mathbf{Q} & \mathbf{X}^{\mathsf{V}} \to \mathbf{e} \\ \\ \mathbf{e} & \mathbf{A}^{\mathsf{V}} & \mathbf{Q} & \mathbf{X} \end{bmatrix}$$

by

6.1 LEMMA In the presence of a twist Ω ,

$$X \approx (^{\vee}X)$$
.

PROOF Consider the composition

$$\begin{array}{c}
X \xrightarrow{R^{-1}} X \otimes e \\
& \xrightarrow{\Omega_{X} \otimes \eta_{V_{X}}} & X \otimes e
\end{array}$$

$$\begin{array}{c}
X \otimes \eta_{V_{X}} & X \otimes e
\end{array}$$

$$\begin{array}{c}
X \otimes \eta_{V_{X}} & X \otimes e
\end{array}$$

$$\begin{array}{c}
X \otimes \eta_{V_{X}} & X \otimes e
\end{array}$$

$$\begin{array}{c}
X \otimes \chi_{X} \otimes \chi_{X} & X \otimes \chi_{X} \otimes \chi_{X} \otimes \chi_{X}
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$$\begin{array}{c}
X \otimes \chi_{X}$$

E.g.:

$$^{\mathsf{V}}$$
e \approx $^{\mathsf{V}}$ e \otimes e \approx $^{\mathsf{V}}$ e \otimes e \otimes $^{\mathsf{V}}$ ($^{\mathsf{V}}$ e) \approx $^{\mathsf{V}}$ ($^{\mathsf{V}}$ e) \approx e.

6.2 <u>LEMMA</u> In the presence of a twist Ω , the left and right dual of every morphism $f:X \to Y$ agree: ${}^Vf = f^V$.

Let \underline{C} be a ribbon category. Given $f \in Mor(X,X)$, define the \underline{trace} of f by

$$\operatorname{tr}_{X}(f) = \varepsilon_{X} \circ \tau_{X, X} \circ \Omega_{X} \otimes \operatorname{id}_{X} \circ (f \otimes \operatorname{id}_{X}) \circ \eta_{X}.$$

[Note:

$$tr_{X}(f) \in Mor(e,e) (= \underline{M}(\underline{C})).]$$

6.3 LEMMA We have

1.
$$tr_{X}(f) = tr_{V_{X}}(^{V}f);$$

2.
$$tr_X(g \circ f) = tr_Y(f \circ g) (f:X \rightarrow Y, g:Y \rightarrow X);$$

3.
$$\operatorname{tr}_{X_1 \boxtimes X_2} (f_1 \boxtimes f_2) = \operatorname{tr}_{X_1} (f_1) \operatorname{tr}_{X_2} (f_2)$$
.

Put

$$\dim X = \operatorname{tr}_{X}(\operatorname{id}_{X}),$$

the dimension of X.

So, on the basis of 6.3,

$$\dim X = \dim^{\vee} X$$

and

$$\dim (X \otimes Y) = (\dim X) (\dim Y)$$
.

N.B. Take $\Omega = id$ — then the <u>categorical dimension</u> of X is the arrow

$$e \xrightarrow{\eta_X} X \boxtimes {}^{\vee}X \xrightarrow{X, {}^{\vee}X} {}^{\vee}X \boxtimes X \xrightarrow{\varepsilon_X} e.$$

6.4 EXAMPLE Consider FDVECk (viewed as an even ribbon category (cf. 5.13)) — then the trace of $f:X \to X$ is the composition

Therefore the abstract definition of $\mathrm{tr}_{\chi}(f)$ is the usual one. In particular:

$$\dim X = (\dim_{\underline{k}} X) 1_{\underline{k}}$$
.

E.g.:

$$\dim \underline{k}^n = nl_{\underline{k}'}$$

the distinction between $n\in N$ and $nl_{\underline{k}}$ being potentially essential if \underline{k} has non-zero characteristic.

6.5 REMARK While evident, it is important to keep in mind that the definitions of trace and dimension depend on all the underlying assumptions, viz. that our monoidal C is symmetric, left autonomous, and has a twist Ω .

Suppose that \underline{C} , \underline{C}' are ribbon categories with respective twists Ω , Ω' — then a symmetric monoidal functor $F:\underline{C} \to \underline{C}'$ is <u>twist preserving</u> if $\forall \ X \in Ob \ \underline{C}$,

$$\mathbf{F}\Omega_{\mathbf{X}} = \Omega_{\mathbf{F}\mathbf{X}}^{\mathbf{I}}$$
.

6.6 $\underline{\text{LFMMA}}$ If $F:\underline{C} \to \underline{C}'$ is twist preserving, then \forall $f \in Mor(X,X)$, the diagram

commutes.

Matters are invariably simpler if \underline{C} is a strict ribbon category, which will be the underlying supposition in 6.7 - 6.9 below.

6.7 LEMMA The arrows

$$\begin{array}{c|c} & & & \\ & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & & \\ & &$$

are mutually inverse isomorphisms.

PROOF Take X = e in the relation

$$(id_X \otimes \epsilon_X) \circ (n_X \otimes id_X) = id_X$$

to see that

$$\varepsilon_{e} \circ \eta_{e} = id_{e}$$
.

Now fix an isomorphism $\phi:e \to {}^{\mathsf{V}}\!e$ — then

$$\begin{bmatrix} - & \phi^{-1} & \circ & \eta_{\mathbf{e}} \\ & & & \in \underline{\mathbf{M}}(\underline{\mathbf{C}}) \\ & & & & \\ - & & \mathbf{e} & \bullet & \end{bmatrix}$$

=>

$$(\phi^{-1} \circ \eta_e) \circ (\epsilon_e \circ \phi) = (\epsilon_e \circ \phi) \circ (\phi^{-1} \circ \eta_e)$$
 (cf. 1.4)
= $\epsilon_e \circ \eta_e = id_e$

=>

$$\eta_e \circ \varepsilon_e = id_e$$
.

6.8 LEMMA $\forall s \in M(C)$,

$$tr_e(s) = s.$$

PROOF In fact,

$$tr_{e}(s) = \varepsilon_{e} \circ \tau_{e, e} \circ \Omega_{e} \otimes id_{e} \circ (s \otimes id_{e}) \circ \eta_{e}$$

$$= \varepsilon_{e} \circ id_{e} \circ id_{e} \circ (s \otimes id_{e}) \circ \eta_{e}$$

$$= \varepsilon_{e} \circ (s \otimes id_{e}) \circ \eta_{e}$$

$$= (id_{e} \otimes \varepsilon_{e}) (s \otimes id_{e} \otimes id_{e}) (id_{e} \otimes \eta_{e})$$

$$= s \otimes (\varepsilon_{e} \circ \eta_{e})$$

$$= s \otimes id_{e}$$

$$= s.$$

[Note: Therefore

$$\dim e = \operatorname{tr}_{e}(\operatorname{id}_{e}) = \operatorname{id}_{e}.$$

6.9 LEMMA $\forall X \in Ob C$,

$$\Omega_{\mathbf{v}_{\mathbf{X}}} = {}^{\mathbf{v}}\Omega_{\mathbf{X}}.$$

PROOF The compositions

are equal, thus the compositions

are equal. Postcompose with $\epsilon_X \otimes \mathrm{id}_{\chi}$ — then the first line gives ${}^{\mathsf{V}}\Omega_X$, while the second line is

$$\epsilon_{\mathbf{X}} \otimes \mathrm{id}_{\mathbf{V}_{\mathbf{X}}} \circ \mathrm{id}_{\mathbf{V}_{\mathbf{X}}} \otimes \mathrm{id}_{\mathbf{X}} \otimes \Omega_{\mathbf{V}_{\mathbf{X}}} \circ \mathrm{id}_{\mathbf{V}_{\mathbf{X}}} \otimes \eta_{\mathbf{X}}$$

or still,

$$\varepsilon_{X} \otimes id_{\chi_{X}} \circ id_{\chi_{X} \otimes X} \otimes \Omega_{\chi_{X}} \circ id_{\chi_{X}} \otimes \eta_{X}$$

or still,

$$(\mathrm{id}_{\mathrm{e}} \ \mathbf{0} \ \mathbf{0}_{\mathsf{v}_{X}}) \ \cdot \ (\mathbf{e}_{\mathsf{X}} \ \mathbf{0} \ \mathrm{id}_{\mathsf{v}_{\mathsf{X}}}) \ \cdot \ \mathrm{id}_{\mathsf{v}_{\mathsf{X}}} \ \mathbf{0} \ \mathbf{0}_{\mathsf{X}}$$

or still,

$$\Omega_{\mathbf{X}} \circ \mathrm{id}_{\mathbf{X}} = \Omega_{\mathbf{X}}$$

6.10 REMARK Let \underline{C} be a ribbon category — then this structure can be transferred to $\underline{C}_{\mathtt{str}}$. That the symmetry τ passes to a symmetry $\tau_{\mathtt{str}}$ of $\underline{C}_{\mathtt{str}}$ was noted already in 4.6. As for the left duality, a generic element of $\underline{C}_{\mathtt{str}}$ is a finite sequence (X_1, \ldots, X_n) and

$$^{\vee}(X_{1},...,X_{n}) = (^{\vee}X_{n},...,^{\vee}X_{1}),$$

where ϵ and η are defined in the obvious way. It is also clear that the twist

on \underline{C} can be brought over to a twist on \underline{C}_{str} . Accordingly, $\gamma:\underline{C}\to\underline{C}_{str}$ is a symmetric monoidal equivalence which is twist preserving, i.e., $\gamma:\underline{C}\to\underline{C}_{str}$ is a ribbon equivalence.

§7. *-CATEGORIES

Let \underline{k} be a commutative ring -- then a category \underline{C} is \underline{k} -enriched if \forall X,Y \in Ob \underline{C} , Mor(X,Y) is a \underline{k} -module and if the composition of morphisms is \underline{k} -bilinear. A functor F between \underline{k} -enriched categories is \underline{k} -linear if the induced maps

$$Mor(X,Y) \rightarrow Mor(FX,FY)$$

are homomorphisms of k-modules.

[Note: If \underline{C} is \underline{k} -enriched and monoidal, then $\underline{C} \times \underline{C}$ is \underline{k} -enriched and the functor $\underline{Q}:\underline{C} \times \underline{C} \to \underline{C}$ is assumed to be \underline{k} -bilinear.]

<u>N.B.</u> An object X in a \underline{k} -enriched category \underline{C} is $\underline{irreducible}$ if $Mor(X,X) = \underline{kid}_X$.

7.1 EXAMPLE Suppose that C is Z-enriched and monoidal. Put $\underline{k} = \underline{M}(\underline{C}).$

Then \underline{k} is a unital commutative ring (cf. 1.4) and \underline{C} is \underline{k} -enriched as a monoidal category (cf. 1.5).

[Note: Suppose in addition that \underline{C} is a ribbon category — then $\forall \ X \in Ob \ \underline{C}$,

$$tr_X:Mor(X,X) \rightarrow \underline{k}$$

is \underline{k} -linear and $\forall X,Y \in Ob \underline{C}$, the map

is <u>k</u>-bilinear.]

A *-category is a pair $(\underline{C},*)$, where \underline{C} is a category enriched over the field of complex numbers and

is an involutive, identity on objects, positive cofunctor. Spelled out: $\forall \ X,Y \in Cb \ \underline{C}, \ Mor(X,Y) \ is \ a \ complex \ vector \ space, \ composition$

$$Mor(X,Y) \times Mor(Y,Z) \rightarrow Mor(X,Z)$$

is complex bilinear,

$$\star:Mor(X,Y) \rightarrow Mor(Y,X)$$

subject to

$$(zf + wg)* = \overline{z}f* + \overline{w}g*$$

and

Finally, the requirement that * be positive means:

$$f* \circ f = 0 \Rightarrow f = 0.$$

[Note: $\forall X \in Ob C$, we have

$$id_{X}^{*} = id_{X} \circ id_{X}^{*}$$

$$= id_{X}^{**} \circ id_{X}^{*}$$

$$= (id_{X} \circ id_{X}^{*})^{*}$$

$$= id_{X}^{**}$$

$$= id_{X}^{*}$$

N.B. A monoidal *-category is a *-category which is monoidal with

$$(f \otimes g)^* = f^* \otimes g^*$$

for all f,g.

[Note: A symmetric monoidal *-category is a monoidal *-category such that $\forall \ X,Y \in Ob \ C$,

$$T_{X,Y}: X \otimes Y \rightarrow Y \otimes X$$

is unitary (see below).]

7.2 EXAMPLE FDHILB is a symmetric monoidal *-category.

[Note: For the record, <u>FDHILB</u> is a construct. As such, it is amnestic and transportable, thus there is no loss of generality in assuming that its monoidal structure is strict (cf. 3.12).]

7.3 REMARK Let A be a complex *-algebra -- then the involution is positive if $A^* \circ A = 0 \Rightarrow A = 0$ (A \in A). To illustrate, take $A = M_2(C)$ and consider the involutions

Then $*_1$ is positive but $*_2$ is not positive since

[Note: It is wellknown that if A is finite dimensional and if the involution is positive, then A is a semisimple algebra, hence "is" a multimatrix algebra.]

Let $f:X \to Y$ be a morphism in a *-category \underline{C} -- then f is an $\underline{isometry}$ if $f^* \circ f = id_X$ and f is $\underline{unitary}$ if both f and f^* are isometries.

Let F be a C-linear functor between *-categories -- then F is *-preserving if \forall f, $F(f^*) = (Ff)^*$.

N.B. Suppose that F is a *-preserving monoidal functor between monoidal *-categories — then F is unitary if the isomorphisms $\xi:e^{t} \to Fe$ and

$$\Xi_{X,Y}$$
:FX &' FY \rightarrow F(X & Y)

are unitary.

Let $p:X \to X$ be a morphism in a *-category C -- then p is a <u>projection</u> if $p = p^*$ and $p \circ p = p$.

[Note: If $g:Y \to X$ is an isometry, then $g \circ g^*:X \to X$ is a projection.] Let C be a *-category and let $X,Y \in Ob$ C -- then X is a subobject of Y if B an isometry $f \in Mor(X,Y)$.

Definition: \underline{C} has subobjects if for any $Y \in Ob$ \underline{C} and any projection $q \in Mor(Y,Y)$, $\exists X \in Ob$ \underline{C} and an isometry $f \in Mor(X,Y)$ such that $f \circ f^* = q$.

Definition: \underline{C} has direct sums if for all $X,Y \in Ob$ \underline{C} , \exists $Z \in Ob$ \underline{C} and isometries $f \in Mor(X,Z)$, $g \in Mor(Y,Z)$ such that $f \circ f^* + g \circ g^* = id_Z$.

E.g.: FDHILB has subobjects and direct sums.

7.4 RAPPEL A category C is essentially small if C is equivalent to a small category.

Suppose that \underline{C} is a *-category which is essentially small -- then \underline{C} is semisimple if the following conditions are met:

$$SS_1: \forall X,Y \in Ob \underline{C}$$

 $\dim Mor(X,Y) < \infty$.

 $\mathrm{SS}_2\colon \ \underline{\mathbf{C}}$ has subobjects and direct sums.

SS3: C has a zero object.

N.B. A monoidal *-category is semisimple if it is semisimple as a *-category and if in addition, e is irreducible.

- 7.5 EXAMPLE FDHILB is a semisimple strict monoidal *-category (cf. 7.2).
- 7.6 <u>LEMMA</u> Suppose that <u>C</u> is a semisimple *-category -- then every nonzero object in <u>C</u> is a finite direct sum of irreducible objects.

[$\forall X \in Ob \ \underline{C}$, Mor(X,X) is a finite dimensional complex *-algebra and the involution *:Mor(X,X) \rightarrow Mor(X,X) is positive (cf. 7.3).]

[Note: Conventionally, zero objects are not irreducible.]

Therefore a semisimple *-category is abelian.

Given a semisimple *-category \underline{C} , denote its set of isomorphism classes of irreducible objects by $\underline{I}_{\underline{C}}$ and let $\{X_{\underline{i}}: i\in \underline{I}_{\underline{C}}\}$ be a set of representatives -- then

$$i \neq j \Rightarrow Mor(X_i, X_j) = \{0\}$$

and $\forall \ X \in Ob \ C$, \exists a finite number of i such that

$$Mor(X_i,X) \neq \{0\},\$$

thereby defining $I_X \subset I_C$.

7.7 REMARK \forall i \in I_X, Mor(X_i,X) is a finite dimensional Hilbert space with inner product

$$<\phi,\psi>$$
 $id_{X_i} = \phi^* \circ \psi$.

7.8 <u>LFMMA</u> Let \underline{C} , \underline{C} ' be semisimple *-categories and suppose that $F:\underline{C} \to \underline{C}$ ' is C-linear -- then F is faithful if FX is nonzero for every irreducible X.

PROOF Consider an $f \in Mor(X,Y)$: Ff = 0, the claim being that f = 0. Fix orthonormal bases

$$\begin{bmatrix} & s_{ik} \in Mor(X_i, X) & (k = 1, ..., dim Mor(X_i, X)) \\ & t_{j\ell} \in Mor(Y_j, Y) & (\ell = 1, ..., dim Mor(Y_j, Y)) \end{bmatrix}$$

such that

$$\begin{bmatrix} \Sigma & s_{ik} & s_{ik}^* = id_X \\ \vdots & \vdots & \vdots \\ \Sigma & t_{j\ell} & t_{j\ell}^* = id_Y. \end{bmatrix}$$

Write

$$f = id_{Y} \circ f \circ id_{X}$$

$$= \sum_{ik,j\ell} t_{j\ell} \circ t_{j\ell}^* \circ f \circ s_{ik} \circ s_{ik}^*$$
$$= \sum_{ik\ell} c_{ik\ell} t_{i\ell} \circ s_{ik}^* (\exists c_{ik\ell} \in C).$$

Then for indices m, μ, ν ,

$$0 = F(t_{mv}^{*}) \circ Ff \circ F(s_{m\mu})$$

$$= \sum_{ik\ell} c_{ik\ell} F(t_{mv}^{*} \circ t_{i\ell} \circ s_{ik}^{*} \circ s_{m\mu})$$

$$= \sum_{k\ell} c_{mk\ell} F(t_{mv}^{*} \circ t_{m\ell} \circ s_{mk}^{*} \circ s_{m\mu})$$

$$= \sum_{k\ell} c_{mk\ell} F(t_{mv}^{*}, t_{m\ell}^{*}) id_{X_{m}} \circ (s_{mk}, s_{m\mu}^{*}) id_{X_{m}^{*}}$$

$$= c_{m\mu\nu} F(id_{X_{m}^{*}})$$

$$= c_{m\mu\nu} id_{FX_{m}^{*}}.$$

But by assumption, $id_{\overrightarrow{FX}_m} \neq 0$, thus the $c_{m\mu\nu}$ vanish, so f = 0.

7.9 <u>LEMMA</u> Let C, C' be semisimple *-categories and suppose that $F:C \to C'$ is C-linear and faithful -- then F is full iff (a) $X \in Ob$ C irreducible => $FX \in Ob$ C' irreducible and (b) $X,Y \in Ob$ C irreducible and nonisomorphic => $FX,FY \in Ob$ C' irreducible and nonisomorphic.

§8. NATURAL TRANSFORMATIONS

Let \underline{C} , \underline{C}' be *-categories and let $F:\underline{C} \to \underline{C}'$ be a *-preserving functor.

8.1 LEMMA Nat(F,F) is a unital *-algebra under the following operations:

$$(\alpha \circ \beta)_{X} = \alpha \alpha_{X} + b \beta_{X}$$

$$(\alpha \circ \beta)_{X} = \alpha_{X} \circ \beta_{X}$$

$$(\alpha^{*})_{X} = (\alpha_{X})^{*}$$

$$1_{X} = id_{FX}.$$

[To check the *-condition, observe that \forall f \in Mor(X,Y),

Ff •
$$(\alpha^*)_X$$
 = Ff • $(\alpha_X)^*$
= $(Ff^*)^* • (\alpha_X)^*$
= $(\alpha_X • Ff^*)^*$
= $(Ff^* • \alpha_Y)^*$
= $(\alpha_Y)^* • (Ff^*)^*$
= $(\alpha^*)_Y • Ff.$

8.2 EXAMPLE Take $\underline{C}' = \underline{FDHILB}$, put $Nat_F = Nat(F,F)$, and let $\underline{Rep}_{fd} Nat_F$ be the category whose objects are the finite dimensional *-representations of

 $\operatorname{Nat}_{\overline{F}}$ and whose morphisms are the intertwining operators. Define a *-preserving functor

$$\Phi:\underline{C} \to \underline{Rep}_{fd} \text{ Nat}_F$$

as follows:

$$- \Phi X = (\pi_X, FX) \quad (X \in Ob C)$$

$$- \Phi f = Ff \quad (f \quad Mor(X,Y)).$$

Here

$$\pi_{\mathbf{X}}(\alpha) = \alpha_{\mathbf{X}}$$

thus the diagram

commutes, so Ff is an intertwining operator.

[Note: If

is the forgetful functor, i.e., $U(\pi,H)=H$, then $U\circ\Phi=F.$]

8.3 THEOREM Let C, C' be *-categories and let $F:C \to C'$ be a *-preserving functor. Assume: C is semisimple -- then there is an isomorphism

$$\Psi_{F}: Nat(F,F) \rightarrow \underset{i \in I_{\underline{C}}}{\uparrow \uparrow} Mor(FX_{i},FX_{i})$$

of unital *-algebras.

 $\underline{\text{PROOF}}$ The definition of $\Psi_{\mathbf{F}}$ is the obvious one:

$$\Psi_{\mathbf{F}}(\alpha) = \prod_{\mathbf{i} \in \mathbf{I}_{\underline{\mathbf{C}}}} \alpha_{\mathbf{X}_{\mathbf{i}}}.$$

 Ψ_{F} is injective:

$$\alpha_{X_{\underline{i}}} = 0 \ \forall \ \underline{i} \in \underline{I}_{\underline{C}} \Rightarrow \alpha_{X} = 0 \ \forall \ X \in Ob \ \underline{C}.$$

To see this, choose the $s_{ik} \in Mor(X_i,X)$ as in the proof of 7.8 -- then

$$\alpha_{X} = \alpha_{X} \circ \text{Fid}_{X}$$

$$= \sum_{ik} \alpha_{X} \circ F(s_{ik} \circ s_{ik}^{*})$$

$$= \sum_{ik} \alpha_{X} \circ F(s_{ik}) \circ F(s_{ik}^{*}).$$

But the diagram

$$F(s_{ik}) \downarrow FX_{i} \longrightarrow FX_{i} \downarrow F(s_{ik})$$

$$FX \longrightarrow FX$$

$$FX \longrightarrow FX$$

commutes, hence

$$\alpha_{X} = \sum_{ik} F(s_{ik}) \circ \alpha_{X_{i}} \circ F(s_{ik}^{*})$$

$$= 0.$$

Y is surjective:

$$\forall \ \{\alpha_{\underline{\mathbf{i}}} \in \mathtt{Mor}(\mathtt{FX}_{\underline{\mathbf{i}}},\mathtt{FX}_{\underline{\mathbf{i}}}) : \underline{\mathbf{i}} \in \mathtt{I}_{\underline{\mathbf{C}}}^{\mathtt{J}}\}, \ \exists \ \alpha \in \mathtt{Nat}(\mathtt{F},\mathtt{F}) : \Psi_{\underline{\mathbf{F}}}(\alpha) = \prod_{\underline{\mathbf{i}} \in \mathtt{I}_{\underline{\mathbf{C}}}} \alpha_{\underline{\mathbf{i}}}.$$

Thus define $\alpha_{X} \in Mor(FX,FX)$ by

$$\alpha_{X} = \sum_{ik} F(s_{ik}) \circ \alpha_{i} \circ F(s_{ik}^{*})$$

$$\alpha_{Y} = \sum_{j\ell} F(t_{j\ell}) \circ \alpha_{j} \circ F(t_{j\ell}^{*}).$$

Then \forall f \in Mor(X,Y),

$$\begin{split} &\text{Ff} \, \circ \, \alpha_{\mathbf{X}} = \, \underset{\mathbf{ik}}{\Sigma} \, \, \mathbf{F}(\mathbf{f} \, \circ \, \mathbf{s}_{\mathbf{ik}}) \, \circ \, \alpha_{\mathbf{i}} \, \circ \, \mathbf{F}(\mathbf{s}_{\mathbf{ik}}^{\star}) \\ &= \, \underset{\mathbf{ik}, \, \mathbf{j}\ell}{\Sigma} \, \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell} \, \circ \, \mathbf{t}_{\mathbf{j}\ell}^{\star} \, \circ \, \mathbf{f} \, \circ \, \mathbf{s}_{\mathbf{ik}}) \, \circ \, \alpha_{\mathbf{i}} \, \circ \, \mathbf{F}(\mathbf{s}_{\mathbf{ik}}^{\star}) \\ &= \, \underset{\mathbf{ik}\ell}{\Sigma} \, \, \mathbf{F}(\mathbf{t}_{\mathbf{i}\ell} \, \circ \, (\mathbf{t}_{\mathbf{i}\ell}^{\star} \, \circ \, \mathbf{f} \, \circ \, \mathbf{s}_{\mathbf{ik}})) \, \circ \, \alpha_{\mathbf{i}} \, \circ \, \mathbf{F}(\mathbf{s}_{\mathbf{ik}}^{\star}) \\ &= \, \underset{\mathbf{ik}\ell}{\Sigma} \, \, \mathbf{F}(\mathbf{t}_{\mathbf{i}\ell}) \, \circ \, \mathbf{F}(\mathbf{t}_{\mathbf{i}\ell}^{\star} \, \circ \, \mathbf{f} \, \circ \, \mathbf{s}_{\mathbf{ik}}) \, \circ \, \mathbf{G}_{\mathbf{i}} \, \circ \, \mathbf{F}(\mathbf{s}_{\mathbf{ik}}^{\star}) \\ &= \, \underset{\mathbf{ik}\ell}{\Sigma} \, \, \mathbf{F}(\mathbf{t}_{\mathbf{i}\ell}) \, \circ \, \alpha_{\mathbf{i}} \, \circ \, \mathbf{F}(\mathbf{t}_{\mathbf{i}\ell}^{\star} \, \circ \, \mathbf{f} \, \circ \, \mathbf{s}_{\mathbf{ik}}) \, \circ \, \mathbf{F}(\mathbf{s}_{\mathbf{ik}}^{\star}) \\ &= \, \underset{\mathbf{ik}, \, \mathbf{j}\ell}{\Sigma} \, \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell}) \, \circ \, \alpha_{\mathbf{j}} \, \circ \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell}^{\star} \, \circ \, \mathbf{f} \, \circ \, \mathbf{s}_{\mathbf{ik}}) \, \circ \, \mathbf{S}_{\mathbf{ik}}^{\star}) \\ &= \, \underset{\mathbf{ik}, \, \mathbf{j}\ell}{\Sigma} \, \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell}) \, \circ \, \alpha_{\mathbf{j}} \, \circ \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell}^{\star} \, \circ \, \mathbf{f} \, \circ \, \mathbf{s}_{\mathbf{ik}}) \, \circ \, \mathbf{S}_{\mathbf{ik}}^{\star}) \\ &= \, \underset{\mathbf{ik}, \, \mathbf{j}\ell}{\Sigma} \, \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell}) \, \circ \, \alpha_{\mathbf{j}} \, \circ \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell}^{\star} \, \circ \, \mathbf{f} \, \circ \, \mathbf{s}_{\mathbf{ik}}) \\ &= \, \underset{\mathbf{ik}, \, \mathbf{j}\ell}{\Sigma} \, \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell}) \, \circ \, \alpha_{\mathbf{j}} \, \circ \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell}^{\star} \, \circ \, \mathbf{f} \, \circ \, \mathbf{s}_{\mathbf{ik}}) \\ &= \, \underset{\mathbf{ik}, \, \mathbf{j}\ell}{\Sigma} \, \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell}) \, \circ \, \alpha_{\mathbf{j}} \, \circ \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell}^{\star} \, \circ \, \mathbf{f} \, \circ \, \mathbf{s}_{\mathbf{ik}}) \\ &= \, \underset{\mathbf{ik}, \, \mathbf{j}\ell}{\Sigma} \, \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell}) \, \circ \, \alpha_{\mathbf{j}} \, \circ \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell}^{\star} \, \circ \, \mathbf{f} \, \circ \, \mathbf{s}_{\mathbf{ik}}) \\ &= \, \underset{\mathbf{ik}, \, \mathbf{j}\ell}{\Sigma} \, \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell}) \, \circ \, \alpha_{\mathbf{j}} \, \circ \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell}^{\star} \, \circ \, \mathbf{f} \, \circ \, \mathbf{f} \, \circ \, \mathbf{f} \, \circ \, \mathbf{s}_{\mathbf{ik}}) \\ &= \, \underset{\mathbf{ik}, \, \mathbf{j}\ell}{\Sigma} \, \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell}) \, \circ \, \alpha_{\mathbf{j}} \, \circ \, \mathbf{F}(\mathbf{t}_{\mathbf{j}\ell}^{\star} \, \circ \, \mathbf{f} \, \bullet \,$$

Accordingly, the diagram

commutes, meaning that $\alpha \in \text{Nat}(\texttt{F},\texttt{F})$. And, by construction, $\alpha_{X_i} = \alpha_i$, so

$$\Psi_{\mathbf{F}}(\alpha) = \prod_{\mathbf{i} \in I_{\underline{\mathbf{C}}}} \alpha_{\mathbf{i}}.$$

[Note: The isomorphism $\boldsymbol{\Psi}_{\!F}$ depends on the choice of the $\mathbf{X}_{\!\mathbf{i}}.]$

8.4 EXAMPLE Take $\underline{C}' = \underline{C}$ and let $F = \mathrm{id}_{\underline{C}}$ (the identity functor) — then

$$Nat(id_{\underline{C}}, id_{\underline{C}}) \approx \prod_{i \in I_{\underline{C}}} C.$$

8.5 EXAMPLE Suppose that C is a semisimple monoidal *-category -- then $C \times C$ is a semisimple *-category with

$$I^{\overline{C}} \times \overline{C} = I^{\overline{C}} \times I^{\overline{C}}$$

And

$$x_i \boxtimes x_j \approx \bigoplus_{k \in I_{\underline{C}}} N_{ij}^k x_k,$$

where

$$N_{ij}^{k} = \dim Mor(X_{k}, X_{i} \otimes X_{j}),$$

so

$$\text{Mor}(\mathbf{X}_{\underline{\mathbf{i}}} \ \underline{\otimes} \ \mathbf{X}_{\underline{\mathbf{j}}}, \mathbf{X}_{\underline{\mathbf{i}}} \ \underline{\otimes} \ \mathbf{X}_{\underline{\mathbf{j}}}) \ \approx \ \underset{k \in \mathbf{I}_{\underline{\mathbf{C}}}}{\oplus} \ \underset{N}{\mathbf{M}_{\underline{\mathbf{k}}}} \ (\mathtt{C}).$$

This said, let

$$F = \Omega: C \times C \rightarrow C$$
.

Then

$$Nat(\mathbf{Q},\mathbf{Q}) \approx \prod_{i,j \in I_{\underline{C}}} Mor(\mathbf{Q}(\mathbf{X}_{i},\mathbf{X}_{j}), \mathbf{Q}(\mathbf{X}_{i},\mathbf{X}_{j}))$$

$$\approx \prod_{i,j \in I_{\underline{C}}} Mor(X_i \boxtimes X_j, X_i \boxtimes X_j)$$

Suppose that \underline{C} is a semisimple *-category, let $F:\underline{C} \to \underline{FDHILB}$ be *-preserving and put

$$A_{F} = \bigoplus_{i \in I_{C}} B(FX_{i})$$

which, of course, can be embedded in

$$\prod_{i \in I_{\underline{C}}} B(FX_i) \ (\approx Nat(F,F)).$$

Needless to say, A_F is a *-algebra, unital iff $I_{\underline{C}}$ is finite. The projections $p_i \colon A_F \to B(FX_i)$ are finite dimensional irreducible *-representations. Moreover, any finite dimensional nondegenerate *-representation of A_F is a direct sum of finite dimensional irreducible *-representations and every finite dimensional irreducible *-representation is unitarily equivalent to a p_i .

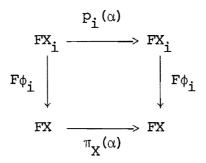
Define now a *-preserving functor

$$\Phi:\underline{C} \to \underline{\text{Rep}}_{fd} A_F$$

as in 8.2 -- then Φ is an equivalence of categories iff F is faithful. In fact, since Φ and F agree on morphisms, it is clear that

Φ faithful <=> F faithful.

Assume therefore that F is faithful. From the definitions, $\pi_{X_{\dot{1}}} = p_{\dot{1}}$ (or still, $\forall \alpha \in A$, $\alpha_{X_{\dot{1}}} = p_{\dot{1}}(\alpha)$), which is a finite dimensional irreducible *-representation of A_{F} . Given an irreducible $X \in Ob \ \underline{C}$, $\exists \ i \in I_{\underline{C}}$ and an isomorphism $\phi_{\dot{1}}: X_{\dot{1}} \to X$. Since the diagram



commutes, π_X is also a finite dimensional irreducible *-representation of $A_F^{}.$ If i \neq j, then

$$Mor(p_{i}, p_{j}) = \{0\},\$$

so if $X,Y\in Ob\ \underline{C}$ are irreducible and nonzero, then

Mor(
$$\pi_{X'}^{\pi_{Y'}} = \{0\}.$$

Because Φ is faithful (and $\operatorname{\underline{Rep}_{fd}} A_F$ is a semisimple *-category), the foregoing considerations imply that Φ is full (cf. 7.9). Finally, Φ has a representative image. Indeed, as mentioned above, every finite dimensional irreducible *-representation of A_F is unitarily equivalent to a P_i .

To recapitulate:

8.6 THEOREM Let \underline{C} be a semisimple *-category and let $F:\underline{C} \to \underline{FDHILB}$ be a *-preserving functor. Put

$$A_{\mathbf{F}} = \bigoplus_{\mathbf{i} \in \mathbf{I}_{\mathbf{C}}} B(\mathbf{F} \mathbf{X}_{\mathbf{i}})$$

and define

$$\Phi: \underline{C} \to \underline{Rep}_{fd} A_F$$

by

$$\Phi X = (\pi_X, FX) \quad (X \in Ob \underline{C})$$

$$\Phi f = Ff \quad (f \in Mor(X,Y)).$$

Then Φ is an equivalence of categories iff F is faithful.

Let C be a semisimple strict monoidal *-category.

Definition: An embedding functor (for C) is a faithful unitary functor

[Note: Recall from §7 that in this context, "unitary" means that F is a *-preserving monoidal functor for which the isomorphisms $\xi:\underline{e} \to Fe$ and

$$\Xi_{X,Y}$$
: FX $\underline{\omega}$ FY \rightarrow F(X $\underline{\omega}$ Y)

are unitary (\underline{e} = standard unit in \underline{FDHILB} , $\underline{\omega}$ = strict monoidal structure in \underline{FDHILB} (cf. 7.5)).]

8.7 LEMMA There is an isomorphism

$$\Psi_{\mathbf{F}}: \mathbb{N} \text{at}(\mathbf{F}, \mathbf{F}) \to \prod_{\mathbf{i} \in \mathbf{I}_{\mathbf{C}}} \mathbb{B}(\mathbf{F} \mathbf{X}_{\mathbf{i}})$$

of unital *-algebras (cf. 8.3).

8.8 LEMMA The map

$$\varepsilon_{\mathbf{F}}$$
:Nat(F,F) \rightarrow Mor(Fe,Fe) \approx C

that sends

$$\alpha = \{\alpha_X^{}\} \text{ to } \alpha_e^{}$$

is a unital *-homomorphism.

8.9 SCHOLIUM The map

$$\bar{\varepsilon}_{\mathbf{F}} : \prod_{\mathbf{i} \in \mathbf{I}_{\underline{\mathbf{C}}}} \mathbf{B}(\mathbf{F} \mathbf{X}_{\mathbf{i}}) \rightarrow \mathbf{C}$$

that sends

T to
$$\varepsilon_{\mathbf{F}} \circ \Psi_{\mathbf{F}}^{-1}(\mathbf{T})$$

is a unital *-homomorphism.

Let

$$\varepsilon = \bar{\varepsilon}_{F} | A_{F}.$$

Then ϵ is a unital *-homomorphism, the counit.

8.10 LEMMA There is an isomorphism

$$\Psi_{F} \circ \mathbf{\Omega}: Nat(F \circ \mathbf{\Omega}, F \circ \mathbf{\Omega}) \longrightarrow \prod_{i,j \in I_{\underline{C}}} B(FX_{i}) \otimes_{C} B(FX_{j})$$

of unital *-algebras.

PROOF In fact,

$$\approx \prod_{i,j \in I_{\underline{C}}} B(F(X_i \boxtimes X_j)) \quad (cf. 8.3)$$

$$\approx \prod_{i,j \in I_{\underline{C}}} B(FX_i \boxtimes' FX_j)$$

$$\underset{i,j \in I_{\underline{C}}}{\text{H}} \text{B}(\text{FX}_{i}) \otimes_{\underline{C}} \text{B}(\text{FX}_{j}).$$

8.11 LEMMA The map

$$\Delta_{F}$$
:Nat(F,F) \rightarrow Nat(F $\circ \otimes$,F $\circ \otimes$)

that sends

$$\alpha = \{\alpha_{\mathbf{X}}\} \text{ to } \{\alpha_{\mathbf{X} \mathbf{Q} \mathbf{Y}}\}$$

is a unital *-homomorphism.

8.12 SCHOLIUM The map

$$\bar{\Delta}_{F}: \prod_{i \in I_{\underline{C}}} B(FX_{i}) \rightarrow \prod_{i,j \in I_{\underline{C}}} B(FX_{i}) \otimes_{C} B(FX_{j})$$

that sends

T to
$$\Psi_{F \circ Q} \circ \Lambda_{F} \circ \Psi_{F}^{-1}$$
 (T)

is a unital *-homomorphism.

Let

$$\Delta = \overline{\Delta}_{\mathbf{F}} | A_{\mathbf{F}}.$$

Then \triangle is a unital *-homomorphism, the coproduct.

Let

$$\begin{array}{c|c}
 & \pi_1: A_F \to B(H_1) \\
 & \pi_2: A_F \to B(H_2)
\end{array}$$

be nondegenerate *-representations of $\mathbf{A}_{\mathbf{F}}$ on finite dimensional Hilbert spaces

(the zero representation is a possibility) — then we can form H_2

$$\pi_1 \boxtimes \pi_2 \colon A_F \boxtimes_{\mathbb{C}} A_F \to B(H_1) \boxtimes_{\mathbb{C}} B(H_2) \approx B(H_1 \boxtimes H_2).$$

Since

$$A_{F} \otimes_{C} A_{F} = \bigoplus_{i,j \in I_{C}} B(FX_{i}) \otimes_{C} B(FX_{j}),$$

it follows that $\pi_1 \ \text{@} \ \pi_2$ admits a unique extension to a unital *-homomorphism

$$\overline{\pi_1 \boxtimes \pi_2} : \prod_{i,j \in I_C} B(FX_i) \boxtimes_C B(FX_j) \rightarrow B(H_1 \boxtimes H_2).$$

This being so, put

$$\pi_1 \times \pi_2 = \overline{\pi_1 \otimes \pi_2} \circ \Delta.$$

Then $\pi_1 \times \pi_2$ is a nondegenerate *-representation of A_F on the finite dimensional Hilbert space $H_1 \ \underline{@} \ H_2$.

8.13 <u>LEMMA</u> The data $(\times, \epsilon, ...)$ is a monoidal structure on $\underline{\text{Rep}}_{fd}$ A_F .

Therefore $\underline{\text{Rep}}_{fd}$ A_F is a semisimple monoidal *-category (the counit ϵ is the irreducible unit).

8.14 THEOREM Let \underline{C} be a semisimple strict monoidal *-category and let $F:\underline{C} \to FDHILB$

be an embedding functor. Put

$$A_{\mathbf{F}} = \bigoplus_{\mathbf{i} \in \mathbf{I}_{\underline{\mathbf{C}}}} \mathbf{B}(\mathbf{F}\mathbf{X}_{\mathbf{i}})$$

and define

by

Then Φ is a monoidal equivalence.

<u>PROOF</u> By hypothesis, F is faithful, hence Φ is an equivalence of categories (cf. 8.6). So, in view of 2.8, it suffices to show that Φ is monoidal. There are two points. First

$$\Phi e = (\pi_{\alpha}, Fe)$$

and $\forall~\alpha \in A_{\mathbf{F}}\text{,}~\text{the diagram}$

$$\begin{array}{ccc}
C & \xrightarrow{\varepsilon(\alpha)} & C \\
\xi & & \downarrow & \xi \\
Fe & \xrightarrow{\pi_{e}(\alpha)} & Fe
\end{array}$$

commutes, i.e., ξ intertwines ϵ and $\pi_{\underline{e}}$. Next, given $X,Y\in Ob$ \underline{C} , consider

Then

$$\Xi_{X,Y}$$
: FX $\underline{\otimes}$ FY \rightarrow F(X $\underline{\otimes}$ Y)

is an intertwining operator: $\forall \ \alpha \in A_{\mathbf{F}}$,

$$\Xi_{X,Y} \circ (\pi_X \times \pi_Y) (\alpha) = \pi_{X \boxtimes Y} (\alpha) \circ \Xi_{X,Y}$$

The interchange $\sigma: A_F \otimes_C A_F \to A_F \otimes_C A_F$ ($\sigma(\alpha \otimes \beta) = \beta \otimes \alpha$) is a nondegenerate *-homomorphism, thus has a unique extension to an involutive *-automorphism

$$\overline{\sigma} \colon \prod_{i,j \in I_{\underline{C}}} B(FX_i) \otimes_{\underline{C}} B(FX_j) \to \prod_{i,j \in I_{\underline{C}}} B(FX_i) \otimes_{\underline{C}} B(FX_j).$$

Let

$$\Delta^{OP} = \bar{\sigma} \circ \Delta$$
.

Then A_F is said to be <u>cocommutative</u> if $\Delta = \Delta^{OP}$.

8.15 $\underline{\text{LEMMA}}$ Suppose that A_{F} is cocommutative -- then

$$\forall \begin{bmatrix} (\pi_1, H_1) \\ \in Ob & \underline{Rep}_{fd} A_F, \\ (\pi_2, H_2) \end{bmatrix}$$

the diagram

commutes.

$$\forall T \in \prod_{i,j \in I_{\underline{C}}} B(FX_i) \otimes_{\underline{C}} B(FX_j),$$

we have

$$\overline{(\pi_1 \boxtimes \pi_2)} \overline{\sigma}(T) = \tau_{2,1} \overline{(\pi_2 \boxtimes \pi_1)} (T) \tau_{1,2}.$$

So, $\forall \alpha \in A_{F}$,

$$\begin{array}{l}
\tau_{1,2}^{(\pi_{1} \times \pi_{2})(\alpha)} \\
= \tau_{1,2}^{(\overline{\pi_{1} \otimes \pi_{2}})(\Delta(\alpha))} \\
= \tau_{1,2}^{(\overline{\pi_{1} \otimes \pi_{2}})(\Delta^{OP}(\alpha))} \\
= \tau_{1,2}^{(\overline{\pi_{1} \otimes \pi_{2}})(\overline{\sigma}(\Delta(\alpha)))} \\
= \tau_{1,2}^{(\overline{\pi_{1} \otimes \pi_{2}})(\overline{\sigma}(\Delta(\alpha)))} \\
= \tau_{1,2}^{(\overline{\pi_{2} \otimes \pi_{1}})(\Delta(\alpha))\tau_{1,2}} \\
= (\overline{\pi_{2} \otimes \pi_{1}})(\Delta(\alpha))\tau_{1,2} \\
= (\pi_{2}^{(\overline{\pi_{2} \otimes \pi_{1}})(\alpha)\tau_{1,2}^{(\overline{\pi_{2} \otimes \pi_{1}})(\alpha)\tau_{1,2}^{(\overline$$

Thus, if A_F is cocommutative, then $\underline{\text{Rep}}_{\mathrm{fd}}$ A_F is a semisimple symmetric monoidal *-category.

8.16 $\underline{\text{REMARK}}$ Assume further that the category $\underline{\text{C}}$ is symmetric and that the embedding functor

$$F:\underline{C} \to \underline{FDHILB}$$

is symmetric monoidal — then ${\bf A_F}$ is cocommutative and $\Phi: \underline{{\bf C}} \to \underline{{\bf Rep}}_{fd} \ {\bf A_F}$ is a symmetric monoidal equivalence.

§9. CONJUGATES

Suppose that \underline{C} is a strict monoidal *-category which is left autonomous. Put $X^V = {}^V X$ -- then

=>

$$\begin{bmatrix} - & \varepsilon_{\mathbf{X}}^{\star} : \mathbf{e} \to {}^{\mathsf{V}} \mathbf{X} & \mathbf{X} \\ \\ & \eta_{\mathbf{X}}^{\star} : \mathbf{X} & \mathbf{X} & \mathbf{Y} \to \mathbf{e} \end{bmatrix}$$

=>

And

$$\begin{bmatrix} -(\operatorname{id}_X \otimes \varepsilon_X) \circ (\eta_X \otimes \operatorname{id}_X) = \operatorname{id}_X \\ (\varepsilon_X \otimes \operatorname{id}_Y) \circ (\operatorname{id}_Y \otimes \eta_X) = \operatorname{id}_Y \\ -(\varepsilon_X \otimes \operatorname{id}_Y) \otimes (\operatorname{id}_Y \otimes \eta_X) = \operatorname{id}_Y \end{bmatrix}$$

=>

$$\begin{bmatrix} - & (\eta_X^* \otimes id_X) & \circ & (id_X \otimes \epsilon_X^*) & = id_X \\ & (id_X \otimes \eta_X^*) & \circ & (\epsilon_X^* \otimes id_X) & = id_X \\ & & & & X^{\vee} & & X^{\vee} . \end{bmatrix}$$

I.e.: The left duality $({}^{\vee}X,\epsilon_{X},\eta_{X})$ automatically leads to a right duality $(X^{\vee},\eta_{X}^{\star},\epsilon_{X}^{\star})\,.$

Now assume in addition that \underline{c} is symmetric (hence that the $\tau_{X,Y}$ are unitary) — then the left duality (${}^{V}X, \epsilon_{X}, \eta_{X}$) gives rise to another right duality, viz.

$$(X^{\vee}, \varepsilon_X \circ \tau_{X, \vee X}, \tau_{X, \vee X} \circ \eta_X).$$

9.1 COHERENCY HYPOTHESIS $\forall X \in Ob C$,

$$\varepsilon_{X}^{\star} = \tau_{X, Y_{X}} \circ \eta_{X}$$

[Note: The asymmetry is only apparent. For

$$\eta_{X} = \tau_{X, X}^{-1} \circ \varepsilon_{X}^{*}$$

$$= \mathsf{T}_{\mathsf{V}_{\mathsf{X},\mathsf{X}}} \circ \varepsilon_{\mathsf{X}}^{\star}$$

=>

$$\eta_{X}^{\star} = \varepsilon_{X} \circ \tau_{X,X}^{\star}$$

$$= \varepsilon_{X} \circ \tau_{X,X}^{-1}$$

$$= \varepsilon_{X} \circ \tau_{X,X}^{-1}$$

In the presence of 9.1, let

$$\vec{x} = {}^{\vee}X (= X^{\vee})$$

$$r_{X} = \varepsilon_{X}^{\star}$$

$$\vec{r}_{X} = \tau_{X} \circ r_{X}^{\prime}$$

thus

and

$$(\bar{X}, r_{\bar{X}}^*, r_{\bar{X}}^*)$$
 is a left duality $(\bar{X}, \bar{r}_{\bar{X}}^*, r_{\bar{X}})$ is a right duality.

Therefore

$$\begin{bmatrix} - & (\operatorname{id}_X \otimes r_X^*) & \circ & (\overline{r}_X \otimes \operatorname{id}_X) & = \operatorname{id}_X \\ (r_X^* \otimes \operatorname{id}_X) & \circ & (\operatorname{id}_X \otimes \overline{r}_X) & = \operatorname{id}_X \\ - & (\overline{r}_X^* \otimes \operatorname{id}_X) & \circ & (\operatorname{id}_X \otimes r_X) & = \operatorname{id}_X \\ & & (\operatorname{id}_X \otimes \overline{r}_X^*) & \circ & (r_X \otimes \operatorname{id}_X) & = \operatorname{id}_X \\ - & & \overline{X} \end{bmatrix}$$

The relations

$$\begin{bmatrix} - & (\mathrm{id}_X \otimes r_X^*) & \circ & (\overline{r}_X \otimes \mathrm{id}_X) = \mathrm{id}_X \\ & (\mathrm{id}_X \otimes \overline{r}_X^*) & \circ & (r_X \otimes \mathrm{id}_X) = \mathrm{id}_X \\ & & & \overline{x} & & \overline{x} \end{bmatrix}$$

are called the conjugate equations, the triple $(\bar{X}, r_{\bar{X}}, \bar{r_{\bar{X}}})$ being a conjugate for X.

N.B. The conjugate equations imply that

$$(\overline{r}_{X}^{*} \otimes \operatorname{id}_{X}) \circ (\operatorname{id}_{X} \otimes r_{X}) = \operatorname{id}_{X}$$

$$(r_{X}^{*} \otimes \operatorname{id}_{X}) \circ (\operatorname{id}_{X} \otimes \overline{r}_{X}) = \operatorname{id}_{X}.$$

Having made these points, matters can be turned around. So start with a symmetric strict monoidal *-category \underline{C} — then \underline{C} has conjugates if one can assign to each $X \in Ob \ \underline{C}$ an object \overline{X} and a morphism

$$r_{x}:e \rightarrow \overline{x} \otimes x$$

such that the triple $(\bar{X}, r_X, \bar{r}_X)$ satisfies the conjugate equations (here, of course, $\bar{r}_X = \tau_X \circ r_X$).

E.g.: FDHILB has conjugates.

9.2 REMARK If C has conjugates, then C is left autonomous (consider $(\bar{X}, r_X^*, \bar{r}_X^*)$) and right autonomous (consider $(\bar{X}, \bar{r}_X^*, r_X^*)$). Moreover, the coherency hypothesis is in force: $(r_X^*)^* = r_X^*$, while

$$T \circ \overline{T}_{X} = T \circ T \circ T_{X} = T_{X}.$$

- 9.3 LEMMA Suppose that C has conjugates.
 - Under the identification

$$Mor(X \otimes Y,Z) \approx Mor(Y,\overline{X} \otimes Z)$$
,

the arrows

are mutually inverse.

• Under the identification

$$Mor(Y \boxtimes X,Z) \approx Mor(Y,Z \boxtimes \overline{X})$$
,

the arrows

are mutually inverse.

E.g.:
$$\forall X \in Ob \underline{C}$$
,

$$Mor(X,X) \approx Mor(e,\overline{X} \otimes X)$$
.

9.4 LEMMA If

$$(\bar{X}, r_{X}, \bar{r}_{X})$$

$$(\bar{X}', r_{X}', \bar{r}_{X}')$$

are conjugates for X, then

$$r_{X}^{\star} \boxtimes id_{\overline{X}'} \circ id_{\overline{X}} \boxtimes \overline{r}_{X}' \in Mor(\overline{X}, \overline{X}')$$

is unitary.

PROOF Put

$$U = r_X^* \boxtimes id \circ id \boxtimes \overline{r}_X'$$

$$(=\operatorname{id}_{\overline{X}'}\otimes \overline{r}_X^{\star}\circ r_X'\otimes \operatorname{id}_{\overline{X}}\ldots).$$

Then the claim is that

$$\begin{array}{c|c} U \circ U^* = id \\ \hline X' \\ U^* \circ U = id \\ \hline X \end{array}$$

And for this, it will be enough to consider U . U*. So write

$$\begin{array}{l} \text{U} \circ \text{U}^{\star} = r_{X}^{\star} \otimes \operatorname{id}_{\overline{X}} \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_{X}^{\star} \circ \operatorname{U}^{\star} \\ \\ = r_{X}^{\star} \otimes \operatorname{id}_{\overline{X}} \circ (\operatorname{id}_{\overline{X}} \otimes \overline{r}_{X}^{\star} \circ \operatorname{U}^{\star} \otimes \operatorname{id}_{e}) \\ \\ = r_{X}^{\star} \otimes \operatorname{id}_{\overline{X}} \circ \operatorname{U}^{\star} \otimes \overline{r}_{X}^{\star} \\ \\ = r_{X}^{\star} \otimes \operatorname{id}_{\overline{X}} \circ \operatorname{U}^{\star} \otimes \operatorname{id}_{\overline{X}} \otimes \overline{r}_{X}^{\star} \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_{X}^{\star}) \\ \\ = r_{X}^{\star} \otimes \operatorname{id}_{\overline{X}} \circ (\operatorname{U}^{\star} \otimes \operatorname{id}_{\overline{X}} \otimes \overline{r}_{X} \otimes \operatorname{id}_{\overline{X}}) \\ \\ = r_{X}^{\star} \otimes \operatorname{id}_{\overline{X}} \circ \operatorname{id}_{\overline{X}} \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_{X}^{\star} \otimes \operatorname{id}_{\overline{X}} \circ \circ \operatorname{id}_{\overline{X}} \circ \circ \circ \operatorname{id}_{\overline{X}} \circ \circ \operatorname{id}_{\overline{X}} \circ \circ \circ \circ \circ \circ \circ$$

$$\begin{array}{c} \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X \otimes \operatorname{id}_{X \otimes \overline{X}} \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X' \\ = \operatorname{id}_e \otimes (r_X^* \otimes \operatorname{id}_{\overline{X}}) \circ r_X^{**} \otimes \operatorname{id}_{\overline{X} \otimes X \otimes \overline{X}}, \\ \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X \otimes \operatorname{id}_{X \otimes \overline{X}}, \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X' \\ \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X \otimes \operatorname{id}_{X \otimes \overline{X}}, \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X' \\ = r_X^{**} \otimes r_X^* \otimes \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{X \otimes \overline{X}}, \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X' \\ = r_X^{**} \otimes \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{X \otimes \overline{X}}, \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X' \\ \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X \otimes \operatorname{id}_{X \otimes \overline{X}}, \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X' \\ = r_X^{**} \otimes \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X' \otimes \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X' \\ \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X' \otimes \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X' \otimes \operatorname{id}_{\overline{X}}, \\ \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X' \otimes \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}}, \otimes (\overline{r}_X \otimes \operatorname{id}_X) \otimes \operatorname{id}_{\overline{X}}, \\ \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X' \otimes \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}}, \otimes (\overline{r}_X \otimes \operatorname{id}_X) \otimes \operatorname{id}_{\overline{X}}, \\ \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X' \otimes \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}}, \otimes (\overline{r}_X \otimes \operatorname{id}_X) \otimes \operatorname{id}_{\overline{X}}, \\ \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X' \otimes \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}}, \otimes (\overline{r}_X \otimes \operatorname{id}_X) \otimes \operatorname{id}_{\overline{X}}, \\ \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_X' \otimes \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_$$

[Note: Evidently,

Conjugates are therefore determined up to "unitary equivalence".

Put

$$\Omega_{X} = r_{X}^{*} \otimes id_{X} \circ id_{X} \otimes r_{X,X} \circ r_{X} \otimes id_{X}.$$

Then

$$\Omega_{X} \in Mor(X,X)$$

is unitary and it can be verified by computation that the assignment $X \to \Omega_X$ defines a twist Ω . This fact, however, is a trivial consequence of the following result.

9.5 LEMMA
$$\forall X \in Ob C$$
,

$$\Omega_{X} = id_{X}$$
.

PROOF We have

$$(\bar{r}_X^* \boxtimes id_X) \circ (id_X \boxtimes r_X) = id_X.$$

On the other hand, there is a commutative diagram

so

$$(\vec{r}_X^* \otimes id_X) \circ (id_X \otimes r_X)$$

$$= \vec{r}_X^* \otimes id_X \circ \tau_X \circ r_X \otimes id_X$$

$$= \vec{r}_X^* \otimes id_X \circ \tau_X \otimes id_X \circ id_X \circ r_X \otimes id_X.$$

And

$$\bar{r}_{X}^{*} \otimes id_{X} \circ \tau_{\bar{X},X} \otimes id_{X}$$

--

$$= (\tau_{\overline{X},X} \circ r_{X})^{*} \otimes id_{X} \circ \tau_{\overline{X},X} \otimes id_{X}$$

$$= r_{X}^{*} \circ \tau_{X,\overline{X}} \otimes id_{X} \circ \tau_{\overline{X},X} \otimes id_{X}$$

$$= r_{X}^{*} \circ \tau_{X,\overline{X}} \circ \tau_{\overline{X},X} \otimes id_{X} \circ id_{X}$$

$$= r_{X}^{*} \otimes id_{X}.$$

[Note: Therefore, in the terminology of §6, C is an even ribbon category.]

9.6 REMARK \forall f \in Mor(X,X), the diagram

commutes. Therefore

$$\vec{r}_{X}^{\star} \circ f \boxtimes id_{\overline{X}} \circ \vec{r}_{X}$$

$$= (\tau_{X,X} \circ r_{X})^{\star} \circ f \boxtimes id_{\overline{X}} \circ (\tau_{X,X} \circ r_{X})$$

$$= r_{X}^{\star} \circ \tau_{X,\overline{X}} \circ f \boxtimes id_{\overline{X}} \circ \tau_{X,X} \circ r_{X}$$

$$= r_{X}^{\star} \circ id_{\overline{X}} \boxtimes f \circ r_{X}.$$

Maintaining the supposition that \underline{C} has conjugates, recall that \underline{C} is left autonomous with left duality $(\bar{X}, r_X^*, \bar{r}_X^*)$ (cf. 9.2), thus by definition the categorical dimension of X is the arrow

But

$$\bar{r}_{X} = \tau_{X,X} \circ r_{X'}$$

so the categorical dimension of X is the composition

$$r_{X}^{\star} \circ r_{X}^{\star} \circ r_{X}^{\star} \circ r_{X}^{\star}$$

$$= r_{X}^{\star} \circ r_{X}^{\star} \in Mor(e,e)$$

$$\equiv \dim X.$$

[Note: Since $\Omega = id$, $\forall f \in Mor(X,X)$,

$$tr_{X}(f) = r_{X}^{*} \circ \tau_{X,\overline{X}} \circ \Omega_{X} \otimes id_{\overline{X}} \circ (f \otimes id_{\overline{X}}) \circ \overline{r}_{X}$$

$$= r_{X}^{*} \circ \tau_{X,\overline{X}} \circ id_{\overline{X}} \circ (f \otimes id_{\overline{X}}) \circ \tau_{\overline{X},X} \circ r_{X}$$

$$= r_{X}^{*} \circ \tau_{X,\overline{X}} \circ (f \otimes id_{\overline{X}}) \circ \tau_{\overline{X},X} \circ r_{X}$$

$$= r_{X}^{*} \circ id_{\overline{X}} \otimes f \circ r_{X} \quad (cf. 9.6).]$$

N.B. dim X does not depend on the choice of a conjugate for X. Indeed, if $U: \bar{X} \to \bar{X}^i$ is unitary, then

$$((U \otimes id_{X}) \circ r_{X}) * \circ ((U \otimes id_{X}) \circ r_{X})$$

$$= r_{X}^{*} \circ U^{*} \otimes id_{X} \circ U \otimes id_{X} \circ r_{X}$$

$$= r_{X}^{*} \circ \operatorname{id}_{\overline{X}} \otimes \operatorname{id}_{X} \circ r_{X}$$

$$= r_{X}^{*} \circ \operatorname{id}_{\overline{X}} \otimes x \circ r_{X}$$

$$= r_{X}^{*} \circ r_{X}^{*}$$

9.7 LEMMA If

$$(\bar{X}, r_{X}, \bar{r}_{X})$$
 is a conjugate for X $(\bar{Y}, r_{Y}, \bar{r}_{Y})$ is a conjugate for Y,

then

$$(\overline{Y} \boxtimes \overline{X}, r_{X \boxtimes Y}, \overline{r}_{X \boxtimes Y})$$

is a conjugate for X & Y, where

[The proof that

$$\begin{bmatrix} -(id_{X \otimes Y} \otimes r_{X \otimes Y}^{*}) \circ (\overline{r}_{X \otimes Y} \otimes id_{X \otimes Y}) = id_{X \otimes Y} \\ (id_{\overline{Y} \otimes \overline{X}} \otimes \overline{r}_{X \otimes Y}^{*}) \circ (r_{X \otimes Y} \otimes id_{\overline{Y} \otimes \overline{X}}) = id_{\overline{Y} \otimes \overline{X}} \end{bmatrix}$$

will be left to the reader but we shall provide the verification that

$$\bar{r}_{X \boxtimes Y} = \bar{r}_{X \boxtimes X, X \boxtimes Y} \circ r_{X \boxtimes Y}$$

Thus write

$$\overset{\mathsf{T}}{\nabla} \mathbf{\underline{\omega}} \, \overline{\mathbf{x}}_{\mathsf{X}} \mathbf{\underline{\omega}} \, \mathbf{\underline{Y}} \overset{\circ}{\nabla} \mathbf{\underline{r}} \mathbf{\underline{X}} \, \mathbf{\underline{\omega}} \, \mathbf{\underline{Y}}$$

$$= \mathsf{T} \qquad \circ \mathsf{T} \qquad \circ \mathsf{r}_{\mathsf{X},\mathsf{X} \boxtimes \mathsf{Y} \boxtimes \bar{\mathsf{Y}}} \circ \mathsf{r}_{\mathsf{X} \boxtimes \mathsf{X} \boxtimes \mathsf{Y}} \quad (cf. 4.3).$$

Then

Therefore

$$= \operatorname{id}_{X} \, \, \overline{\operatorname{r}}_{Y} \, \, \underline{\otimes} \, \operatorname{id}_{\overline{X}} \, \circ \, \, \overline{\operatorname{r}}_{X} = \, \overline{\operatorname{r}}_{X \, \, \underline{\otimes} \, \, Y}.]$$

For all $X,Y \in Ob \ \underline{C}$, the map

$$Mor(X,Y) \rightarrow Mor(\overline{Y},\overline{X})$$

that sends f to 'f is a linear bijection.

N.B. Here, as will be recalled from §5,

$$v_f = r_Y^* \otimes id_{\overline{X}} \circ id_{\overline{X}} \otimes f \otimes id_{\overline{X}} \circ id_{\overline{X}} \otimes \overline{r}_{X}.$$

Now put

$$f^+ = (^{\vee}f)*,$$

thus

$$f^+ = id_{\overline{Y}} \otimes \overline{r}_{X}^* \circ id_{\overline{Y}} \otimes f^* \otimes id_{\overline{X}} \circ r_{Y} \otimes id_{\overline{X}}$$

and

$$f^+ \in Mor(\bar{X}, \bar{Y})$$
.

Properties:

1.
$$id_X^+ = id_{\overline{x}}$$
;

2.
$$(f^+)^* = (f^*)^+$$
;

3.
$$(f \circ g)^+ = f^+ \circ g^+$$
.

9.8 LEMMA Given $f \in Mor(X,Y)$, we have

$$f^+ \otimes id_X \circ r_X = id_{\overline{Y}} \otimes f^* \circ r_Y$$

PROOF Start with the LHS and write

$$f^{\dagger} \otimes \operatorname{id}_{X} \circ r_{X}$$

$$= (\operatorname{id}_{\overline{Y}} \otimes \overline{r}_{X}^{*} \circ \operatorname{id}_{\overline{Y}} \otimes f^{*} \otimes \operatorname{id}_{\overline{X}} \circ r_{Y} \otimes \operatorname{id}_{\overline{X}}) \otimes \operatorname{id}_{X} \circ r_{X}$$

$$= (\operatorname{id}_{\overline{Y}} \otimes \overline{r}_{X}^{*} \circ \operatorname{id}_{\overline{Y}} \otimes f^{*} \otimes \operatorname{id}_{\overline{X}} \circ r_{Y} \otimes \operatorname{id}_{\overline{X}}) \otimes (\operatorname{id}_{X} \circ \operatorname{id}_{X} \circ \operatorname{id}_{X}) \circ r_{X}$$

$$= (\operatorname{id}_{\overline{Y}} \otimes \overline{r}_{X}^{*} \otimes \operatorname{id}_{X} \circ \operatorname{id}_{\overline{Y}} \otimes f^{*} \otimes \operatorname{id}_{\overline{X}} \circ r_{Y} \otimes \operatorname{id}_{X} \circ r_{X} \circ \operatorname{id}_{X} \circ r_{X}.$$

But

$$r_{Y} \overset{\text{did}}{\overline{X}} \overset{\text{did}}{X} \circ r_{X}$$

$$= r_{Y} \overset{\text{did}}{\overline{X}} \overset{\text{did}}{X} \circ id_{e} \overset{\text{did}}{X} r_{X}$$

$$= id_{\overline{Y}} \overset{\text{did}}{X} \overset{\text{did}}{X} r_{X} \circ r_{Y} \overset{\text{did}}{R} r_{X}$$

$$= id_{\overline{Y}} \overset{\text{did}}{X} r_{X} \circ r_{Y} \overset{\text{did}}{R} r_{X} \circ r_{Y}$$

=>

=>

$$f^{+} \otimes id_{X} \circ r_{X}$$

$$= id_{\overline{Y}} \otimes \overline{r}_{X}^{*} \otimes id_{X} \circ id_{\overline{Y}} \otimes id_{X} \otimes r_{X} \circ id_{\overline{Y}} \otimes f^{*} \circ r_{Y}$$

$$= id_{\overline{Y}} \otimes (\overline{r}_{X}^{*} \otimes id_{X} \circ id_{X} \otimes r_{X} \circ f^{*}) \circ r_{Y}$$

$$= \operatorname{id}_{\overline{Y}} \otimes \operatorname{id}_{X} \circ f^{*} \circ r_{Y}$$

$$= \operatorname{id}_{\overline{Y}} \otimes f^{*} \circ r_{Y}.$$

9.9 REMARK Suppose that $T \in Mor(\bar{X}, \bar{Y})$ satisfies the equation

$$T \otimes id_X \circ r_X = id_{\overline{Y}} \otimes f^* \circ r_Y$$

Then

$$T = f^+$$
.

Proof:

$$f^{+} = id_{\overline{Y}} \otimes \overline{r_{X}^{*}} \circ id_{\overline{Y}} \otimes f^{*} \otimes id_{\overline{X}} \circ r_{Y} \otimes id_{\overline{X}}$$

$$= id_{\overline{Y}} \otimes \overline{r_{X}^{*}} \circ id_{\overline{Y}} \otimes f^{*} \circ r_{Y} \otimes id_{\overline{X}} \circ id_{\overline{X}}$$

$$= id_{\overline{Y}} \otimes \overline{r_{X}^{*}} \circ T \otimes id_{X} \circ r_{X} \otimes id_{\overline{X}}.$$

On the other hand,

$$T = T \circ id_{\overline{X}}$$

$$= T \circ id_{\overline{X}} \otimes \overline{r_{X}} \circ r_{X} \otimes id_{\overline{X}}$$

$$= T \otimes id_{e} \circ id_{\overline{X}} \otimes \overline{r_{X}} \circ r_{X} \otimes id_{\overline{X}}$$

$$= id_{e} \circ id_{\overline{X}} \otimes \overline{r_{X}} \circ r_{X} \otimes id_{\overline{X}}$$

$$= id_{\overline{Y}} \otimes \overline{r_{X}} \circ T \otimes id_{X} \otimes \overline{x} \circ r_{X} \otimes id_{\overline{X}}$$

$$= id_{\overline{Y}} \otimes \overline{r_{X}} \circ T \otimes id_{X} \otimes id_{\overline{X}} \circ r_{X} \otimes id_{\overline{X}}$$

$$= id_{\overline{Y}} \otimes \overline{r_{X}} \circ T \otimes id_{X} \otimes id_{X} \circ r_{X} \otimes id_{\overline{X}} \circ id_{\overline{X}}$$

$$= id_{\overline{Y}} \otimes \overline{r_{X}} \circ T \otimes id_{X} \circ r_{X} \otimes id_{\overline{X}} \circ id_{\overline{X}}$$

$$=\operatorname{id}_{\overline{Y}} \, \boxtimes \, \overline{r}_{X}^{\star} \, \circ \, T \, \boxtimes \, \operatorname{id}_{X} \, \circ \, r_{X} \, \boxtimes \, \operatorname{id}_{\overline{X}}.$$

[Note: It is thus a corollary that if

$$id_{\overline{Y}} \otimes f^* \circ r_{Y} = 0,$$

then $f^+ = 0$, so

$$(f^{V})^{*} = 0 \Rightarrow (f^{V})^{**} = 0 \Rightarrow f^{V} = 0 \Rightarrow f = 0.$$

9.10 SCHOLIUM f^+ is the unique element of Mor $(\overline{X},\overline{Y})$ such that

$$f^+ \otimes id_X \circ r_X = id_{\overline{V}} \otimes f^* \circ r_Y$$

[Note: $^{\vee}$ f is the unique element of Mor(\overline{Y} , \overline{X}) such that

$$id_{\underline{Y}} \otimes {}^{\vee}f \circ \overline{r}_{\underline{Y}} = f \otimes id_{\overline{\underline{X}}} \circ \overline{r}_{\underline{X'}}$$

so f^+ is the unique element of $Mor(\bar{X},\bar{Y})$ such that

$$\bar{r}_{Y}^{*} \circ id_{Y} \otimes f^{+} = \bar{r}_{X}^{*} \circ f^{*} \otimes id_{\bar{X}}.$$

9.11 LEMMA Suppose that

is symmetric and unitary. Given $X \in Ob C$, put

Then the triple $(F\overline{X},r_{FX},\overline{r}_{FX})$ is a conjugate for FX.

PROOF What we know is that

$$\begin{bmatrix} (\operatorname{id}_X \otimes r_X^*) & \circ & (\overline{r}_X \otimes \operatorname{id}_X) & = \operatorname{id}_X \\ (\operatorname{id}_X \otimes \overline{r}_X^*) & \circ & (r_X \otimes \operatorname{id}_X) & = \operatorname{id}_X \\ \end{bmatrix}$$

hence

$$\begin{bmatrix}
F(id_X \otimes r_X^*) & \circ F(\overline{r}_X \otimes id_X) = id_{FX} \\
F(id_X \otimes \overline{r}_X^*) & \circ F(r_X \otimes id_X) = id_{\overline{X}}, \\
\end{bmatrix}$$

and what we want to prove is that

$$(id_{FX} \overset{\underline{\otimes}}{=} r_{FX}^{\star}) \circ (\overline{r}_{FX} \overset{\underline{\otimes}}{=} id_{FX}) = id_{FX}$$

$$(id_{F\overline{X}} \overset{\underline{\otimes}}{=} \overline{r}_{FX}^{\star}) \circ (r_{FX} \overset{\underline{\otimes}}{=} id_{F\overline{X}}) = id_{F\overline{X}}.$$

The LHS of the first of these is the composition

$$\operatorname{id}_{\operatorname{FX}} \overset{\underline{\omega}}{=} \xi^{-1} \circ \operatorname{Fr}_{X}^{*} \circ \overset{\Xi}{\times}_{,X}$$

$$\circ \tau \circ (\Xi_{X})^{-1} \circ \operatorname{Fr}_{X} \circ \xi \overset{\underline{\omega}}{=} \operatorname{id}_{\operatorname{FX}},$$

F being unitary. Write

Taking into account the commutative diagrams

$$Fe \ \underline{\underline{\omega}} \ FX \longrightarrow F(e \ \underline{\omega} \ X)$$

$$Fr_{X} \ \underline{\underline{\omega}} \ id_{FX} \downarrow \qquad \qquad \downarrow F(r_{X} \ \underline{\omega} \ id_{X})$$

$$F(\overline{X} \ \underline{\omega} \ X) \ \underline{\underline{\omega}} \ FX \longrightarrow F(\overline{X} \ \underline{\omega} \ X \ \underline{\omega} \ X),$$

we have

$$\begin{aligned} \operatorname{Fr}_{X} & \overset{\boxtimes}{=} \operatorname{id}_{FX} \circ \xi \overset{\boxtimes}{=} \operatorname{id}_{FX} \\ & = (\operatorname{\Xi}_{\overline{X}} \otimes \operatorname{X}, \operatorname{X})^{-1} \circ \operatorname{F}(\operatorname{r}_{X} \otimes \operatorname{id}_{X}) \circ \operatorname{\Xi}_{e,X} \circ \xi \overset{\boxtimes}{=} \operatorname{id}_{FX} \\ & = (\operatorname{\Xi}_{\overline{X}} \otimes \operatorname{X}, \operatorname{X})^{-1} \circ \operatorname{F}(\operatorname{r}_{X} \otimes \operatorname{id}_{X}) \circ \operatorname{id}_{FX}. \end{aligned}$$

This leaves

$$\tau_{F\overline{X},FX} \overset{\underline{\omega}}{=} id_{FX} \circ (\Xi_{\overline{X},X})^{-1} \overset{\underline{\omega}}{=} id_{FX} \circ (\Xi_{\overline{X}} \overset{\underline{\omega}}{=} X,X)^{-1} \circ F(r_{X} \overset{\underline{\omega}}{=} id_{X}) \circ id_{FX}.$$

Next

$$\begin{split} \mathbf{F}(\overline{\mathbf{r}}_{\mathbf{X}} & \otimes \mathrm{id}_{\mathbf{X}}) &= \mathbf{F}(\mathbf{\tau}_{\mathbf{X},\mathbf{X}} \circ \mathbf{r}_{\mathbf{X}} \otimes \mathrm{id}_{\mathbf{X}}) \\ &= \mathbf{F}(\mathbf{\tau}_{\mathbf{X},\mathbf{X}} \otimes \mathrm{id}_{\mathbf{X}}) \circ \mathbf{F}(\mathbf{r}_{\mathbf{X}} \otimes \mathrm{id}_{\mathbf{X}}) \,. \end{split}$$

Since F is symmetric, there is a commutative diagram

Here "top" is the composition

$$F\overline{X} \overset{\underline{\otimes}}{\longrightarrow} FX \overset{\underline{\otimes}}{\longrightarrow} F(\overline{X} \overset{\underline{\otimes}}{\times} X) \overset{\underline{\otimes}}{\longrightarrow} FX$$

$$\xrightarrow{\overline{X} \overset{\underline{\otimes}}{\times} X, X} \longrightarrow F(\overline{X} \overset{\underline{\otimes}}{\times} X \overset{\underline{\otimes}}{\times} X)$$

and "bttm" is the composition

$$FX \overset{\underline{\otimes}}{\longrightarrow} FX \xrightarrow{\underline{\otimes}} FX \xrightarrow{\underline$$

Therefore

=>

.

$$= (\Xi)^{-1} \underline{\otimes} \operatorname{id}_{FX} \circ (\Xi)^{-1} \circ F(\tau \underline{\otimes} \operatorname{id}_{X})$$

$$X, \overline{X} \underline{\otimes} \overline{X}, X \underline{X}, X \underline{X}, X$$

=>

Analogously,

$$id_{FX} \overset{\mathbf{Q}}{=} \xi^{-1} \circ Fr_{X}^{*} \circ \Xi_{\overline{X},X}$$

$$= id_{FX} \circ F(id_{X} \overset{\mathbf{Q}}{=} r_{X}^{*}) \circ \Xi_{X,\overline{X} \overset{\mathbf{Q}}{=} X} \circ id_{FX} \overset{\mathbf{Q}}{=} \Xi_{\overline{X},X}.$$

So, in summary,

$$(\mathrm{id}_{\mathrm{FX}} \ \underline{\otimes} \ r_{\mathrm{FX}}^{*}) \circ (\overline{r}_{\mathrm{FX}} \ \underline{\otimes} \ \mathrm{id}_{\mathrm{FX}})$$

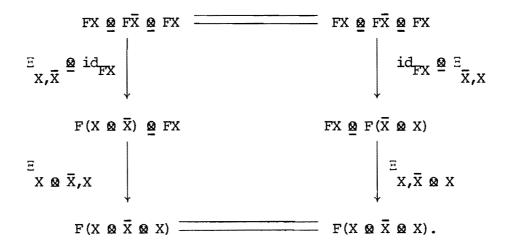
$$= \mathrm{id}_{\mathrm{FX}} \circ \mathrm{F}(\mathrm{id}_{\mathrm{X}} \ \underline{\otimes} \ r_{\mathrm{X}}^{*})$$

$$\circ \ \underline{\mathrm{E}}_{\mathrm{X}, \overline{\mathrm{X}}} \otimes \mathrm{X} \circ \mathrm{id}_{\mathrm{FX}} \ \underline{\otimes} \ \underline{\mathrm{E}}_{\mathrm{X}, \mathrm{X}} \circ (\underline{\mathrm{E}}_{\mathrm{X}})^{-1} \ \underline{\otimes} \ \mathrm{id}_{\mathrm{FX}} \circ (\underline{\mathrm{E}}_{\mathrm{X}})^{-1}$$

$$\circ \ \mathrm{F}(\overline{r}_{\mathrm{X}} \ \underline{\otimes} \ \mathrm{id}_{\mathrm{X}}) \circ \mathrm{id}_{\mathrm{FX}},$$

thus to finish, it need only be shown that

This, however, follows from the commutative diagram



9.12 REMARK We have

$$\bar{r}_{FX} = (\bar{z}_{X,\bar{X}})^{-1} \circ F\bar{r}_{X} \circ \xi.$$

In fact, the RHS equals

$$(\Xi_{X,\overline{X}})^{-1} \circ F_{\overline{X},X} \circ F_{X} \circ \xi$$

and there is a commutative diagram

§10. TANNAKIAN CATEGORIES

Let <u>C</u> be a symmetric strict monoidal *-category which is essentially small -then <u>C</u> is said to be tannakian if the following conditions are met:

$$\underline{\underline{T_1}}$$
: $\forall X,Y \in Ob \underline{C}$,

 $\dim Mor(X,Y) < \infty$.

 \mathbf{T}_2 : $\underline{\mathbf{C}}$ has subobjects, direct sums, and conjugates.

 T_3 : \underline{C} has a zero object.

T₄: e is irreducible.

- 10.1 REMARK A tannakian category is necessarily semisimple, hence is abelian.
- 10.2 EXAMPLE Let CPTGRP be the category whose objects are the compact Hausdorff topological groups (in brief, the "compact groups") and whose morphisms are the continuous homomorphisms. Given an object G in this category, let $\underline{\text{Rep}}$ G be the category whose objects are the finite dimensional continuous unitary representations of G and whose morphisms are the intertwining operators then $\underline{\text{Rep}}$ G is tannakian (define r and $\overline{\text{r}}$ by

where $\{e_{\underline{i}}\}$ \in ${\it H}$ is an orthonormal basis for the representation space and $\{\bar{e}_{\underline{i}}\}$ \in $\bar{{\it H}}$ is

its conjugate). In particular: FDHILB is tannakian (take $G = \{*\}$).

[Note: The construct Rep G is amnestic and transportable, so we can and will assume that its monoidal structure is strict (cf. 3.12).]

10.3 RAPPEL An additive functor $F:A \rightarrow B$ between abelian categories A and B is exact if it preserves finite limits and finite colimits.

Accordingly, since a tannakian category is not only abelian but also autonomous, $\forall \ X \in Ob \ C, \ the \ functors$

are exact.

If C is tannakian, then e is irreducible and

has the following properties.

- 1. dim $X = \dim \overline{X}$.
- 2. $\dim(X \otimes Y) = (\dim X)(\dim Y)$.
- 3. $\dim(X \oplus Y) = \dim X + \dim Y$.
- 4. $\dim e = 1$, $\dim 0 = 0$.
- 10.4 <u>LEMMA</u> If X is not a zero object, then dim X (= $r_X^* \circ r_X$) ≥ 1 .

<u>PROOF</u> First, from the positivity of the involution, dim X > 0. But X \boxtimes \overline{X} contains e as a direct summand, thus

$$(\dim X)^2 \ge 1 \Longrightarrow \dim X \ge 1.$$

[Note: If dim X = 1, then X $\otimes \overline{X} \approx e \approx \overline{X} \otimes X$.]

Given $X \neq 0$ in Ob C, define

$$\Pi_n^X: S_n \to \text{Aut } X^{\otimes n}$$

as in 4.5.

N.B. Π_n^X is a homomorphism from $\$_n$ to the unitary group of Mor(X,X).

Put

$$x^{\Omega 0} = e$$
,
$$\begin{bmatrix} - & \text{Sym}_0^X = id_e \\ & & \\ - & \text{Alt}_0^X = id_e \end{bmatrix}$$

and for $n \in N$, put

$$\operatorname{Sym}_{n}^{X} = \frac{1}{n!} \sum_{\sigma \in \mathcal{S}_{n}} \Pi_{n}^{X}(\sigma)$$

$$\operatorname{Alt}_{n}^{X} = \frac{1}{n!} \sum_{\sigma \in \mathcal{S}_{n}} (\operatorname{sgn} \sigma) \Pi_{n}^{X}(\sigma).$$

Then

are projections.

10.5 LEMMA We have

$$tr_{X} (Alt_{n}^{X})$$

$$= \frac{1}{n!} (\dim X) (\dim X - 1) \dots (\dim X - n + 1).$$

PROOF The key preliminary is the observation that

$$\operatorname{tr}_{X}^{\Omega n}(\Pi_{n}^{X}(\sigma)) = (\dim X)^{\#\sigma},$$

where $\#\sigma$ is the number of cycles into which σ decomposes, thus

$$\operatorname{tr}_{X}^{\otimes n}(\operatorname{Alt}_{n}^{X}) = \frac{1}{n!} \sum_{\sigma \in \mathcal{S}_{n}} (\operatorname{sgn} \sigma) (\dim X)^{\#\sigma}.$$

But for every complex number z,

$$\sum_{\sigma \in \mathcal{G}_n} (\operatorname{sgn} \sigma) z^{\#\sigma} = z(z-1) \dots (z-n+1).$$

10.6 THEOREM \forall nonzero X in Ob C,

$$\dim X \in N$$
.

 $\underline{PROOF} \text{ Let } A_n(X) \text{ be the subobject of } X^{\boxtimes n} \text{ corresponding to } Alt_n^X. \text{ Fix an}$ isometry $f:A_n(X) \to X^{\boxtimes n}$ such that $f \circ f^* = Alt_n^X$ — then

$$tr_{X}^{\otimes n}(Alt_{n}^{X})$$

$$= tr_{X}^{\otimes n}(f \circ f^{*})$$

$$= tr_{A_{n}(X)}(f^{*} \circ f) \quad (cf. 6.3)$$

=
$$\operatorname{tr}_{A_{n}(X)}(\operatorname{id}_{A_{n}(X)})$$

= $\operatorname{dim} A_{n}(X) \ge 1$ (cf. 10.4).

On the other hand, thanks to 10.5,

$$\operatorname{tr}_{X^{\boxtimes n}} (\operatorname{Alt}_{n}^{X})$$

is negative for some $n \in N$ unless dim $X \in N$.

10.7 LEMMA Let $d = \dim X$ -- then

$$\dim A_{d}(X) = \operatorname{tr}_{X} \operatorname{Alt}_{d}^{X} = \frac{d!}{d!} = 1.$$

The isomorphism class of $A_d(X)$ is called the <u>determinant</u> of X (written det (X)). Properties:

- 1. $\det(\overline{X}) \approx \overline{\det(X)}$;
- 2. $det(X \oplus Y) \approx det(X) \otimes det(Y)$;
- 3. $\det(X \oplus \overline{X}) \approx e$.

§11. FIBER FUNCTORS

Let $\underline{\mathtt{C}}$ be a tannakian category — then a symmetric embedding functor

is called a fiber functor.

E.g.: Take $\underline{C} = \underline{Rep} G$ (cf. 10.2) — then the forgetful functor

is a fiber functor.

 $\underline{\text{N.B.}}$ It is a nontrivial result that every tannakian category admits a fiber functor (proof omitted).

11.1 REMARK Let

be a fiber functor. Consider

$$A_{F} = \bigoplus_{i \in I_{\underline{C}}} B(FX_{i}),$$

viewed as a subset of $\underline{\text{Nat}}(\mathcal{F},\mathcal{F})$ — then the <u>coinverse</u> is the map $S:A_{\mathcal{F}} \to A_{\mathcal{F}}$ defined by

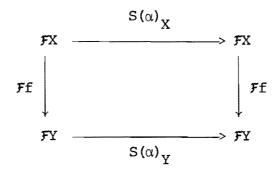
$$\mathtt{S}\left(\alpha\right)_{X} = \mathtt{F}(\mathtt{id}_{X} \ \boxtimes \ r_{X}^{\star}) \ \circ \ \mathtt{id}_{\mathtt{F}X} \ \underline{\boxtimes} \ \alpha_{\overline{X}} \ \underline{\boxtimes} \ \mathtt{id}_{\mathtt{F}X} \ \circ \ \mathtt{F}(\overline{r}_{X} \ \boxtimes \ \mathtt{id}_{X}) \,,$$

matters being slightly imprecise in that the identification

$$F(X \boxtimes \overline{X} \boxtimes X) \approx FX \boxtimes F\overline{X} \boxtimes FX$$

has been suppressed. It is not difficult to see that the equation defining $S(\alpha)_X$ is independent of the choice $(\bar{X}, r_X, \bar{r}_X)$ of a conjugate for X and \forall $f \in Mor(X,Y)$,

the diagram



commutes. Algebraically, S is linear and antimultiplicative. Moreover,

$$S \circ * \circ S \circ * = id_{A_{\mathbf{F}}}$$

hence S is invertible.

[Note: There are various relations among Δ, ϵ, S which, however, need not be detailed. Still, despite appearances, in general $(A_{\mathfrak{f}}, \Delta, \epsilon, S)$ is not a Hopf *-algebra but rather in the jargon of the trade is a "cocommutative discrete algebraic quantum group".]

Write ff(C) for the full subcategory of

whose objects are the fiber functors -- then ff(C) is a groupoid (cf. 5.12).

11.2 THEOREM ff(C) is a transitive groupoid, i.e., if F_1 , F_2 are fiber functors, then F_1 , F_2 are isomorphic.

Definition: Given fiber functors F_1, F_2 , a <u>unitary</u> monoidal natural transformation $\alpha: F_1 \to F_2$ is a monoidal natural transformation such that $\forall \ X \in Ob \ \underline{C}$,

$$\alpha_{x}:\mathcal{F}_{1}X \rightarrow \mathcal{F}_{2}X$$

is unitary.

Write $ff^*(\underline{C})$ for the category whose objects are the fiber functors and whose morphisms are the unitary monoidal natural transformations — then $ff^*(\underline{C})$ is a subcategory of ff(C).

11.3 THEOREM ff*(\underline{C}) is a transitive groupoid, i.e., if F_1 , F_2 are fiber functors, then F_1 , F_2 are unitarily isomorphic.

Obviously,

$$11.3 \Rightarrow 11.2.$$

As for the proof of 11.3, there will be three steps.

Step 1: Construct a commutative unital *-algebra $A(F_1,F_2)$ whose dual space is in a one-to-one correspondence with the natural transformations $F_1 \to F_2$, to wit:

$$Nat(\mathcal{F}_1,\mathcal{F}_2) \iff A(\mathcal{F}_1,\mathcal{F}_2)^*.$$

Step 2: Under this bijection, prove that the monoidal natural transformations correspond to the nonzero multiplicative linear functionals on $A(\mathcal{F}_1,\mathcal{F}_2)$ and the unitary monoidal natural transformations correspond to the *-preserving multiplicative linear functionals on $A(\mathcal{F}_1,\mathcal{F}_2)$.

Step 3: Establish that $A(\mathcal{F}_1,\mathcal{F}_2)$ admits a C*-norm, thus is a pre-C*-algebra. Therefore, since the structure space $\Delta(\bar{A}(\mathcal{F}_1,\mathcal{F}_2))$ of the C*-completion $\bar{A}(\mathcal{F}_1,\mathcal{F}_2)$ of $A(\mathcal{F}_1,\mathcal{F}_2)$ is not empty, it follows that $Mor(\mathcal{F}_1,\mathcal{F}_2)$ is also not empty, from which 11.3.

[Note: Here, of course, Mor is computed in $ff*(\underline{C})$.] To fix notation, bear in mind that there are isomorphisms

$$\begin{bmatrix} \xi^{1} : \underline{e} \rightarrow F_{1} \underline{e} & & & & & \\ & \xi^{2} : \underline{e} \rightarrow F_{2} \underline{e} & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & \\ & & & \\ & & & \\ & & \\ & & & \\ & & \\ & & & \\ & & \\ & & \\ & & & \\$$

subject to the compatibility conditions enumerated in §2.

Let $A_0(\mathcal{F}_1,\mathcal{F}_2)$ be the complex vector space

$$x \in Ob \subseteq Mor(\mathcal{F}_2X, \mathcal{F}_1X).$$

Given $X \in Ob \subseteq A$ and $\phi \in Mor(\mathcal{F}_2X, \mathcal{F}_1X)$, write $[X, \phi]_0$ for the element of $A_0(\mathcal{F}_1, \mathcal{F}_2)$ that is ϕ at X and is zero elsewhere — then $A_0(\mathcal{F}_1, \mathcal{F}_2)$ is simply the complex linear span of the $[X, \phi]_0$. Define a product in $A_0(\mathcal{F}_1, \mathcal{F}_2)$ by stipulating that

$$[X,\phi]_0 \cdot [Y,\psi]_0 = [X \otimes Y,u]_0$$

where u is the composition

11.4 $\underline{\text{LEMMA}}$ $A_0(\mathcal{F}_1, \mathcal{F}_2)$ is associative.

11.5 <u>LEMMA</u> $A_0(\mathcal{F}_1, \mathcal{F}_2)$ is unital.

PROOF Let

$$1_{A_0} = [e, \xi^1 \circ (\xi^2)^{-1}]_0.$$

Then l_{A_0} is the unit. E.g.: Consider

$$[x,\phi]_0 \cdot [e,\xi^1 \circ (\xi^2)^{-1}]_0 = [x,u]_0$$

the claim being that the composite

$$\mathcal{F}_{2} = \mathcal{F}_{2} (X \boxtimes e) \xrightarrow{(\Xi_{X,e}^{2})^{-1}} \qquad \qquad \mathcal{F}_{2} \times \underline{\otimes} \mathcal{F}_{2} e$$

$$\xrightarrow{\Phi \boxtimes (\xi^{1} \circ (\xi^{2})^{-1})} \qquad \qquad \mathcal{F}_{1} \times \underline{\otimes} \mathcal{F}_{1} e$$

$$\xrightarrow{\Xi_{X,e}^{1}} \qquad \qquad \mathcal{F}_{1} (X \boxtimes e) = \mathcal{F}_{1} \times e$$

reduces to ϕ itself. To see this, recall that the composition

$$\begin{split} \mathfrak{F}_1^{} x &= \mathfrak{F}_1^{} x \ \underline{\otimes} \ \underline{e} & \xrightarrow{\mathrm{id}_{\mathfrak{F}_1^{} X}} \ \underline{\otimes} \ \xi^1 \\ & \\ \mathcal{F}_1^{} x &= \mathfrak{F}_1^{} x \ \underline{\otimes} \ \underline{e} & \xrightarrow{\Xi^1_{X,e}} \\ \end{split} \rightarrow \begin{split} \mathfrak{F}_1^{} x &= \mathfrak{F}_1^{} x \ \underline{\otimes} \ \underline{e} & \xrightarrow{\Xi^1_{X,e}} \end{split} \rightarrow \mathcal{F}_1^{} (x \ \underline{\otimes} \ \underline{e}) = \mathcal{F}_1^{} x \end{split}$$

is the identity morphism of $\mathcal{F}_1^{}\mathbf{X}$ and the composition

$$f_{2} = f_{2} \times \underline{\underline{Q}} = \underbrace{f_{2}^{2} \times \underline{\underline{Q}}}_{\xi_{2}} = \underbrace{f_{$$

is the identity morphism of \mathcal{F}_2X . Now write

$$\begin{split} & \Xi_{X,e}^{1} \circ \phi \ \underline{\mathfrak{Q}} \ (\xi^{1} \circ (\xi^{2})^{-1}) \circ (\Xi_{X,e}^{2})^{-1} \\ &= \operatorname{id}_{\mathcal{F}_{1}X} \circ \operatorname{id}_{\mathcal{F}_{1}X} \ \underline{\mathfrak{Q}} \ (\xi^{1})^{-1} \circ \phi \ \underline{\mathfrak{Q}} \ (\xi^{1} \circ (\xi^{2})^{-1}) \circ \operatorname{id}_{\mathcal{F}_{2}X} \ \underline{\mathfrak{Q}} \ \xi^{2} \circ \operatorname{id}_{\mathcal{F}_{2}X} \end{split}$$

$$\begin{split} &=\mathrm{id}_{F_1X}\circ(\mathrm{id}_{F_1X}\circ\phi\circ\mathrm{id}_{F_2X}\ \underline{\otimes}\ (\xi^{-1})^{-1}\circ(\xi^1\circ(\xi^2)^{-1}\circ\xi^2)\circ\mathrm{id}_{F_2X}\\ &=\mathrm{id}_{F_1X}\circ\phi\ \underline{\otimes}\ \mathrm{id}_{E_2X}=\mathrm{id}_{F_1X}\circ\phi\circ\mathrm{id}_{F_2X}=\phi. \end{split}$$

Let $I_0(\mathcal{F}_1,\mathcal{F}_2)$ be the complex linear span of the

$$[X,a \circ F_2f]_0 - [Y,F_1f \circ a]_0$$

where

$$\mathtt{f} \in \mathtt{Mor}(\mathtt{X},\mathtt{Y})\,,\ \mathtt{a} \in \mathtt{Mor}(\mathtt{F}_{2}\mathtt{Y},\mathtt{F}_{1}\mathtt{X})\,.$$

Then $I_0(\mathcal{F}_1,\mathcal{F}_2)$ is an ideal in $A_0(\mathcal{F}_1,\mathcal{F}_2)$.

Denote by $A(\mathcal{F}_1, \mathcal{F}_2)$ the quotient algebra

$$A_0(F_1,F_2)/I_0(F_1,F_2)$$
,

let

$$\operatorname{pr}: A_0(\mathcal{F}_1, \mathcal{F}_2) \to A(\mathcal{F}_1, \mathcal{F}_2)$$

be the projection, and put

$$[X, \phi] = pr[X, \phi]_0.$$

11.6 EXAMPLE Let $f:X \to X$ be an isomorphism — then

$$[X,\phi] = [X,\mathcal{F}_1 f \circ \mathcal{F}_1 f^{-1} \circ \phi]$$
$$= [X,\mathcal{F}_1 f^{-1} \circ \phi \circ \mathcal{F}_2 f].$$

11.7 EXAMPLE Let

$$\Phi \in Mor(\mathcal{F}_2(\overline{X} \boxtimes X), \mathcal{F}_1(\overline{X} \boxtimes X))$$

Then

$$[\overline{X} \boxtimes X, \mathcal{F}_{1}r_{X} \circ \mathcal{F}_{1}r_{X}^{*} \circ \Phi]$$

$$= [e, \mathcal{F}_{1}r_{X}^{*} \circ \Phi \circ \mathcal{F}_{2}r_{X}].$$

[Note: We also have

$$[\overline{X} \otimes X, \mathcal{F}_{1}(r_{X} \circ r_{X}^{*}) \circ \Phi]$$

$$= [\overline{X} \otimes X, \Phi \circ \mathcal{F}_{2}(r_{X} \circ r_{X}^{*})].]$$

11.8 REMARK Every $A \in A(\mathcal{F}_1, \mathcal{F}_2)$ can be written as $[X, \phi]$ for a suitable choice of X and ϕ . Thus suppose that $A = \sum\limits_{i} [X_i, \phi_i]$, put $X = \emptyset$ X_i , and choose isometries $v_i : X_i \to X$ such that $\sum\limits_{i} v_i \circ v_i^* = \mathrm{id}_X$ — then

$$a_{i} = \phi_{i} \circ \mathcal{F}_{2}v_{i}^{*} \in Mor(\mathcal{F}_{2}X, \mathcal{F}_{1}X_{i})$$

=>

$$A = \sum_{i} [X_{i}, \phi_{i}]$$

$$= \sum_{i} [X_{i}, \phi_{i} \circ id_{\mathcal{F}_{2}X_{i}}]$$

$$= \sum_{i} [X_{i}, \phi_{i} \circ \mathcal{F}_{2}(v_{i}^{*} \circ v_{i})]$$

$$= \sum_{i} [X_{i}, \phi_{i} \circ \mathcal{F}_{2}v_{i}^{*} \circ \mathcal{F}_{2}v_{i}]$$

$$= \sum_{i} [X_{i}, a_{i} \circ \mathcal{F}_{2}v_{i}]$$

$$= \sum_{i} [X, \mathcal{F}_{1}v_{i} \circ a_{i}]$$

$$= \sum_{i} [X, \mathcal{F}_{1} v_{i} \circ \phi_{i} \circ \mathcal{F}_{2} v_{i}^{*}]$$

$$= [X, \sum_{i} \mathcal{F}_{1} v_{i} \circ \phi_{i} \circ \mathcal{F}_{2} v_{i}^{*}]$$

$$= [X, \phi],$$

where

$$\phi = \sum_{i} \mathcal{F}_{1} v_{i} \circ \phi_{i} \circ \mathcal{F}_{2} v_{i}^{*} \in Mor(\mathcal{F}_{2} X, \mathcal{F}_{1} X).$$

11.9 <u>LEMMA</u> $A(\mathcal{F}_1, \mathcal{F}_2)$ is commutative.

PROOF Let

$$\begin{bmatrix} \begin{bmatrix} [X,\phi]_0 & (\phi:\mathcal{F}_2X \to \mathcal{F}_1X) \\ [Y,\psi]_0 & (\psi:\mathcal{F}_2Y \to \mathcal{F}_1Y) \end{bmatrix}$$

be elements of $A_0(\mathfrak{F}_1,\mathfrak{F}_2)$ — then

$$\left[\mathbf{X}, \phi \right]_0 \cdot \left[\mathbf{Y}, \psi \right]_0 = \left[\mathbf{X} \ \mathbf{\underline{Q}} \ \mathbf{Y}, \Xi_{\mathbf{X}, \mathbf{Y}}^1 \circ \phi \ \mathbf{\underline{\underline{Q}}} \ \psi \circ \left(\Xi_{\mathbf{X}, \mathbf{Y}}^2 \right)^{-1} \right].$$

On the other hand,

$$[Y,\psi]_0 \cdot [X,\phi]_0 = [Y \otimes X,\Xi^1_{Y,X} \circ \psi \underline{\otimes} \phi \circ (\Xi^2_{Y,X})^{-1}]$$

and there is a commutative diagram

Thus

$$\Xi_{\mathbf{Y},\mathbf{X}}^{1} \circ \psi \underline{\otimes} \phi \circ (\Xi_{\mathbf{Y},\mathbf{X}}^{2})^{-1}$$

$$= \Xi_{\mathbf{Y},\mathbf{X}}^{1} \circ \tau_{\mathcal{F}_{1}\mathbf{X},\mathcal{F}_{1}\mathbf{Y}} \circ \phi \underline{\otimes} \psi \circ \tau_{\mathcal{F}_{2}\mathbf{Y},\mathcal{F}_{2}\mathbf{X}} \circ (\Xi_{\mathbf{Y},\mathbf{X}}^{2})^{-1}.$$

But there are also commutative diagrams

$$\begin{array}{c|c} \mathcal{F}_{1}^{X} & \underline{\otimes} & \mathcal{F}_{1}^{Y} & & & & \\ & & & & & \\ \uparrow_{f_{1}^{X}, \mathcal{F}_{1}^{Y}} & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ &$$

and

Thus

$$\Xi_{Y,X}^{1} \circ \tau_{\mathcal{F}_{1}X,\mathcal{F}_{1}Y} \circ \phi \underline{\otimes} \psi \circ \tau_{\mathcal{F}_{2}Y,\mathcal{F}_{2}X} \circ (\Xi_{Y,X}^{2})^{-1}$$

$$= \mathcal{F}_{1}\tau_{X,Y} \circ \Xi_{X,Y}^{1} \circ \phi \underline{\otimes} \psi \circ (\Xi_{X,Y}^{2})^{-1} \circ \mathcal{F}_{2}\tau_{Y,X}^{2}$$

Let $f = \tau_{X,Y}$ and put

$$a = \Xi_{X,Y}^1 \circ \phi \underline{\omega} \psi \circ (\Xi_{X,Y}^2)^{-1} \circ \mathcal{F}_{2^T Y,X}.$$

Then

$$f \in Mor(X \otimes Y, Y \otimes X)$$

and

$$\mathtt{a} \in \mathtt{Mor}(\mathtt{F}_2(\mathtt{Y} \ \mathtt{Q} \ \mathtt{X})\,,\mathtt{F}_1(\mathtt{X} \ \mathtt{Q} \ \mathtt{Y})\,)\,.$$

Moreover

$$[Y,\psi]_0 \cdot [X,\phi]_0 = [Y \otimes X,F_1f \circ a]_0.$$

Meanwhile

$$F_{2}^{\mathsf{T}}_{\mathsf{Y},\mathsf{X}} \circ F_{2}^{\mathsf{T}}_{\mathsf{X},\mathsf{Y}}$$

$$= F_{2}(\mathsf{T}_{\mathsf{Y},\mathsf{X}} \circ \mathsf{T}_{\mathsf{X},\mathsf{Y}})$$

$$= F_{2}(\mathrm{id}_{\mathsf{X}} \otimes \mathsf{Y})$$

$$= \mathrm{id}_{\mathsf{F}_{2}}(\mathsf{X} \otimes \mathsf{Y})'$$

SO

$$[\mathbf{X}, \phi]_0 \cdot [\mathbf{Y}, \psi]_0 = [\mathbf{X} \mathbf{Q} \mathbf{Y}, \mathbf{a} \circ \mathbf{F}_2 \mathbf{f}]_0$$

Therefore

$$\left[\mathbf{x},\boldsymbol{\varphi}\right]_{0} \; \cdot \; \left[\mathbf{y},\boldsymbol{\psi}\right]_{0} \; - \; \left[\mathbf{y},\boldsymbol{\psi}\right]_{0} \; \cdot \; \left[\mathbf{x},\boldsymbol{\varphi}\right]_{0} \in \mathcal{I}_{0}(\mathcal{F}_{1},\mathcal{F}_{2})$$

=>

$$[A_0(F_1,F_2),A_0(F_1,F_2)] \subset I_0(F_1,F_2).$$

And this implies that $A(\mathcal{F}_1,\mathcal{F}_2)$ is commutative.

Given $[X,\phi]_0$, choose a conjugate (\bar{X},r_X,\bar{r}_X) for X and let $[X,\phi]_0^\star = [\bar{X},\bar{\phi}]_0^\star,$

where $\bar{\phi}$ is the composition

N.B. We have

$$(\Xi_{\overline{X},X}^{1})^{-1} \overset{\underline{\omega}}{=} \overset{id}{F_{2}\overline{X}} \circ F_{1}r_{X} \overset{\underline{\omega}}{=} \overset{id}{F_{2}\overline{X}} \circ \xi^{1} \overset{\underline{\omega}}{=} \overset{id}{F_{2}\overline{X}}$$

$$= (\Xi_{\overline{X},X}^{1})^{-1} \circ F_{1}r_{X} \circ \xi^{1} \overset{\underline{\omega}}{=} \overset{id}{F_{2}\overline{X}}$$

$$= r_{F_{1}X} \overset{\underline{\omega}}{=} \overset{id}{F_{2}\overline{X}} \quad (cf. 9.10)$$

and

$$\operatorname{id}_{F_{1}\bar{X}} \overset{\underline{\omega}}{=} (\xi^{2})^{-1} \circ \operatorname{id}_{F_{1}\bar{X}} \overset{\underline{\omega}}{=} F_{2}\bar{x}^{*} \circ \operatorname{id}_{F_{1}\bar{X}} \overset{\underline{\omega}}{=} \Xi^{2}_{X,\bar{X}}$$

$$= \operatorname{id}_{F_{1}\bar{X}} \overset{\underline{\omega}}{=} (\xi^{2})^{-1} \circ F_{2}\bar{x}^{*} \circ \Xi^{2}_{X,\bar{X}}$$

$$= \operatorname{id}_{F_{1}\bar{X}} \overset{\underline{\omega}}{=} \bar{x}^{*}_{F_{2}X} \quad (\text{cf. 9.11}).$$

Therefore

$$\bar{\phi} = \operatorname{id}_{\mathfrak{F}_1 \overline{X}} \ \underline{\otimes} \ \bar{r}_{\mathfrak{F}_2 X}^{\star} \circ \operatorname{id}_{\mathfrak{F}_1 \overline{X}} \ \underline{\otimes} \ \phi^{\star} \ \underline{\otimes} \ \operatorname{id}_{\mathfrak{F}_2 \overline{X}} \circ r_{\mathfrak{F}_1 X} \ \underline{\otimes} \ \operatorname{id}_{\mathfrak{F}_2 \overline{X}}.$$

[Note: By definition, $\phi \in Mor(\mathcal{F}_2X,\mathcal{F}_1X)$, so $^{\vee}\phi \in Mor(\mathcal{F}_1\bar{X},\mathcal{F}_2\bar{X})$, where, as in §5,

$${}^{\vee}\varphi = \epsilon_{\mathcal{F}_1X} \overset{\underline{\omega}}{=} \overset{\mathrm{id}}{\mathfrak{F}_2\overline{x}} \circ \overset{\mathrm{id}}{=}_{\mathcal{F}_1\overline{x}} \overset{\underline{\omega}}{=} \varphi \overset{\underline{\omega}}{=} \overset{\mathrm{id}}{=}_{\mathcal{F}_2\overline{x}} \circ \overset{\mathrm{id}}{=}_{\mathcal{F}_1\overline{x}} \overset{\underline{\omega}}{=} {}^{\eta}_{\mathcal{F}_2X}$$

or still,

$${}^{\vee} \varphi = r_{\mathcal{F}_{1}X}^{*} \overset{\underline{\omega}}{=} \operatorname{id}_{\mathcal{F}_{2}\overline{X}} \circ \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \overset{\underline{\omega}}{=} \varphi \overset{\underline{\omega}}{=} \operatorname{id}_{\mathcal{F}_{2}\overline{X}} \circ \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \overset{\underline{\omega}}{=} r_{\mathcal{F}_{2}X}.$$

Therefore

$$\bar{\Phi} = (^{\vee} \Phi) *.]$$

Replacing $(\bar{X}, r_{\bar{X}}, \bar{r_{\bar{X}}})$ by $(\bar{X}', r_{\bar{X}}', \bar{r_{\bar{X}}}')$ and using 9.4, one finds that

$$[\overline{x},\overline{\phi}]_0 - [\overline{x}',\overline{\phi}']_0 \in \mathcal{I}_0(\mathcal{F}_1,\mathcal{F}_2)$$
.

Therefore the image of $[X,\phi]_0^*$ in $A(\mathcal{F}_1,\mathcal{F}_2)$ is independent of the choice of a conjugate for X.

11.10 LEMMA $I_0(\mathcal{F}_1, \mathcal{F}_2)$ is *-invariant.

Consequently, $*:A_0(\mathcal{F}_1,\mathcal{F}_2) \to A_0(\mathcal{F}_1,\mathcal{F}_2)$ induces a map $*:A(\mathcal{F}_1,\mathcal{F}_2) \to A(\mathcal{F}_1,\mathcal{F}_2)$.

11.11 <u>LEMMA</u> $A(\mathcal{F}_1, \mathcal{F}_2)$ is a *-algebra.

Summary: $A(F_1,F_2)$ is a commutative unital *-algebra.

Accordingly, to complete Step 1, it remains to construct an isomorphism between $A(F_1,F_2)$ * and $Nat(F_1,F_2)$.

On general grounds,

$$A_0(\mathcal{F}_1,\mathcal{F}_2)^* = \prod_{X \in Ob} \underline{C} \operatorname{Mor}(\mathcal{F}_2^X,\mathcal{F}_1^X)^*.$$

But the pairing

$$Mor(\mathcal{F}_2X, \mathcal{F}_1X) \times Mor(\mathcal{F}_1X, \mathcal{F}_2X) \rightarrow C$$

that sends $\phi \times \psi$ to $tr(\phi \circ \psi)$ is nondegenerate, thus

$$\mathsf{A}_0(\mathsf{F}_1,\mathsf{F}_2) \, * \, \approx \, \mathop{\textstyle \prod}_{\mathsf{X} \in \mathsf{Ob}} \, \, _{\mathsf{C}} \, \, \mathsf{Mor} \, (\mathsf{F}_1 \mathsf{X},\mathsf{F}_2 \mathsf{X}) \, .$$

On the other hand, $\operatorname{Nat}(\mathcal{F}_1,\mathcal{F}_2)$ consists of those elements

$$\alpha \in \mathop{\textstyle \prod}_{X \in Ob} \mathop{\underline{C}} \mathop{\texttt{Mor}} ({\mathcal{F}}_1 x, {\mathcal{F}}_2 x)$$

such that $\forall f \in Mor(X,Y)$,

$$\mathcal{F}_2 f \circ \alpha_X = \alpha_V \circ \mathcal{F}_1 f$$

and the dual of $A(\mathcal{F}_1,\mathcal{F}_2)$ is the subspace of $A_0(\mathcal{F}_1,\mathcal{F}_2)^*$ comprised of those elements that vanish identically on $I_0(\mathcal{F}_1,\mathcal{F}_2)$. To characterize the latter, take an

and suppose that $\forall \ \mathtt{A} \in \mathcal{I}_0(\mathbf{F}_1,\mathbf{F}_2)$,

$$\langle A, \alpha \rangle = 0$$

or still,

$$< [X,a \circ F_2f]_0 - [Y,F_1f \circ a]_0,\alpha > = 0$$

for all

$$\mathtt{f} \in \mathtt{Mor}(\mathtt{X},\mathtt{Y})\,,\ \mathtt{a} \in \mathtt{Mor}(\mathtt{F}_2\mathtt{Y},\mathtt{F}_1\mathtt{X})\,.$$

I.e.:

$$\mathsf{tr}_{\mathcal{F}_1 X}(\mathsf{a} \, \circ \, \mathcal{F}_2 \mathsf{f} \, \circ \, \alpha_X) \, = \, \mathsf{tr}_{\mathcal{F}_1 Y}(\mathcal{F}_1 \mathsf{f} \, \circ \, \mathsf{a} \, \circ \, \alpha_Y) \, .$$

From the nondegeneracy of the trace, it then follows that

$$\mathcal{F}_2 f \circ \alpha_x = \alpha_y \circ \mathcal{F}_1 f$$

implying thereby that

$$\alpha \in \text{Nat}(F_1,F_2)$$
.

11.12 LEMMA Under the bijection

$$Nat(F_1,F_2) \iff A(F_1,F_2)*,$$

the monoidal natural transformations correspond to the nonzero multiplicative linear functionals on $A(\mathcal{F}_1,\mathcal{F}_2)$.

<u>PROOF</u> To say that a linear functional on $A(F_1,F_2)$ corresponding to an $\alpha \in \operatorname{Nat}(F_1,F_2)$ is multiplicative amounts to saying that

$$\langle [X, \phi] \cdot [Y, \psi], \alpha \rangle$$

$$= \langle [X, \phi], \alpha \rangle \cdot \langle [Y, \psi], \alpha \rangle$$

for all

$$[X,\phi],[Y,\psi] \in A(\mathcal{F}_1,\mathcal{F}_2).$$

Since $<--, \alpha>$ is null on $I_0(\mathcal{F}_1, \mathcal{F}_2)$, it suffices to work upstairs, hence explicated we have

$$\begin{split} \operatorname{tr}_{F_{1}(X \boxtimes Y)} & \stackrel{(\Xi_{X,Y}^{1})}{=} \operatorname{color}_{X,Y} \circ \varphi \ \underline{\otimes} \ \psi \circ (\Xi_{X,Y}^{2})^{-1} \circ \alpha_{X \boxtimes Y}) \\ & = \operatorname{tr}_{F_{1}X} (\varphi \circ \alpha_{X}) \operatorname{tr}_{F_{1}Y} (\psi \circ \alpha_{Y}) \\ & = \operatorname{tr}_{F_{1}X \boxtimes F_{1}Y} ((\varphi \circ \alpha_{X}) \ \underline{\otimes} \ (\psi \circ \alpha_{Y})) \\ & = \operatorname{tr}_{F_{1}X \boxtimes F_{1}Y} ((\varphi \otimes \psi \circ \alpha_{X} \ \underline{\otimes} \ \alpha_{Y}). \end{split}$$

Therefore

$$\alpha_{X \otimes Y} = \Xi_{X,Y}^2 \circ \alpha_{X \otimes \alpha_{Y}} \circ (\Xi_{X,Y}^1)^{-1},$$

the condition that α be monoidal.

[Note: Tacitly,

$$<1_{A_0}, \alpha> = 1$$

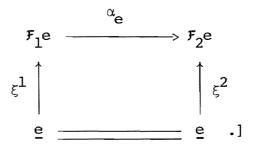
or still,

$$<[e,\xi^1 \circ (\xi^2)^{-1}]_{0},\alpha>=1$$

or still,

$$\operatorname{tr}_{F_1} e^{(\xi^1 \circ (\xi^2)^{-1} \circ \alpha_e)} = 1,$$

from which the commutativity of the diagram



11.13 LEMMA Under the bijection

$$Nat(F_1,F_2) \iff A(F_1,F_2)*,$$

the unitary monoidal natural transformations correspond to the *-preserving non-zero multiplicative linear functionals on $A(\mathcal{F}_1,\mathcal{F}_2)$.

 $\underline{\text{PROOF}}$ Given $[\textbf{X}, \boldsymbol{\varphi}] \, \in \, \textbf{A}(\textbf{\textit{F}}_1, \textbf{\textit{F}}_2) \,,$ the claim is that

$$\langle [X, \phi]^*, \alpha \rangle = \overline{\langle [X, \phi], \alpha \rangle} \ (= \langle [X, \phi], \alpha \rangle^* \ldots)$$

iff $\alpha_X^* = \alpha_X^{-1}$.

From the definitions,

$$\langle [X, \phi], \alpha \rangle = tr_{\mathcal{F}_1} X^{(\phi \circ \alpha_X)}$$

=>

$$\overline{\langle [X, \phi], \alpha \rangle} = \operatorname{tr}_{\mathcal{F}_{1}X}(\phi^{*} \circ \alpha_{X}^{*}).$$

In the other direction,

$$\langle [X, \phi]^*, \alpha \rangle = \langle [\overline{X}, \overline{\phi}], \alpha \rangle$$

$$= \operatorname{tr}_{f_1} (\overline{\phi} \circ \alpha)$$

$$= r^* \circ \operatorname{id}_{f_1} X \stackrel{\underline{\omega}}{=} (\overline{\phi} \circ \alpha) \circ r$$

$$f_1 \overline{X}$$

$$=\bar{r}_{\mathfrak{F}_{1}X}^{\star}\circ\mathrm{id}_{\mathfrak{F}_{1}X}\stackrel{\underline{\omega}}{-}(\bar{\phi}\circ\alpha_{\overline{X}})\circ\bar{r}_{\mathfrak{F}_{1}X}.$$

But

I.e.:

$$\operatorname{tr}_{\mathcal{F}_{1}\bar{\mathbf{X}}}(\bar{\boldsymbol{\varphi}} \circ \alpha_{\bar{\mathbf{X}}}) = \bar{\mathbf{r}}_{\mathcal{F}_{2}\mathbf{X}}^{\star} \circ \boldsymbol{\varphi}^{\star} \overset{\boldsymbol{\underline{\omega}}}{=} \alpha_{\bar{\mathbf{X}}} \circ \bar{\mathbf{r}}_{\mathcal{F}_{1}\mathbf{X}}^{\star}$$

Proceeding, write

$$\begin{split} \vec{r}_{\mathcal{F}_{2}X}^{\star} & \circ \ \varphi^{\star} \ \underline{\otimes} \ \alpha_{\overline{X}} \ \circ \ \vec{r}_{\mathcal{F}_{1}X} \\ & = \ \vec{r}_{\mathcal{F}_{2}X}^{\star} \ \circ \ ((\alpha_{X} \circ \alpha_{X}^{-1} \circ \varphi^{\star}) \ \underline{\otimes} \ \alpha_{\overline{X}} \circ \ \vec{r}_{\mathcal{F}_{1}\overline{X}}) \ \circ \ \vec{r}_{\mathcal{F}_{1}X} \\ & = \ \vec{r}_{\mathcal{F}_{2}X}^{\star} \circ \alpha_{X} \ \underline{\otimes} \ \alpha_{\overline{X}} \circ \alpha_{X}^{-1} \circ \varphi^{\star} \ \underline{\otimes} \ \mathrm{id}_{\mathcal{F}_{1}\overline{X}} \circ \ \vec{r}_{\mathcal{F}_{1}X}. \end{split}$$

We then claim that

$$\bar{r}_{\mathcal{F}_{2}X}^{\star} \circ \alpha_{X} \stackrel{\underline{\omega}}{=} \alpha_{\overline{X}} = \bar{r}_{\mathcal{F}_{1}X}^{\star}$$

implying thereby that

$$\operatorname{tr}_{\mathcal{F}_{1}\overline{X}}(\overline{\phi} \circ \alpha_{\overline{X}}) = \operatorname{tr}_{\mathcal{F}_{1}X}(\alpha_{X}^{-1} \circ \phi^{*})$$

which, when combined with the initial observation, renders the contention of the

lemma manifest. From the commutative diagram

we see that

$$\alpha_{X} \stackrel{\triangle}{=} \alpha_{\overline{X}} = (\Xi^{2})^{-1} \circ \alpha \qquad \circ \Xi^{1}$$

and from the commutative diagram

we see that

$$\mathcal{F}_{2}\overline{r}_{X}^{\star} \circ \alpha = \alpha_{e} \circ \mathcal{F}_{1}\overline{r}_{X}^{\star}.$$

Recalling now that

$$\bar{r}_{f_1X} = (\Xi^1_{X,\bar{X}})^{-1} \circ F_1\bar{r}_X \circ \xi^1$$

$$\bar{r}_{f_2X} = (\Xi^2_{X,\bar{X}})^{-1} \circ F_2\bar{r}_X \circ \xi^2$$
(cf. 9.12)

we have

$$\begin{array}{l}
\bar{r}_{\mathcal{F}_{2}X}^{\star} \circ \alpha_{X} \overset{\triangle}{=} \alpha_{\bar{X}} \\
= (\xi^{2})^{-1} \circ \mathcal{F}_{2}\bar{r}_{X}^{\star} \circ \Xi^{2} \times (\Xi^{2})^{-1} \circ \alpha_{X} \overset{\triangle}{=} \bar{x}_{X}^{\bar{x}} \\
= (\xi^{2})^{-1} \circ \mathcal{F}_{2}\bar{r}_{X}^{\star} \circ \alpha_{X} \overset{\triangle}{=} \bar{x}_{X}^{\bar{x}} \\
= (\xi^{2})^{-1} \circ \alpha_{e} \circ \mathcal{F}_{1}\bar{r}_{X}^{\star} \circ \Xi^{1} \times (\bar{x}_{X}^{\bar{x}}) \\
= (\xi^{2})^{-1} \circ \mathcal{F}_{1}\bar{r}_{X}^{\star} \circ \Xi^{1} \times (\bar{x}_{X}^{\bar{x}}) \\
= (\xi^{1})^{-1} \circ \mathcal{F}_{1}\bar{r}_{X}^{\star} \circ \Xi^{1} \times (\bar{x}_{X}^{\bar{x}}) \\
= \bar{r}_{\mathcal{F}_{1}X'}^{\star}
\end{array}$$

as claimed.

The results embodied in 11.12 and 11.13 finish Step 2 of the program, which leaves Step 3 to be dealt with.

Put

$$A_{\mathcal{F}_{1},\mathcal{F}_{2}} = \bigoplus_{\mathbf{i} \in \mathcal{I}_{\underline{C}}} Mor(\mathcal{F}_{2}X_{\mathbf{i}},\mathcal{F}_{1}X_{\mathbf{i}}).$$

11.14 LEMMA The linear map

$$^{\mathrm{q}: \mathsf{A}_{\mathcal{F}_1, \mathcal{F}_2}} \rightarrow ^{\mathsf{A}(\mathcal{F}_1, \mathcal{F}_2)}$$

that sends

$$\phi_{\mathbf{i}} \in Mor(\mathcal{F}_{2}X_{\mathbf{i}}, \mathcal{F}_{1}X_{\mathbf{i}})$$

to $[X_{\mathbf{i}},\phi_{\mathbf{i}}]$ is an isomorphism of vector spaces.

<u>PROOF</u> Every $A \in A(\mathcal{F}_1, \mathcal{F}_2)$ is an $[X, \phi]$ (cf. 11.8) and every $[X, \phi]$ is a sum of elements $[X_i, \phi_i]$ with X_i irreducible. Therefore Y is surjective. That Y is injective is a consequence of the fact that

$$i \neq j => Mor(X_{i}, X_{j}) = \{0\}.$$

Put

$$A_{i} = \Psi(Mor(\mathcal{F}_{2}X_{i},\mathcal{F}_{1}X_{i})).$$

Then there is a direct sum decomposition

$$A(\mathcal{F}_1,\mathcal{F}_2) = \bigoplus_{\mathbf{i} \in \mathbf{I}_{\underline{C}}} A_{\mathbf{i}}.$$

Define a linear functional

$$\omega:A(\mathcal{F}_1,\mathcal{F}_2)\to C$$

by taking it to be zero on A_i if i does not correspond to e but on A_e , let

$$\omega([e,\phi]) = (\xi^{1})^{-1} \circ \phi \circ \xi^{2} \in C.$$

11.15 LEMMA \forall A \neq 0, ω (A*A) > 0.

PROOF Write

$$A = \sum_{i} [X_{i}, \phi_{i}],$$

where the $\mathbf{X}_{\mathbf{i}}$ are irreducible and distinct -- then

$$\mathtt{i} \neq \mathtt{j} \Rightarrow \omega([\mathtt{X}_{\mathtt{i}},\phi_{\mathtt{i}}]^{*} \cdot [\mathtt{X}_{\mathtt{j}},\phi_{\mathtt{j}}]) = 0.$$

In fact,

$$Mor(e, \vec{X}_{i} \otimes X_{j}) \approx Mor(X_{i}, X_{j}) = \{0\}$$
 (cf. 9.3),

so e is not a subobject of $\bar{X}_i \otimes X_j$. One can therefore assume that $A = [X, \phi] \neq 0$ with X irreducible. Recall now that

$$r_X^{\star} \circ r_X = \text{dim } X = n_X \text{id}_e \ (n_X \in N)$$
.

This said, let

$$p_{X} = \frac{r_{X} \circ r_{X}^{*}}{n_{X}}.$$

Then $p_X^* = p_X$ and

$$p_{X} \circ p_{X} = \frac{r_{X} \circ r_{X}^{*}}{n_{X}} \circ \frac{r_{X} \circ r_{X}^{*}}{n_{X}}$$

$$= \frac{1}{n_{X}^{2}} (r_{X} \circ r_{X}^{*} \circ r_{X} \circ r_{X}^{*})$$

$$= \frac{1}{n_{X}^{2}} r_{X} \circ n_{X} id_{e} \circ r_{X}^{*}$$

$$= \frac{r_{X} \circ r_{X}^{*}}{n_{X}} = p_{X}.$$

I.e.:

$$p_{X} \in Mor(\bar{X} \boxtimes X, \bar{X} \boxtimes X)$$

is a projection. Write

$$A^*A = [X, \phi]^* \cdot [X, \phi]$$

$$= [\overline{X}, \overline{\phi}] \cdot [X, \phi]$$

$$= [\overline{X} \boxtimes X, \Xi^{1}_{\overline{X}, X} \circ \overline{\phi} \boxtimes \phi \circ (\Xi^{2}_{\overline{X}, X})^{-1}]$$

$$= [\overline{X} \boxtimes X, \mathcal{F}_{1}(p_{X}) \circ \Xi^{1}_{\overline{X}, X} \circ \overline{\phi} \boxtimes \phi \circ (\Xi^{2}_{\overline{X}, X})^{-1}]$$

$$\begin{split} &+ [\overline{X} \boxtimes X, \mathcal{F}_{1}(\mathrm{id}_{\overline{X} \boxtimes X} - \mathrm{p}_{X}) \circ \Xi_{\overline{X}, X}^{1} \circ \overline{\phi} \ \underline{\otimes} \ \phi \circ (\Xi_{\overline{X}, X}^{2})^{-1}] \\ &= [\overline{X} \boxtimes X, \mathcal{F}_{1}(\mathrm{p}_{X}) \circ \Xi_{\overline{X}, X}^{1} \circ \overline{\phi} \ \underline{\otimes} \ \phi \circ (\Xi_{\overline{X}, X}^{2})^{-1}] \\ &= \frac{1}{n_{X}} [\mathrm{e}, \mathcal{F}_{1} r_{X}^{\star} \circ (\Xi_{\overline{X}, X}^{1}) \circ \overline{\phi} \ \underline{\otimes} \ \phi \circ (\Xi_{\overline{X}, X}^{2})^{-1} \circ \mathcal{F}_{2} r_{X}] \quad (\mathrm{cf. 11.7}) \\ &= \frac{1}{n_{X}} ((\xi^{1})^{-1} \circ \mathcal{F}_{1} r_{X}^{\star} \circ (\Xi_{\overline{X}, X}^{1}) \circ \overline{\phi} \ \underline{\otimes} \ \phi \circ (\Xi_{\overline{X}, X}^{2})^{-1}) \circ \mathcal{F}_{2} r_{X} \circ \xi^{2}) [\mathrm{e}, \xi^{1} \circ (\xi^{2})^{-1}] \\ &= \frac{1}{n_{X}} (r_{\mathcal{F}_{1} X}^{\star} \circ \overline{\phi} \ \underline{\otimes} \ \phi \circ r_{\mathcal{F}_{2} X}) [\mathrm{e}, \xi^{1} \circ (\xi^{2})^{-1}] \\ &= \frac{1}{n_{X}} (r_{\mathcal{F}_{1} X}^{\star} \circ \overline{\phi} \ \underline{\otimes} \ \phi \circ \phi^{\star}) \circ r_{\mathcal{F}_{1} X}) [\mathrm{e}, \xi^{1} \circ (\xi^{2})^{-1}] \\ &= \frac{1}{n_{Y}} (\Phi^{\star} \circ \Phi) [\mathrm{e}, \xi^{1} \circ (\xi^{2})^{-1}], \end{split}$$

where

$$\Phi = \operatorname{id}_{\mathcal{F}_{1} \overline{X}} \underline{\otimes} \, \varphi^{*} \, \circ \, r_{\mathcal{F}_{1} X}.$$

Then

$$\Phi^* \circ \Phi : \underline{e} \to \underline{e}$$

when viewed as a constant, is nonnegative. But $\phi \neq 0 \Rightarrow \Phi \neq 0$. Proof: $\bar{\phi}$ is the unique element of Mor $(\mathcal{F}_2\bar{X},\mathcal{F}_1\bar{X})$ such that

$$\overline{\phi} \overset{\boxtimes}{\cong} \operatorname{id}_{\mathcal{F}_{2}X} \circ r_{\mathcal{F}_{2}X} = \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \overset{\boxtimes}{\cong} \phi^{*} \circ r_{\mathcal{F}_{1}X} \quad (cf. 9.10),$$
so $\Phi = 0 \Rightarrow \overline{\phi} = 0$

$$\Rightarrow (^{\vee}\phi)^{*} = 0 \Rightarrow (^{\vee}\phi)^{*} = 0 \Rightarrow {}^{\vee}\phi = 0 \Rightarrow \phi = 0.$$

[Note: To justify the equation

$$\overline{\phi} \overset{\underline{\omega}}{=} \phi \circ r_{\mathcal{F}_{\mathbf{2}}X} = \mathrm{id}_{\mathcal{F}_{\mathbf{1}}\overline{X}} \overset{\underline{\omega}}{=} (\phi \circ \phi^{*}) \circ r_{\mathcal{F}_{\mathbf{1}}X'}$$

write

$$\bar{\phi} \ \underline{\omega} \ \phi = \mathrm{id}_{\mathcal{F}_1 \bar{X}} \ \underline{\omega} \ \phi \circ \bar{\phi} \ \underline{\omega} \ \mathrm{id}_{\mathcal{F}_2 X}.$$

Then

$$\begin{array}{l} \overline{\phi} \ \underline{\otimes} \ \phi \ \circ \ r_{\mathcal{F}_{2}X} = \ \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \ \underline{\otimes} \ \phi \ \circ \ \overline{\phi} \ \underline{\otimes} \ \operatorname{id}_{\mathcal{F}_{2}X} \circ \ r_{\mathcal{F}_{2}X} \\ \\ = \ \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \ \underline{\otimes} \ \phi \ \circ \ \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \ \underline{\otimes} \ \phi^{\star} \circ \ r_{\mathcal{F}_{1}X} \\ \\ = \ \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \circ \ \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \ \underline{\otimes} \ (\phi \circ \phi^{\star}) \circ \ r_{\mathcal{F}_{1}X} \\ \\ = \ \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \ \underline{\otimes} \ (\phi \circ \phi^{\star}) \circ \ r_{\mathcal{F}_{1}X} \\ \\ = \ \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \ \underline{\otimes} \ (\phi \circ \phi^{\star}) \circ \ r_{\mathcal{F}_{1}X} \\ \end{array}$$

Given $A,B \in A(F_1,F_2)$, let

$$\langle A,B \rangle = \omega(A*B)$$
.

Then < , > equips $A(\mathcal{F}_1,\mathcal{F}_2)$ with the structure of a pre-Hilbert space w.r.t. which the left multiplication operators

$$A(F_1,F_2) \rightarrow A(F_1,F_2)$$

are continuous. Denoting by $\mathcal{H}(\mathcal{F}_1,\mathcal{F}_2)$ the Hilbert space completion of $\mathcal{A}(\mathcal{F}_1,\mathcal{F}_2)$, it thus follows that $\mathcal{A}(\mathcal{F}_1,\mathcal{F}_2)$ admits a faithful *-representation

$$L:A(\mathcal{F}_1,\mathcal{F}_2) \rightarrow B(\mathcal{H}(\mathcal{F}_1,\mathcal{F}_2))$$
,

hence $A(\mathcal{F}_1,\mathcal{F}_2)$ admits a C*-norm as claimed in Step 3.

§12. THE INTRINSIC GROUP

Let C be a tannakian category and suppose that

is a fiber functor — then its intrinsic group $G_{\mathfrak{F}}$ is the group of unitary monoidal natural transformations $\alpha: \mathfrak{F} \to \mathfrak{F}$, i.e., in the notation of §11,

$$G_{\mathfrak{F}} = Mor(\mathfrak{F}, \mathfrak{F})$$
,

where $Mor(\mathcal{F}, \mathcal{F})$ is computed in $ff^*(\underline{C})$.

So

$$G_{F} \subset \prod_{X \in Ob} U(FX),$$

 $\mathcal{U}(\mathit{FX})$ the compact group of unitary operators $\mathit{FX} \to \mathit{FX}$. And G_{F} is closed if

$$TT U(FX)$$

X \in Ob C

is equipped with the product topology, thus $\mathbf{G}_{\mathbf{F}}$ is a compact group.

N.B. Define

$$\pi_{X}:G_{F} \to U(FX)$$

by $\pi_{X}(\alpha) = \alpha_{X}$ -- then

$$(\pi_{X}, FX) \in Ob \ \underline{Rep} \ G_{F}.$$

12.1 <u>LEMMA</u> \exists a faithful symmetric monoidal *-preserving functor $\Phi:\underline{C} \to \underline{\text{Rep}}\ G_{F}$ such that $U \circ \Phi = F$, where

$$U: \underline{Rep} G_{\mathcal{F}} \rightarrow \underline{FDHILB}$$

is the forgetful functor.

PROOF Define Φ on objects by

$$\Phi X = (\pi_{X}, FX)$$

and on morphisms $f:X \to Y$ by $\Phi f = \mathcal{F} f$ (cf. 8.2) and take for ξ , Ξ the corresponding entities per \mathcal{F} . To see that this makes sense for Ξ say, one must check that $\Xi_{X,Y}$ is a morphism in $\underline{\operatorname{Rep}}\ G_{\mathcal{F}}$, viz.:

$$\Xi_{X,Y} \circ (\pi_X(\alpha) \stackrel{\underline{\omega}}{=} \pi_Y(\alpha)) = \pi_{X \stackrel{\underline{\omega}}{=} Y}(\alpha) \circ \Xi_{X,Y}.$$

But this is obvious since the diagram

commutes. That Φ is symmetric is equally clear.

More is true: Φ is an equivalence of categories. Because Φ is faithful, it remains to establish that Φ is full and has a representative image (details below).

12.2 REMARK The category $\operatorname{Rep}_{\mathrm{fd}} A_{\mathbf{F}}$ is a semisimple symmetric monoidal *-category which can be shown to have conjugates, thus $\operatorname{Rep}_{\mathrm{fd}} A_{\mathbf{F}}$ is "almost" tannakian. Specializing 8.14, it was pointed out in 8.16 that the " Φ " defined there is a symmetric monoidal equivalence $\underline{C} \to \operatorname{Rep}_{\mathrm{fd}} A_{\mathbf{F}}$. Denote now by $\operatorname{Rep}_{\mathrm{fd}} G_{\mathbf{F}}$ the category whose objects are the finite dimensional continuous representations

of $G_{\overline{F}}$ and whose morphisms are the intertwining operators — then the inclusion functor

$$\underline{\text{Rep}} \ G_{\mathcal{F}} \to \underline{\text{Rep}}_{\text{fd}} \ G_{\mathcal{F}}$$

is an equivalence. On the other hand, there is a canonical functor

$$\underline{\text{Rep}}_{\text{fd}} \stackrel{A_{\text{F}}}{\rightarrow} \underline{\text{Rep}}_{\text{fd}} \stackrel{G_{\text{F}}}{\rightarrow}$$

and it too is an equivalence (a nontrivial fact).

- 12.3 <u>LEMMA</u> If $X \in Ob$ <u>C</u> is irreducible, then the complex linear span of the $\pi_X(\alpha)$ ($\alpha \in G_{\mathfrak{p}}$) is dense in B(FX).
- 12.4 <u>LFMMA</u> If X,Y \in Ob <u>C</u> are irreducible and nonisomorphic, then the complex linear span of the $\pi_X(\alpha)$ \oplus $\pi_Y(\alpha)$ ($\alpha \in G_F$) is dense in B(FX) \oplus B(FY).
- 12.5 REMARK If X_1, \dots, X_n are distinct elements of $I_{\underline{C}'}$ then the complex linear span of the

$$\pi_{X_{1}}(\alpha) \oplus \cdots \oplus \pi_{X_{n}}(\alpha) \qquad (\alpha \in G_{\mathcal{F}})$$

is dense in

$$B(FX_1) \oplus \cdots \oplus B(FX_n)$$
.

To prove that Φ is full, we shall appeal to 7.9.

- (a) X irreducible => Φ X irreducible. In fact, thanks to 12.3, the only $T \in B(FX) \text{ that intertwine the } \pi_X(\alpha) \ (\alpha \in G_F) \text{ are the scalar multiples of the identity.}$
 - (b) X,Y irreducible and nonisomorphic => $\Phi X, \Phi Y$ irreducible and nonisomorphic.

For suppose that $T: FX \to FY$ intertwines π_X and π_Y , thus $T\pi_X(\alpha) = \pi_Y(\alpha)T$ ($\alpha \in G_F$). But then Tu = vT for all $u \in B(FX)$, $v \in B(FY)$ (cf. 12.4). Now take u = 0, v = 1 to conclude that T = 0, hence $\Phi X, \Phi Y$ are nonisomorphic.

The final claim is that Φ has a representative image. To see this, consider the map

defined by the rule

$$\gamma_{\mathcal{F}}(X_{\mathbf{i}}) = (\pi_{X_{\mathbf{i}}}, \mathcal{F}X_{\mathbf{i}}).$$

Then $\gamma_{\mathbf{F}}$ is injective.

12.6 LEMMA γ_F is surjective.

PROOF The complex linear span of the matrix elements of the $\pi_{X_{\hat{1}}}$ as i ranges over $I_{\underline{C}}$ is a unital *-subalgebra of $C(G_{\underline{F}})$ which separates the points of $G_{\underline{F}}$, thus is dense in $C(G_{\underline{F}})$. Accordingly, there can be no irreducible object in $\underline{\text{Rep}}\ G_{\underline{F}}$ which is not unitarily equivalent to a $\pi_{X_{\hat{1}}}$ for some i, so $\gamma_{\underline{F}}$ is surjective.

Therefore $\gamma_{\mathbf{F}}$ is bijective and Φ has a representative image.

12.7 REMARK Suppose that

$$f_1:\underline{C} \to \underline{FDHILB}$$

$$f_2:\underline{C} \to \underline{FDHILB}$$

are fiber functors — then as objects of $ff^*(C)$, F_1 , F_2 are isomorphic (cf. 11.3),

so $\mathbf{G}_{\mathbf{F}_1}, \mathbf{G}_{\mathbf{F}_2}$ are isomorphic (in the category $\underline{\mathbf{CPTGRP}}$).

Let G be a compact group -- then the forgetful functor

is a fiber functor. Define a map $\Gamma \colon\! G \to G_{\mbox{U}}$ by sending $\sigma \in G$ to the string

$$\{\pi(\sigma): (\pi, H_{\pi}) \in Ob \text{ Rep } G\}.$$

That this is meaningful follows upon noting that if

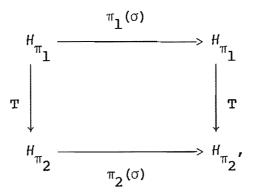
$$(\pi_{1}, H_{\pi_{1}}) \in Ob \underline{Rep} G,$$

$$(\pi_{2}, H_{\pi_{2}})$$

then

$$\forall \ \mathtt{T} \in \mathtt{Mor}((\pi_1, H_{\pi_1}), (\pi_2, H_{\pi_2}))$$

there is a commutative diagram



thus the string

$$\{\pi(\sigma): (\pi, \mathcal{H}_{\pi}) \in Ob \ \underline{Rep} \ G\}$$

defines an element

$$\alpha(\sigma) \in Mor(U,U)$$
,

where technically

$$\alpha(\sigma)_{(\pi,H_{\pi})} = \pi(\sigma).$$

12.8 LEMMA I is a continuous injective homomorphism.

[This is immediate from the definitions.]

In fact, Γ is surjective, hence ${\tt G}$ and ${\tt G}_{{\tt II}}$ are isomorphic.

[If Γ were not surjective, replace G by ΓG and think of G as a proper closed subgroup of G_U — then there would be an irreducible representation of G_U that contains a nonzero vector invariant under G but not under G_U . This, however, is impossible:

is bijective.]

12.9 THEOREM Up to isomorphism in CPTGRP, G is the "intrinsic group" of Rep G.

[If

is a fiber functor, then $G_{F} \approx G_{\overline{U}}$ (cf. 12.7).]

12.10 <u>REMARK</u> Compact groups G,G' are said to be <u>isocategorical</u> if <u>Rep</u> G, <u>Rep</u> G' are equivalent as monoidal categories. In general, this does not mean that <u>Rep</u> G, <u>Rep</u> G' are equivalent as symmetric monoidal categories and G,G' may very well be isocategorical but not isomorphic.

§13. CLASSICAL THEORY

A <u>character</u> of a commutative unital C*-algebra A is a nonzero homomorphism $\omega:A\to \mathbb{C}$ of algebras. The set of all characters of A is called the <u>structure space</u> of A and is denoted by $\Delta(A)$.

N.B. We have

$$\begin{bmatrix} - & \Delta(A) = \emptyset & (A = \{0\}) \\ & \Delta(A) \neq \emptyset & (A \neq \{0\}). \end{bmatrix}$$

13.1 LEMMA Let $\omega \in \Delta(A)$ — then ω is necessarily bounded. In fact,

$$||\omega|| = 1 = \omega(1_A).$$

N.B. The elements of $\Delta(A)$ are the pure states of A, hence, in particular, are *-homomorphisms: $\forall \ A \in A$,

$$\omega(A^*) = \overline{\omega(A)}$$
.

Given $A \in A$, define

$$\hat{A}:\Delta(A) \rightarrow C$$

by

$$\hat{\mathbf{A}}(\omega) = \omega(\mathbf{A}).$$

Equip $\Delta(A)$ with the initial topology determined by the A, i.e., equip $\Delta(A)$ with the relativised weak* topology.

13.2 LEMMA $\Delta(A)$ is a compact Hausdorff space.

If X is a compact Hausdorff space, then C(X) equipped with the supremum norm

$$||f|| = \sup_{x \in X} |f(x)|$$

and involution

$$f^*(x) = \overline{f(x)}$$

is a commutative unital C*-algebra. Moreover, \forall x \in X, the Dirac measure $\delta_{\bf x}$ \in Δ (C(X)) and the arrow

$$\begin{array}{c} - & X \to \Delta(C(X)) \\ & \times \to \delta_{X} \end{array}$$

is a homeomorphism.

13.3 LEMMA $\hat{A} \in C(\Delta(A))$ and the arrow

$$\begin{array}{c} - & A \rightarrow C(\Delta(A)) \\ & \hat{A} \rightarrow \hat{A} \end{array}$$

is a unital *-isomorphism.

N.B. If $A = \{0\}$, then $\Delta(A) = \emptyset$ and there is exactly one map $\emptyset \to \mathbb{C}$, namely the empty function $(\emptyset = \emptyset \times \mathbb{C})$, which we shall take to be 0.

Notation: Let <u>CPTSP</u> be the category whose objects are the compact Hausdorff spaces and whose morphisms are the continuous functions.

Notation: Let <u>COMUNC*ALG</u> be the category whose objects are the commutative unital C*-algebras and whose morphisms are the unital *-homomorphisms.

Let X and Y be compact Hausdorff spaces. Suppose that $\phi: X \to Y$ is a continuous function — then ϕ induces a unital *-homomorphism

$$\phi*:C(Y) \to C(X)$$
,

viz. $\phi^*(f) = f \circ \phi$. Therefore the association that sends X to C(X) defines a cofunctor

Let A and B be commutative unital C*-algebras. Suppose that $\Phi: A \to B$ is a unital *-homomorphism -- then Φ induces a continuous function

$$\Phi^*:\Delta(\mathcal{B}) \to \Delta(A)$$
,

viz. $\Phi^*(\omega) = \omega \circ \Phi$. Therefore the association that sends A to $\Delta(A)$ defines a cofunctor

$$\Delta$$
: COMUNC*ALG \rightarrow CPTSP.

13.4 <u>THEOREM</u> The category <u>CPTSP</u> is coequivalent to the category <u>COMUNC*ALG</u>.

PROOF Define

$$\Xi_X:X \to \Delta(C(X))$$

by the rule $\Xi_X(x) = \delta_x$ — then Ξ_X is a homeomorphism and there is a commutative diagram

$$\begin{array}{ccc}
X & \xrightarrow{\Xi_{X}} & \Delta(C(X)) \\
 & \downarrow & & \downarrow & \phi^{**} \\
 & Y & \xrightarrow{\Xi_{Y}} & \Delta(C(Y)).
\end{array}$$

Define

$$\Xi_A:A \rightarrow C(\Delta(A))$$

by the rule $\Xi_{A}(A) = \hat{A}$ -- then Ξ_{A} is a unital *-isomorphism and there is a commutative

diagram

$$A \xrightarrow{\Xi_{A}} C(\Delta(A))$$

$$\Phi \downarrow \qquad \qquad \downarrow \Phi^{**}$$

$$B \xrightarrow{\Xi_{B}} C(\Delta(B)).$$

Therefore

id
$$\approx \Delta \circ C$$
id $\approx C \circ \Delta$.

The category <u>CPTSP</u> has finite products with final object {*}. Therefore the category <u>COMUNC*ALG</u> has finite coproducts with initial object C. To explicate the latter, invoke the nuclearity of the objects of <u>COMUNC*ALG</u>, thus

$$A \otimes_{\max} B = A \otimes_{\min} B$$
,

call it A <u>®</u> B -- then

$$A \mid \mid B = A \otimes B$$

and there are arrows

$$\begin{bmatrix} A \rightarrow A & \underline{\omega} & B \\ A \rightarrow A & \underline{\omega} & B \end{bmatrix} \xrightarrow{B \rightarrow A} A & \underline{\omega} & B \\ B \rightarrow 1_{A} & \underline{\omega} & B.$$

13.5 EXAMPLE We have

13.6 <u>REMARK</u> Let A be a commutative unital C*-algebra — then the algebraic tensor product A @ A can be viewed as an involutive subalgebra of A @ A. Another point is this: Since A @ A is the coproduct, there is a canonical arrow A @ A \longrightarrow A with m(A @ B) = AB, i.e., the restriction of m to A @ A is the multiplication in A.

[Note: If A_1, A_2, B are commutative unital C*-algebras and if

$$\begin{bmatrix}
 \Phi_1 : A_1 \to B \\
 \Phi_2 : A_2 \to B
\end{bmatrix}$$

are unital *-homomorphisms, then the diagram

admits a unique filler

$$\Phi_1 \overset{\triangle}{=} \Phi_2 : A_1 \overset{\triangle}{=} A_2 \rightarrow \mathcal{B}$$

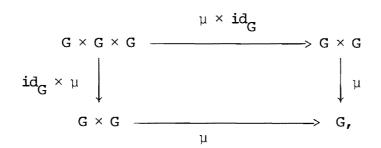
such that

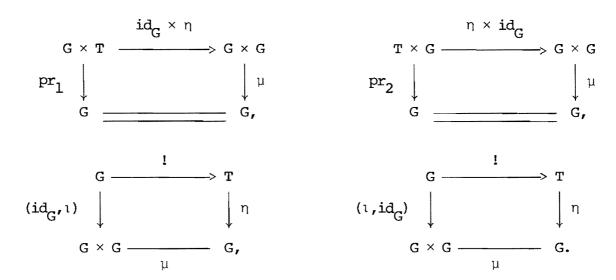
$$(\Phi_1 \ \underline{\&} \ \Phi_2) \ (A_1 \ \& \ A_2) \ = \ \Phi_1 \ (A_1) \ \Phi_2 \ (A_2) \quad (A_1 \ \in \ A_1, \ A_2 \ \in \ A_2) \ .]$$

13.7 RAPPEL Let \underline{C} be a category with finite products and final object \underline{T} -then a group object in \underline{C} consists of an object \underline{G} and morphisms

$$\mu:G \times G \rightarrow G$$
, $\eta:T \rightarrow G$, $\iota:G \rightarrow G$

such that the following diagrams commute:





There are obvious definitions of internal group homomorphism $G \to G'$, composition of internal group homomorphisms $G \to G'$, $G' \to G''$, and the identity internal group homomorphism $\mathrm{id}_{G}: G \to G$. Accordingly, there is a category $\underline{GRP}(\underline{C})$ whose objects are the group objects in C and whose morphisms are the internal group homomorphisms.

[Note: If instead \underline{C} is a category with finite coproducts and initial object I, then we put

$$\underline{\text{COGRP}}(\underline{\mathbf{C}}) = \underline{\text{GRP}}(\underline{\mathbf{C}}^{\text{OP}})^{\text{OP}}$$

and call the objects the cogroup objects in C and the morphisms the internal cogroup homomorphisms.]

13.8 EXAMPLE Take
$$\underline{C} = \underline{SET}$$
 — then
$$\underline{GRP}(\underline{SET}) = \underline{GRP}.$$

13.9 LEMMA We have

$$GRP(CPTSP) = CPTGRP.$$

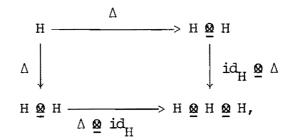
13.10 REMARK The forgetful functor

has a left adjoint. Proof: Given a set X, equip it with the discrete topology, form the associated free topological group $F_{\rm gr}\left(X\right)$, and consider its Bohr compactification.

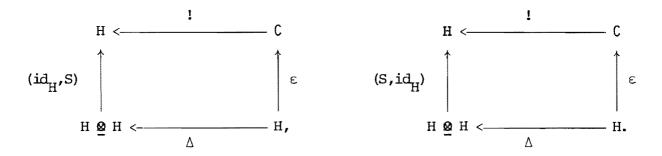
A <u>commutative Hopf C*-algebra</u> is commutative unital C*-algebra H together with unital *-homomorphisms

$$\Delta: H \rightarrow H \otimes H$$
, $\epsilon: H \rightarrow C$, $S: H \rightarrow H$

for which the following diagrams commute:







[Note: Such an H is not necessarily a Hopf algebra (in general, \triangle takes values in H Ω H rather than H Ω H).]

 $\underline{\text{N.B.}}$ Consider, e.g., (id_H,S) -- then in terms of the coproduct diagram

the arrow

$$(id_{H},S):H \otimes H \rightarrow H$$

is characterized by the condition that

$$(id_{H},S) \circ in_{1} = id_{H}$$

$$(id_{H},S) \circ in_{2} = S.$$

On the other hand, there is an arrow

$$\mathrm{id}_{\mathrm{H}} \ \underline{\underline{\omega}} \ \mathrm{S:H} \ \underline{\underline{\omega}} \ \mathrm{H} \ \rightarrow \ \mathrm{H} \ \underline{\underline{\omega}} \ \mathrm{H}$$

characterized by the condition that

And

$$m \circ id_{H} \underline{\otimes} S = (id_{H}, S).$$

Proof:

Denote by <u>COMHOPFC*AIG</u> the category whose objects are the commutative Hopf C*-algebras and whose morphisms $f:H \to H'$ are the unital *-homomorphisms such that $f \otimes f \circ \Delta = \Delta' \circ f$, $\epsilon = \epsilon' \circ f$, $f \circ S = S' \circ f$.

13.11 LEMMA We have

$$\underline{\text{COGRP}}(\underline{\text{COMUNC*ALG}}) = \underline{\text{COMHOPFC*ALG}}.$$

Let G be a compact group — then the group operations in G induce operations Δ , ϵ , S in C(G) w.r.t. which C(G) acquires the structure of a commutative Hopf C*-algebra. And the association that sends G to C(G) defines a cofunctor

Let H be a commutative Hopf C*-algebra — then the cogroup operations in H induce operations μ , η , ι in Δ (H) w.r.t. which Δ (H) acquires the structure of a compact group. And the association that sends H to Δ (H) defines a cofunctor Δ :COMHOPFC*ALG \rightarrow CPTGRP.

- 13.12 <u>THEOREM</u> The category <u>CPTGRP</u> is coequivalent to the category COMHOPFC*ALG (cf. 13.4).
- 13.13 <u>RAPPEL</u> Given a compact group G, let A(G) be its set of representative functions then A(G) is a unital *-subalgebra of C(G) and when endowed with the

restrictions of Δ , ϵ , S forms a commutative Hopf *-algebra.

[Note: Recall that A(G) is dense in C(G).]

- Let $\Delta\left(A\left(G\right)\right)$ be the set of nonzero multiplicative linear functionals on $A\left(G\right)$.
- Let $\Delta^*(A(G))$ be the set of *-preserving nonzero multiplicative linear functionals on A(G).

Then

$$\Delta$$
* (A(G)) $\subset \Delta$ (A(G))

and the containment is proper in general.

Equip $\Delta(A(G))$ (and hence $\Delta^*(A(G))$) with the topology of pointwise convergence and introduce the following operations:

(i)
$$(\omega_1 \cdot \omega_2) = (\omega_1 \overset{\triangle}{=} \omega_2) \circ \Delta$$
; (ii) $1_{A(G)} = \varepsilon$; (iii) $\omega^{-1} = \omega \circ S$.

Then $\Delta(A(G))$ is a group containing $\Delta^*(A(G))$ as a subgroup (in this connection, note that $\Delta(f^*) = \Delta(f)^*$ and $S(f^*) = S(f)^*$).

13.14 LEMMA \triangle *(A(G)) is a compact group.

13.15 THEOREM Define

ev:
$$G \rightarrow \Delta^*(A(G))$$

by

$$\operatorname{ev}(\sigma) = \delta_{\sigma} (\delta_{\sigma}(f) = f(\sigma)).$$

Then ev is an isomorphism in CPTGRP.

Let

be the forgetful functor.

13.16 LEMMA The arrow

$$\rho:A(U,U) \rightarrow A(G)$$

that sends $[\mathcal{H}_\pi^{},\varphi]$ $(\varphi\colon\!\mathcal{H}_\pi^{}\to\mathcal{H}_\pi^{})$ to the representative function

$$\sigma \rightarrow tr(\pi(\sigma)\phi)$$
 $(\sigma \in G)$

is a linear bijection.

[Note: This can be sharpened in that A(U,U) carries a canonical Hopf algebra structure which is preserved by ρ , i.e., ρ is an isomorphism of Hopf algebras.]